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வெளி ப்படை

# Open



USE



IMPROVE



EVANGELIZE

## Solaris 10 & OpenSolaris Performance, Observability & Debugging (POD)

Jim Mauro  
Principal Engineer  
Sun Microsystems, Inc  
[james.mauro@sun.com](mailto:james.mauro@sun.com)

Significant Material/Content contributions:  
Richard McDougall  
Chief Performance Architect  
VMware  
[rmc@vmware.com](mailto:rmc@vmware.com)





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## About The Authors

Jim Mauro is a Principle Engineer in the Systems Group Quality Office at Sun Microsystems, where he focuses on systems performance with real customer workloads. Jim also dabbles in performance for ZFS and Virtualization.

Richard McDougall is the Chief Performance Architect at VMware.

Richard and Jim authored **Solaris Internals: Solaris 10 and Open Solaris Kernel Architecture**.

Prentice Hall, 2006. ISBN 0-13-148209-2

Richard and Jim (with Brendan Gregg) authored **Solaris Performance and Tools: DTrace and MDB Techniques for Solaris 10 and Open Solaris**

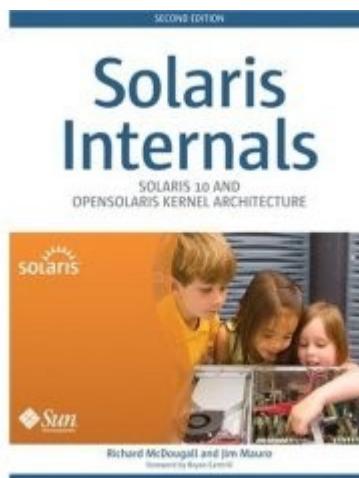
Prentice Hall, 2006. ISBN 0-13-156819-1

Richard and Jim authored **Solaris Internals:Core Kernel Architecture**,

Prentice Hall, 2001. ISBN 0-13-022496-0

## Solaris Performance and Tools

DTRACE AND MDB TECHNIQUES FOR SOLARIS 10  
AND OPENSOLARIS

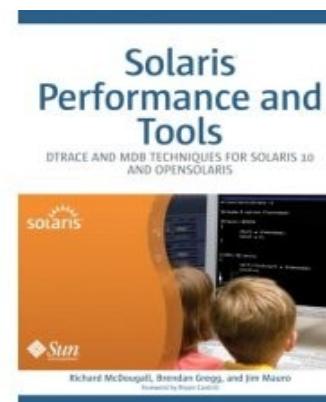
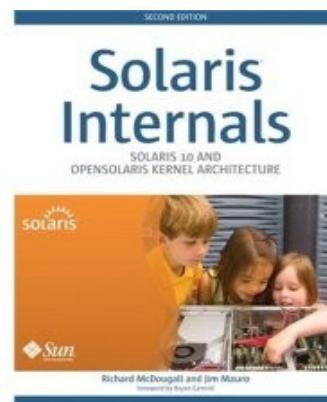


james.mauro@sun.com

rmc@vmware.com

# The Books (Shameless Plug)

- Solaris Internals: Solaris 10 and OpenSolaris Kernel Architecture
  - Community effort: over 35 contributing authors
  - Kernel data structures and algorithms
  - A lot of DTrace and mdb(1) examples to support the text
- Solaris Performance and Tools: DTrace and MDB Techniques for Solaris 10 and OpenSolaris
  - Guide to using the tools and utilities, methods, examples, etc





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# Coming Soon!

# DTrace

DYNAMIC TRACING IN SOLARIS,  
MAC OS X AND FREEBSD



Jim Mauro, Brendan Gregg, Chad Mynhier, Tariq Magdon-Ismail



# Before We Begin...

# IT DEPENDS

*What was the question...?*

**Batteries Not Included**

**Your Mileage May Vary (YMMV)**



# Solaris Performance

- Resources
  - [www.solarisinternals.com](http://www.solarisinternals.com)
    - Wikipedia of Solaris Performance
  - [www.opensolaris.org](http://www.opensolaris.org) / [www.opensolaris.com](http://www.opensolaris.com)
    - Downloads, communities, documentation, discussion groups
      - DTrace, ZFS, Performance, etc
  - [www.sun.com/bigadmin/collections/performance.html](http://www.sun.com/bigadmin/collections/performance.html)
    - Technical articles, tips, etc
  - [www.brendangregg.com](http://www.brendangregg.com)
    - DTrace Toolkit (over 230 scripts!)
    - Other Goodies
  - [www.cooltools.sunsource.net](http://www.cooltools.sunsource.net)
    - Development, optimized opensource software, analysis
  - [blogs.sun.com](http://blogs.sun.com)
    - Too much to summarize here – search for what you're interested in
  - [sunsolve.sun.com](http://sunsolve.sun.com)
    - Search for bugs related to what you're chasing...

# Agenda

- Session 1 - 9:00AM to 10:30PM
  - Goals, non goals and assumptions
  - OpenSolaris
  - Solaris 10 Kernel Overview
  - Solaris 10 Features
  - The Tools of the Trade
- Session 2 - 11:00PM to 12:30PM
  - Memory
    - Virtual Memory
    - Physical Memory
    - Memory dynamics
    - Performance and Observability
    - Memory Resource Management

# Agenda

- Session 3 - 2:00PM to 3:30PM
  - Processes, Threads & Scheduling
    - Processes, Threads, Priorities & Scheduling
    - Performance & Observability
      - Load, apps & the kernel
    - Processor Controls and Binding
    - Resource Pools, Projects & Zones
- Session 4 - 4:00PM to 5:30PM
  - File Systems and I/O
    - I/O Overview
    - UFS
    - ZFS
    - Performance & Observability
  - Network & Miscellanea



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# Session 1 - Intro, Tools, Stuff

# Goals, Non-goals & Assumptions

- Goals
  - Architectural overview of the Solaris kernel
  - The tools – what they are, what they do, when and how to use them
  - Correlate performance & observability to key functions
  - Resource control & management framework
- Non-goals
  - Detailed look at core kernel algorithms
  - Networking internals
- Assumptions
  - General familiarity with the Solaris environment
  - General familiarity with operating systems concepts

## OpenSolaris - [www.opensolaris.com](http://www.opensolaris.com)

- An open source operating system providing for community collaboration and development
- Source code released under the Common Development & Distribution License ( CDDL – pronounced “cuddle”)
- Based on “Nevada” Solaris code-base (Solaris 10+)
- New features added to OpenSolaris, then back-ported to Solaris 10
- **OpenSolaris 2008.05**
  - First supported OpenSolaris distro with many new features
    - Live CD and easy-to-use graphical installer
    - ZFS default for root
    - Network-based package management (IPS)
    - Lots of apps
- **OpenSolaris 2009.06 – current release**
  - 2010.03 next planned release (subject to change)

# Solaris 10 – Update Releases

- New features, new hardware support bug fixes
  - Check out the “What's New” Document;  
<http://docs.sun.com/app/docs/coll/1531.1?l=en>
- Solaris 10 3/05 – First release of S10
- Solaris 10 1/06 – Update 1
- Solaris 10 6/06 – Update 2
- Solaris 10 11/06 – Update 3
- Solaris 10 8/07 – Update 4
- Solaris 10 5/08 – Update 5
- Solaris 10 10/08 – Update 6
- Solaris 10 5/09 – Update 7
- Solaris 10 10/09 – Update 8



# Solaris Kernel Features

- Dynamic
- Multithreaded
- Preemptive
- Multithreaded Process Model
- Multiple Scheduling Classes
  - Including realtime support, fixed priority and fair-share scheduling
- Tightly Integrated File System & Virtual Memory
- Virtual File System
- 64-bit kernel
  - 32-bit and 64-bit application support
- Resource Management
- Service Management & Fault Handling
- Integrated Networking

# Solaris 10 & OpenSolaris

## The Headline Grabbers

- Solaris Containers (Zones)
  - Solaris Dynamic Tracing (DTrace)
  - Predictive Self Healing
    - System Management Framework (SMF)
    - Fault Management Architecture (FMA)
  - Process Rights Management (aka Least Privilege)
  - Premier x86 support
    - Optimized 64-bit Opteron support (x64)
    - Optimized Intel support
  - Zetabyte Filesystem (ZFS)
  - Network Stack – TCP/IP, UDP/IP, GLDv3
- ... and much, much more!



# Solaris Performance Optimizations

- Threads, threads, threads...
- FULLY 64-Bit
- CPU Scalability
- UFS, NFS & ZFS File System Performance
- JVM Optimizations
- Memory Scalability
- Networking
- Resource Management
- Compiler Technology
- Observability: DTrace, Analyzer, and Perf Tools

# Threads, threads, threads....

- Solaris's big-bet: large numbers of CPUs
  - SMP systems in the early days
  - Now thread-rich multicore CPUs w/hardware strands
    - 256-way T5440 (4 N2 sockets, each 8x8), 256GB RAM
    - 512-way DC3-64 (64 sockets, 256 cores, 512 strands)
- Fully preemptable kernel
- Architected around threads:
  - Kernel threads scheduled, executed
- I/O is thread-rich:
  - (Blocking) Interrupts are handled as threads
  - Worker-pools (tasks-queues) of threads avail for driving I/O
- Rich thread-development environment

# The 64-bit Revolution

Solaris 2.6

ILP32 Apps

ILP32 Libs

ILP32 Kernel

ILP32 Drivers

32-bit HW

Solaris 7, 8, 9, 10, OpenSolaris

ILP32 Apps

LP64 Apps

ILP32 Libs

LP64 Libs

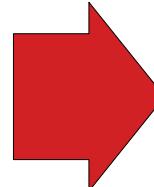
ILP32 Kernel

LP64 Kernel

ILP32 Drivers

LP64 Drivers

64-bit HW  
(SPARC, X64)



# CPU Scalability

- Per-cpu dispatcher queues
  - Fine-grained locking on thread enqueue/dequeue
- Slab / vmem allocator in the kernel
  - Adopted by most other major kernels
  - Ported to user-land – libumem.so (scalable malloc(3C))
- CMP & CMT Optimizations
  - Chip MultiProcessing/Chip MultiThreading
    - Multi-strand, Multi-core designs
    - Optimize thread placement on cores
- NUMA Optimizations (MPO)
  - Locality groups (CPUs and Memory)

# Scheduler Enhancements

- FX – Fixed Priority Scheduler
  - Integrated into Solaris 9
  - Provides fixed quantum scheduling
  - Fixed priorities
  - Eliminates unnecessary context switches for server-style apps
  - Recommend setting as the default for Databases/Oracle
- FSS – Fair Share Schedule
  - Integrated into Solaris 9
  - Replaces SRM 1.X
  - Shares of CPU allocated
  - Adds Projects and Tasks for administration / management

# File System Performance

- UFS & Databases
  - Direct I/O enables scalable database performance
  - Enhanced Logging Support introduced in S9
- NFS
  - Fireengine + RPC optimizations provide high throughput:
    - 108MB/s on GBE, 910MB/s on 10GBE, Solaris 10, x64
  - NFS for Databases Optimizations
    - 50,000+ Database I/O's per second via Direct I/O
- ZFS
  - Adaptive Replacement Cache (ARC)
  - Dynamic space management for metadata and data
  - Copy-On-Write (COW) – in-place data is never overwritten
  - Still evolving - new features and performance enhancements

# Java VM Performance

- Java SE 6
  - Lock enhancements
  - GC improvements
  - A lot more;
    - [http://java.sun.com/performance/reference/whitepapers/6\\_performance.html](http://java.sun.com/performance/reference/whitepapers/6_performance.html)
- DTrace & Java
  - jstack() (Java 5)
    - jstackstrsize for more buffer space
  - dvm provider
    - Java 1.4.2 (libdvm.so)
    - Java 1.5 (libdvm.so)
    - <https://solaris10-dtrace-vm-agents.dev.java.net>
  - Integrated HotSpot provider in Java 6
    - All DVM probes, plus extensions
  - Additional DTrace probes coming in Java 7

# Memory Scalability

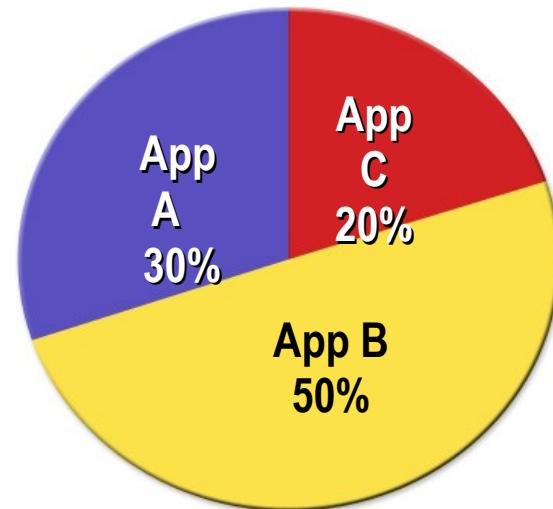
- Large Memory Optimizations
  - Solaris 9 & 10
  - 1TB shipping today. 4TB coming soon
  - 64GB hardly considered large anymore...
- Large Page Support
  - Evolved since Solaris 2.6
    - Large (4MB) pages with ISM/DISM for shared memory
  - Solaris 9/10
    - Multiple Page Size Support (MPSS)
      - Optional large pages for heap/stack
      - Programmatically via madvise()
      - Shared library for existing binaries (LD\_PRELOAD)
      - Tool to observe potential gains
        - # trapstat -t
  - Solaris 10 Updates and OpenSolaris
    - Large Pages Out-Of-The-Box (LPOOB)

# Networking

- Fire-engine in Solaris 10
  - New “vertical perimiter” scaling model
  - 9Gbits/s on 10GBE, @50% of a 2-way x64 system
- Application to application round trip latency close to 40usec
- Nemo: High performance drivers in Solaris 1 Update 2
  - GLDv3 NIC Driver Interface
  - Enables multiple-ring support
  - Generic VLan and Trunking Support
- Yosemite: High performance UDP
  - Enabled in Solaris 10 Update 2
- IP Instances
  - Unique per-Zone TCP/IP stack
- Crossbow
  - Virtualization – VNICs
  - Resource management

# Resource Management

- CPU & Network Resource Control Framework
  - Often used to control resource allocation to a container/zone.
- Processor Sets/Pools
  - Partition the machine into sets of processor resources
- Fair-Share Scheduler

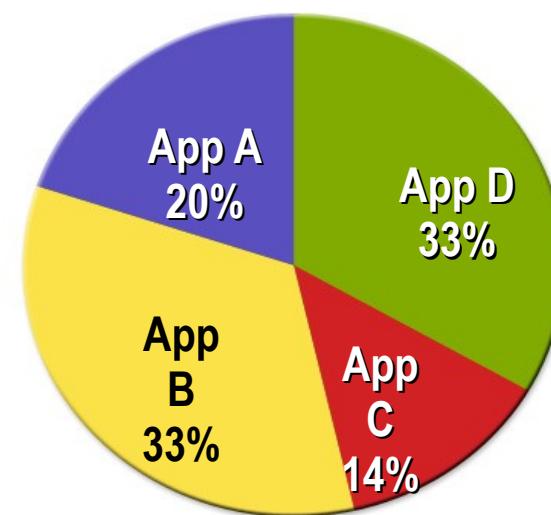


 App A (3 shares)

 App B (5 shares)

 App C (2 shares)

 App D (5 shares)



# Compiler Technology

- Studio Compilers performance optimizations
  - Highest level of optimized code for SPARC & x64
  - Studio 11 and Studio 12 **available free**
    - Outstanding profiler/performance analysis tool
    - collect(1) for collecting experiments
    - analyzer GUI, er\_print(1) for command line
  - sgcc – gcc front-end, with an optimized SPARC code generator
    - Use cc(1) or gcc(1) to build apps
  - Analyzer Tool

# Observability

- Performance Observability is the primary means to optimization
  - “That which is not measured, gets worse”
- Key Observability Features:
  - DTrace
  - Thread Microstate Measurements
  - \*stat tools, ptools, kstats, truss, prstat
  - Lock Analysis Tools: lockstat, plockstat
  - System Statistics: sar, mpstat, vmstat etc...

# Why Performance, Observability & Debugging?

- Reality, what a concept
  - Chasing performance problems
    - Sometimes they are even well defined
  - Chasing pathological behaviour
    - My app should be doing X, but it's doing Y
      - It's only doing it sometimes
- Understand utilization
  - Resource consumption
    - CPU, Memory, IO (Disk and Network)
  - Capacity planning
- In general, attaining a good understanding of the system, the workload, and how they interact
- 90% of system activity falls into one of the above categories, for a variety of roles
  - Admins, DBA's, Developers, etc...



# Before You Begin...

*“Would you tell me, please, which way I ought to go from here?” asked Alice*

*“That depends a good deal on where you want to get to” said the Cat*

*“I don't much care where...” said Alice*

*“Then it doesn't matter which way you go” said the Cat*

**Lewis Carroll**

***Alice's Adventures in Wonderland***

# General Methods & Approaches

- Define the problem
  - In terms of a business metric
  - Something measurable
- System View
  - Resource usage/utilization
    - CPU, Memory, Network, IO
- Process View
  - Execution profile
    - Where's the time being spent
  - May lead to a thread view
- Drill down depends on observations & goals
  - The path to root-cause has many forks
  - “bottlenecks” move
    - Moving to the next knee-in-the-curve

# The Utilization Conundrum

- What is utilization?
  - The most popular metric on the planet for determining if something on your system is potentially a bottleneck or out of capacity
  - Properly defined as the amount of time something is busy relative to wall clock (elapsed) time
    - N is busy for .3 seconds over 1 second sampling periods, it's  $((0.30 / 1) * 100)$  30% utilized
- So... What's the problem?
  - Basic utilization metrics assume simple devices capable of only doing 1 thing at a time
    - Old disks, old networks (NICs), old CPUs
  - Bottom Line – 100% utilized is NOT necessarily a pain point



# The Utilization Conundrum (cont)

- Modern Times
  - Disks, CPUs, NICs are all very sophisticated, with concurrency built-in at the lowest levels
    - Disks – integrated controllers with deep queues and NCQ
    - NICs – multiple ports, multiple IO channels per port
- Case in point, iostat "%b"
  - We've been ignoring it for years – it's meaningless because it simply means that an IO thread is in the disks queue every time it looks
  - "100% busy" Disks, or NICs, may be able to do more work with acceptable latency
  - It's all about meeting business requirements

# The Utilization Conundrum (cont)

- CPUs
  - Multicore. Multiple execution cores on a chip
  - Multithread (hyperthreads) – multiple threads-per-core
  - CMT – Chip Multithreading
    - Combining multicore and multithread.
- CPU Utilization
  - Each thread (or strand) appears as a CPU to Solaris
  - Each CPU maintains its own set of utilization metrics
    - Derived from CPU microstates – sys, usr, idle
    - Multiple threads sharing the same core can each appear 100% utilized
- A CPU that shows 100% utilization (0% idle) has about as much meaning as a disk or NIC that shows 100% utilization
  - More to the point, a CPU that is observed to be 100% utilized may be capable of doing more work without a tradeoff in latency
    - e.g. a multi-execution unit pipeline running 1 thread all the time is 100% utilized, but capable of running another thread while maintaining the same service level

Google “*Utilization is Virtually Useless as a Metric*”

# CPU Utilization

- Traditional “stat” tools
  - threads are CPUs
  - CPU microstates
- corestat
  - Unbundled script that uses cpustat(1)
    - cpustat(1) programs hardware counters (PICs) to gather chip statistics
      - Very hardware-specific
  - corestat reports and vmstat/mpstat reports may vary due to the very different methods of data gathering



## vmstat

```

kthr      memory          page          disk      faults      cpu
r b w    swap   free  re mf pi po fr de sr m1 m2 m2  in   sy   cs us sy id
0 11 0 83283928 30844904 10682 15143 42254 121 106 0 0 59 0 60 0 18928 145389 24371 52 21 27
0 16 0 83250248 30829608 9300 13420 37344 27797 27789 0 0 170 0 168 0 20669 149382 26695 49 22 29
0 11 0 83237872 30848544 10700 13440 28631 76 76 0 0 170 0 169 0 20531 158387 27474 42 21 37
0 13 0 83248120 30840824 11433 14640 42167 23755 23755 0 0 24 0 25 0 20447 139348 26573 47 27 26
0 21 0 83207744 30843816 11944 15716 50408 19854 19854 0 0 47 0 46 0 20261 148615 25345 51 26 23
0 11 0 83231536 30861544 10106 15260 90549 811 811 0 0 43 0 43 0 17916 170772 23022 50 25 26
0 14 0 83241696 30866376 7904 12898 27841 100 100 0 0 36 0 36 0 16016 168074 20462 45 20 35
0 15 0 83257208 30888504 9293 15014 28643 10077 10077 0 0 41 0 41 0 15377 120768 19819 32 17 51
0 12 0 83219888 30894256 8621 16163 28433 27427 27419 0 0 30 0 30 0 14892 144676 17142 44 28 28
0 14 0 83225824 30911520 8931 15740 26061 7036 7036 0 0 48 0 48 0 17729 143259 23048 37 20 43
0 14 0 83241384 30914304 9709 14068 32892 177 177 0 0 42 0 42 0 16522 144625 21967 36 19 45
0 12 0 83239696 30931664 9197 16691 31077 13915 13892 0 0 34 0 34 0 15312 141997 19194 42 19 39
4 9 0 83255376 30947384 8979 15789 31323 54 54 0 0 34 0 34 0 15345 147188 19970 44 28 29
2 12 0 83280600 30962944 9593 11575 27278 162 162 0 0 54 0 54 0 14589 129520 18370 37 16 47
0 10 0 83279888 30960968 10216 11699 36746 1385 1385 0 0 35 0 35 0 13142 135268 16318 31 18 51
0 13 0 83289392 30969488 11409 13129 44079 0 0 0 32 0 32 0 13245 130531 16526 24 19 56
0 6 0 83304648 30977048 10653 10764 38600 162 162 0 0 30 0 29 0 12565 113020 15400 23 13 64
0 12 0 83293552 30967384 10505 14347 33439 8 8 0 0 42 0 43 0 13933 116965 17959 27 15 57
0 9 0 83316160 30977416 9349 11402 28741 447 447 0 0 35 0 35 0 13688 91570 17723 29 16 55

```

**avg idle – 40.42%**

**avg sys - 20.53%**

**avg usr - 39.05%**

## corestat

CPU Utilization Mon May 11 11:00:58 2009				
CPU	(Thd)	%Usr	%Sys	%Usr+Sys
0	(0,1)	28.56	30.77	59.33
1	(2,3)	28.59	23.12	51.71
4	(8,9)	45.72	21.97	67.68
5	(10,11)	18.43	31.62	50.05
8	(16,17)	13.51	32.17	45.69
9	(18,19)	17.99	16.03	34.02
12	(24,25)	25.04	27.74	52.78
13	(26,27)	13.48	18.45	31.93
16	(32,33)	54.45	33.69	88.13
17	(34,35)	47.48	24.56	72.05
20	(40,41)	42.69	43.48	86.17
21	(42,43)	25.89	39.09	64.98
24	(48,49)	63.57	35.05	98.62
25	(50,51)	23.27	29.89	53.17
28	(56,57)	71.71	43.87	100.00
29	(58,59)	33.65	19.33	52.98
	Avg	34.63	29.43	64.05

	vmstat	corestat
%sys	20.53	29.43
%usr	39.05	34.63
%idle	40.42	35.95



## vmstat

```

kthr      memory          page          disk        faults       cpu
r b w  swap   free  mf pi po fr de sr m1 m2 m2  in   sy   cs us sy id
0 8 0 83293328 30951456 8724 12802 30608 4000 4000 0 0 25 0 25 0 13271 87958 17034 28 14 58
0 6 0 83298144 30966096 9688 10972 30122 25454 25438 0 0 47 0 47 0 15283 115241 19714 28 17 55
0 8 0 83307976 30980096 8732 11147 30765 2521 2444 0 0 39 0 39 0 13151 96896 17352 26 13 61
0 10 0 83301456 30967616 8312 9286 29132 255 255 0 0 27 0 27 0 12856 89560 15560 27 20 53
0 12 0 83295096 30956080 8820 8621 29532 14728 14728 0 0 50 0 51 0 13639 123786 16865 32 13 55
0 13 0 83255472 30956744 8936 10414 28178 23928 23920 0 0 32 0 31 0 15620 142711 20481 31 20 49
0 15 0 83215552 30959632 9234 9270 37623 21136 21128 0 0 31 0 31 0 17276 140842 22307 35 20 45
0 14 0 83191296 30951064 9249 12303 40026 185 185 0 0 28 0 28 0 16325 143003 20585 40 16 43
0 13 0 83184352 30937424 8859 9732 37956 1182 1182 0 0 30 0 30 0 15235 151314 17797 34 26 39
0 16 0 83208648 30960648 9249 10079 35849 23026 23026 0 0 29 0 29 0 16332 128670 20900 32 17 51
0 15 0 83185728 30980040 9866 6413 39944 11595 11580 0 0 17 0 17 0 17262 103816 22081 25 14 60
0 17 0 83180464 30968800 11455 6607 115326 15 15 0 0 22 0 22 0 15087 92050 19142 25 15 60
0 13 0 83186320 30963096 12146 6460 63788 46 46 0 0 20 0 20 0 15810 82579 20030 25 17 58
0 15 0 83184984 30942064 12559 11172 60716 23 23 0 0 31 0 32 0 15018 82876 17789 24 15 61
0 16 0 83192240 30920472 10457 7314 48095 2990 2990 0 0 39 0 38 0 14787 97524 18930 21 15 64
0 18 0 83190632 30917808 9801 9156 44190 529 529 0 0 30 0 30 0 15252 110036 19890 19 18 63
0 10 0 83178232 30890776 9107 7755 46624 272 272 0 0 23 0 23 0 15567 84655 17958 25 14 61
0 11 0 83197376 30911640 9985 13448 42818 12116 12110 0 0 40 0 40 0 15324 101347 18084 29 23 48
0 9 0 83187208 30916064 9256 12272 30049 13563 13555 0 0 34 0 34 0 12918 99437 15620 26 15 58

```

**avg idle – 54.84%**

**avg sys - 16.95%**

**avg usr - 28.00%**

## corestat

CPU Utilization Mon May 11 11:02:58 2009				
CPU	(Thd)	%Usr	%Sys	%Usr+Sys
0	(0,1)	22.85	19.04	41.89
1	(2,3)	8.76	15.82	24.58
4	(8,9)	6.68	24.22	30.91
5	(10,11)	88.10	12.89	100.00
8	(16,17)	61.92	20.17	82.10
9	(18,19)	4.17	18.45	22.61
12	(24,25)	60.86	14.97	75.84
13	(26,27)	2.79	14.66	17.45
16	(32,33)	62.64	17.38	80.02
17	(34,35)	11.88	18.07	29.95
20	(40,41)	24.16	29.25	53.41
21	(42,43)	19.97	21.51	41.48
24	(48,49)	38.06	26.48	64.54
25	(50,51)	12.46	31.90	44.36
28	(56,57)	40.40	26.78	67.18
29	(58,59)	33.50	35.24	68.75
	Avg	31.20	21.68	52.88

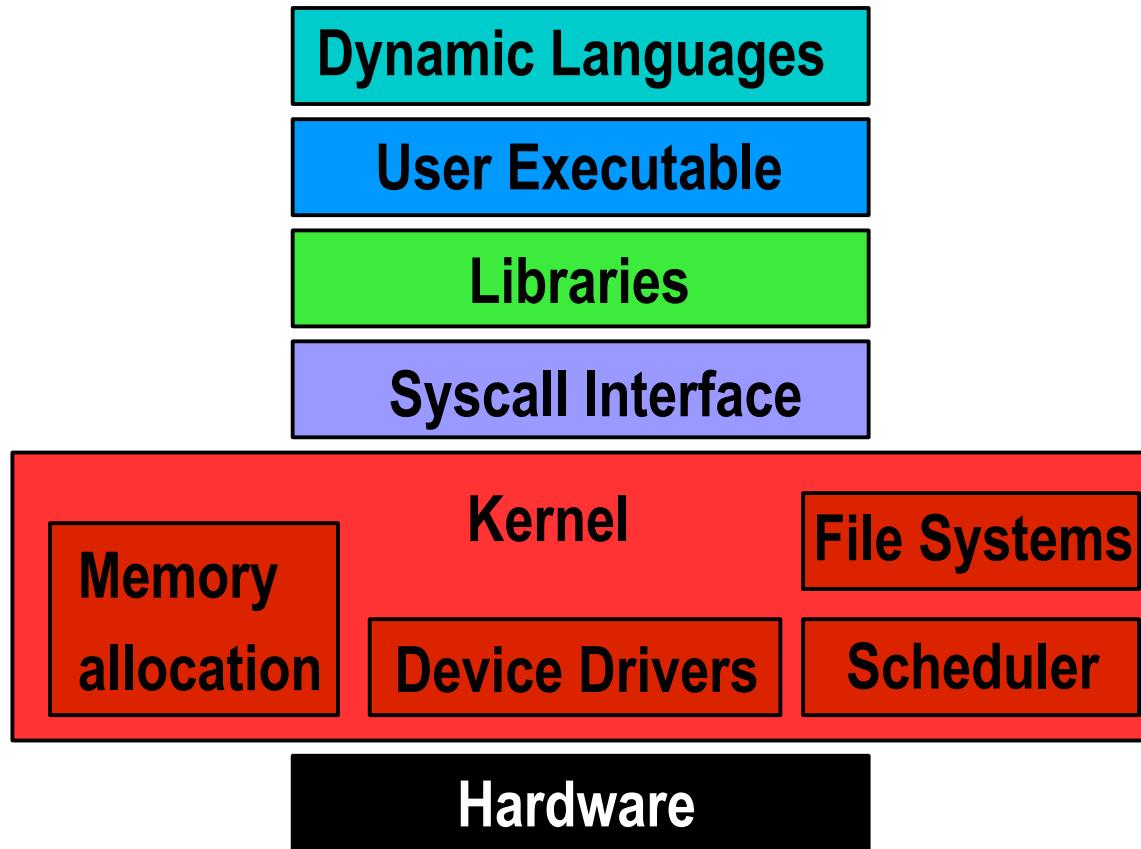
	vmstat	corestat
%sys	16.95	21.68
%usr	28.00	31.20
%idle	54.84	47.12

# CPU Utilization/Capacity

- vmstat/mpstat and corestat will vary depending on the load
  - corestat will generally be more accurate
- Use “prstat -m” LAT category, in conjunction with utilization measurements, delivered workload throughput and run queue depth (vmstat “r” column) to determine for CPU capacity planning
-

# The Workload Stack

- All stack layers are observable



## Little's Law

- A relatively simple queueing theory theorem that relates response time to throughput
  - The throughput of a system (Q) is a factor of the rate of incoming work (N), and the average amount of time required to complete the work (R – response time)
  - Independent of any underlying probability distribution for the arrival of work or the performance of work

throughput = arrival rate / avg processing time ... or

$$Q = N / R$$

e.g

if  $N = 100$  and  $R = 1$  second,  $Q = 100$  TPS

*More compelling, it makes it easy to see how these critical performance metrics relate to each other....*



# Amdahl's Law

- In general terms, defines the expected speedup of a system when part of the system is improved
- As applied to multiprocessor systems, describes the expected speedup when a unit of work is parallelized
  - Factors in degree of parallelization

$$S = \frac{1}{\left(F + \frac{(1-F)}{N}\right)}$$

S is the speedup  
F is the fraction of the work that is serialized  
N is the number of processors

$$S = \frac{1}{\left(0.5 + \frac{(1-0.5)}{4}\right)}$$

S = 1.6     4 processors,  $\frac{1}{2}$  of the work is serialized

$$S = \frac{1}{\left(0.25 + \frac{(1-0.25)}{4}\right)}$$

S = 2.3     4 processors,  $\frac{1}{4}$  of the work is serialized



USE



IMPROVE



EVANGELIZE

# Performance & Observability Tools



# Solaris Performance and Tracing Tools

## Process stats

- cputrack / cpustat – processor hw counters
- flockstat – process locks
- pargs – process arguments
- pflags – process flags
- pcred – process credentials
- pldd – process's library dependencies
- psig – process signal disposition
- pstack – process stack dump
- pmap – process memory map
- pfiles – open files and names
- prstat – process statistics
- ptree – process tree
- ptime – process microstate times
- pwdx – process working directory

## Process control

- pgrep – grep for processes
- pkill – kill processes list
- pstop – stop processes
- prun – start processes
- prctl – view/set process resources
- pwait – wait for process
- preap\* – reap a zombie process

## System Stats

- acctcom – process accounting
- busstat – Bus hardware counters
- cpustat – CPU hardware counters
- iostat – IO & NFS statistics
- kstat – display kernel statistics
- mpstat – processor statistics
- netstat – network statistics
- nfsstat – nfs server stats
- sar – kitchen sink utility
- vmstat – virtual memory stats

## Process Tracing/ debugging

- abitrace – trace ABI interfaces
- dtrace – trace the world
- mdb – debug/control processes
- truss – trace functions and system calls

## Kernel Tracing/ debugging

- dtrace – trace and monitor kernel
- lockstat – monitor locking statistics
- lockstat -k – profile kernel
- mdb – debug live and kernel cores

\*why did Harry Cooper & Ben wish they had preap?

# Solaris Dynamic Tracing - DTrace

“ *[expletive deleted] It's like they saw inside my head and gave me The One True Tool.* ”

- A Slashdotter, in a post referring to DTrace

“ *With DTrace, I can walk into a room of hardened technologists and get them giggling* ”

- Bryan Cantrill, Inventor of DTrace

# DTrace

## Solaris Dynamic Tracing – An Observability Revolution

- Ease-of-use and instant gratification engenders serious hypothesis testing
- Instrumentation directed by high-level control language (not unlike AWK or C) for easy scripting and command line use
- Comprehensive probe coverage and powerful data management allow for concise answers to arbitrary questions

# What is DTrace

- DTrace is a dynamic troubleshooting and analysis tool first introduced in the Solaris 10 and OpenSolaris operating systems.
- DTrace is many things, in particular:
  - A tool
  - A programming language interpreter
  - An instrumentation framework
- DTrace provides observability across the entire software stack from one tool. This allows you to examine software execution like never before.

# The Entire Software Stack

- How did you analyze these?

Examples:

Dynamic Languages

Java, JavaScript, ...

User Executable

native code, /usr/bin/\*

Libraries

/usr/lib/\*

Syscall Interface

man -s2

Kernel

File Systems

Memory allocation

Device Drivers

Scheduler

VFS, DNLC, UFS,

ZFS, TCP, IP, ...

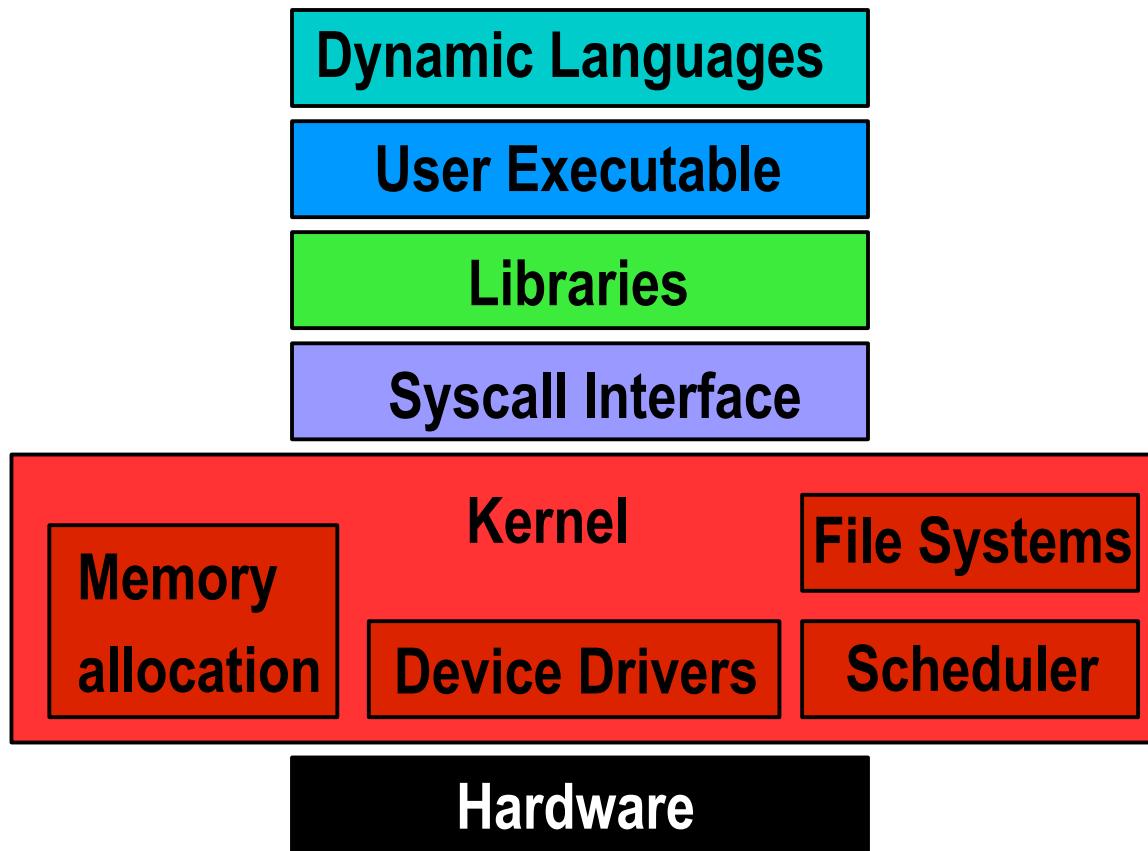
sd, st, hme, eri, ...

Hardware

NIC, Disk HBA, Processors, etc

# The Entire Software Stack

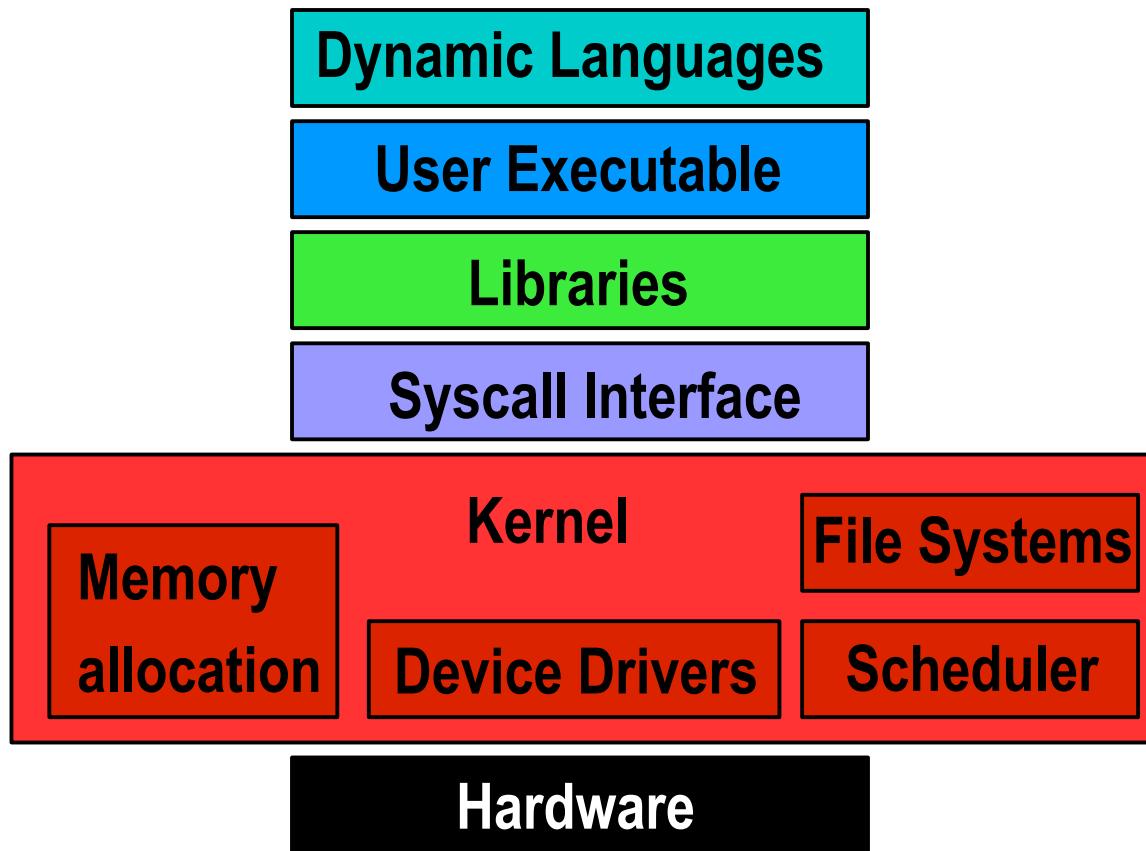
- It was possible, but difficult.



Previously:  
debuggers  
`truss -ua.out`  
`apptrace, sotruss`  
`truss`  
`prex; tnf*`  
`lockstat`  
`mdb`  
`kstat, PICs, guesswork`

# The Entire Software Stack

- DTrace is all seeing:



**DTrace visibility:**

Yes, with providers

Yes

Yes

Yes

Yes

No. Indirectly, yes

# What DTrace is like

- DTrace has the combined capabilities of numerous previous tools and more,

Tool	Capability
<b>truss -u a.out</b>	tracing user functions
<b>apptrace</b>	tracing library calls
<b>truss</b>	tracing system calls
<b>prex; tnf*</b>	tracing some kernel functions
<b>lockstat</b>	profiling the kernel
<b>mdb -k</b>	accessing kernel VM
<b>mdb -p</b>	accessing process VM

Plus a programming language similar to C and awk.



# Syscall Example

- Using truss,

Only examine 1 process

```
$ truss date
```

execve("/usr/bin/date", 0x08047C9C, 0x08047CA4) argc = 1  
resolvepath("/usr/lib/ld.so.1", "/lib/ld.so.1", 1023) = 12  
resolvepath("/usr/bin/date", "/usr/bin/date", 1023) = 13  
xstat(2, "/usr/bin/date", 0x08047A58) = 0  
open("/var/ld/ld.config", O\_RDONLY) = 3  
fxstat(2, 3, 0x08047988) = 0  
mmap(0x00000000, 152, PROT\_READ, MAP\_SHARED, 3, 0) = 0xFEFB0000  
close(3) = 0  
mmap(0x00000000, 4096, PROT\_READ|PROT\_WRITE|PROT\_EXEC, MAP\_PRIVATE|MAP\_ANON, -1  
sysconfig(\_CONFIG\_PAGESIZE) = 4096

[...]

*truss slows down the target*

Output is limited to provided options



# Syscall Example

- Using DTrace,

You choose the output

```
# dtrace -n 'syscall:::entry { printf("%16s %x %x", execname, arg0, arg1); }'
dtrace: description 'syscall:::entry' matched 233 probes

CPU      ID          FUNCTION:NAME
        1  75943        read:entry           Xorg f 8047130
        1  76211        setitimer:entry       Xorg 0 8047610
        1  76143        writev:entry         Xorg 22 80477f8
        1  76255        pollsys:entry        Xorg 8046da0 1a
        1  75943        read:entry           Xorg 22 85121b0
        1  76035        ioctl:entry          soffice.bin 6 5301
        1  76035        ioctl:entry          soffice.bin 6 5301
        1  76255        pollsys:entry        soffice.bin 8047530 2
```

[...]

Minimum performance cost

Watch every process



# What is DTrace for

- Troubleshooting software bugs
  - Proving what the problem is, and isn't.
  - Measuring the magnitude of the problem.
- Detailed observability
  - Observing devices, such as disk or network activity.
  - Observing applications, whether they are from Sun, 3<sup>rd</sup> party, or in-house.
- Capturing profiling data for performance analysis
  - If there is latency somewhere, DTrace can find it

# Terminology

- Example #1



```
# dtrace -n 'syscall::exece:return { trace(execname); }'
dtrace: description 'syscall::exece:return' matched 1 probe
CPU      ID          FUNCTION:NAME
 0    76044        exece:return  man
 0    76044        exece:return  sh
 0    76044        exece:return  neqn
 0    76044        exece:return  tbl
 0    76044        exece:return  nroff
[...]
```

# Consumer

- Consumers of libdtrace(3LIB),

dtrace	command line and scripting interface
lockstat	kernel lock statistics
plockstat	user-level lock statistics
intrstat	run-time interrupt statistics
- libdtrace is currently a private interface and not to be used directly (nor is there any great reason to); the supported interface is dtrace(1M).
  - NOTE: You are still encouraged to use libkstat(3LIB) and proc(4) directly, rather than wrapping /usr/bin consumers.

# Privileges

```
$ id  
uid=1001(user1) gid=1(other)  
$ /usr/sbin/dtrace -n 'syscall::exece:return'  
dtrace: failed to initialize dtrace: DTrace requires additional privileges
```

- Non-root users need certain DTrace privileges to be able to use DTrace.
- These privileges are from the Solaris 10 “Least Privilege” feature.

```
root::::auths=solaris.*;solaris.grant;profiles=Web Console Management,All;lock_after_retries=no  
mauroj::::defaultpriv=basic,dtrace_user,dtrace_proc,dtrace_kernel,proc_prioctl,  
proc_clock_highres;project=laptop
```



# Probes

- Data is generated from instrumentation points called “probes”.
- DTrace provides thousands of probes.
- Probe examples:

Probe Name	Description
syscall::read:entry	A read() syscall began
proc:::exec-success	A process created successfully
io:::start	An I/O was issued (disk)
io:::done	An I/O completed

# Probe Names

- Probe names are a four-tuple,

Provider	Module	Function	Name
			
<b>syscall::exec::return</b>			

- Provider A library of related probes.
- Module The module the function belongs to, either a kernel module or user segment.
- Function The function name that contains the probe.
- Name The name of the probe.



# Listing Probes

- `dtrace -l` lists all currently available probes that you have privilege to see, with one probe per line,
- Here the root user sees 69,879 available probes.
- The probe count changes – it is dynamic (DTrace).

```
# dtrace -l
      ID  PROVIDER          MODULE           FUNCTION NAME
      1    dtrace
      2    dtrace
      3    dtrace
      4    sched            FX                fx_yield schedctl-yi
[...]
# dtrace -l | wc -l
69880
```



# Tracing Probes

- `dtrace -n` takes a probe name and enables tracing,

```
# dtrace -n syscall::exece:return
dtrace: description 'syscall::exece:return' matched 1 probe
CPU      ID          FUNCTION:NAME
  0    76044          exece:return
  0    76044          exece:return
^C
```

- The default output contains,
  - CPU            CPU id that event occurred on (if this changes, the output may be shuffled)
  - ID            DTrace numeric probe id
  - FUNCTION:NAME        Part of the probe name

# Providers

- Examples of providers,

Provider	Description
<b>syscall</b>	system call entries and returns
<b>proc</b>	process and thread events
<b>sched</b>	kernel scheduling events
<b>sysinfo</b>	system statistic events
<b>vminfo</b>	virtual memory events
<b>io</b>	system I/O events
<b>profile</b>	fixed rate sampling
<b>pid</b>	user-level tracing
<b>fbt</b>	raw kernel tracing



# Providers

- Example of probes,

Provider	Example probe
syscall	<code>syscall::read:entry</code>
proc	<code>proc:::exec-success</code>
sched	<code>sched::on-cpu</code>
sysinfo	<code>sysinfo:::readch</code>
vminfo	<code>vminfo:::maj_fault</code>
io	<code>io:::start</code>
profile	<code>profile:::profile-1000hz</code>
pid	<code>pid172:libc:fopen:entry</code> <code>pid172:a.out:main:entry</code>
fbt	<code>fbt::bdev_strategy:entry</code>

# Providers

- Providers are documented in the DTrace Guide as separate chapters.
- Providers are dynamic; the number of available probes can vary.
- Some providers are “unstable interface”, such as fbt and sdt.
  - This means that their probes, while useful, may vary in name and arguments between Solaris versions.
  - Try to use stable providers instead (if possible).
  - Test D scripts that use unstable providers across target Solaris releases

# Provider Documentation

- Some providers assume a little background knowledge, other providers assume a lot. Knowing where to find supporting documentation is important.
- Where do you find documentation on -
  - Syscalls?
  - User Libraries?
  - Application Code?
  - Kernel functions?

# Provider Documentation

- Additional documentation may be found here,

Target	Provider	Additional Docs
syscalls	<b>syscall</b>	man(2)
libraries	<b>pid:lib*</b>	man(3C)
app code	<b>pid:a.out</b>	source code, ISV, developers
raw kernel	<b>fbt</b>	Solaris Internals 2 <sup>nd</sup> Ed, <a href="http://cvs.opensolaris.org">http://cvs.opensolaris.org</a>

# Actions

- When a probe fires, an action executes.
- Actions are written in the D programming language.
- Actions can,
  - print output
  - save data to variables, and perform calculations
  - walk kernel or process memory
- With destruction actions allowed, actions can,
  - raise signals on processes
  - execute shell commands
  - write to some areas of memory



# trace() Example

The trace() action accepts one argument and prints it when the probe fired.

```
# dtrace -n 'syscall::exece:return { trace(execname); }'  
dtrace: description 'syscall::exece:return ' matched 1 probe  


| CPU   | ID    | FUNCTION : NAME    |
|-------|-------|--------------------|
| 0     | 76044 | exece:return man   |
| 0     | 76044 | exece:return sh    |
| 0     | 76044 | exece:return neqn  |
| 0     | 76044 | exece:return tbl   |
| 0     | 76044 | exece:return nroff |
| 0     | 76044 | exece:return col   |
| [...] |       |                    |


```



# printf() Example

DTrace ships with a powerful printf(), to print formatted output.

```
# dtrace -n 'syscall::exece:return { printf("%6d %s\n", pid, execname); }'
dtrace: description 'syscall::exece:return ' matched 1 probe
CPU      ID          FUNCTION:NAME
  0    74415          exece:return    4301 sh
  0    74415          exece:return    4304 neqn
  0    74415          exece:return    4305 nroff
  0    74415          exece:return    4306 sh
  0    74415          exece:return    4308 sh
[...]
```

# DTrace Built-In Variables

- Numerous predefined variables can be used, e.g.,
  - pid, tid      Process ID, Thread ID
  - timestamp    Nanosecond timestamp since boot
  - probefunc     Probe function name (3<sup>rd</sup> field)
  - execname      Process name
  - arg0, ...     Function arguments and return value
  - errno          Last syscall failure error code
  - curpsinfo     Struct containing current process info, e.g.,  
                      curpsinfo->pr\_psargs – process + args
- Pointers and structs! DTrace can walk memory using C syntax, and has kernel types predefined.



# User-Defined Variable Types

- DTrace supports the following variable types
  - Integers
  - Structs
  - Pointers
  - Strings
  - Associative arrays
  - Aggregates
- Including types from /usr/include/sys
  - e.g. uint32\_t.



# Predicates

- DTrace predicates are used to filter probes, so that the action fires when a conditional is true.

```
probename /predicate/ { action }
```

- E.g., syscalls for processes called “bash”,

```
# dtrace -n 'syscall:::entry /execname == "bash"/ { @num[probefunc] =  
count(); }'
```

```
dtrace: description 'syscall:::entry ' matched 233 probes
```

```
^C
```

exec	2
[...]	
read	29
write	31
lwp_sigmask	42
sigaction	62



# DTrace – command line

```
usenix> dtrace -n 'syscall:::entry { @scalls[probefunc] = count() }'  
dtrace: description 'syscall:::entry' matched 228 probes  
^C
```

lwp_self	1
fork1	1
fdsync	1
sigpending	1
rexit	1
fxstat	1
...	
write	205
writev	234
brk	272
munmap	357
mmap	394
read	652
pollsys	834
ioctl	
usenix>	1116

## The D language

- D is a C-like language specific to DTrace, with some constructs similar to awk(1)
- Complete access to kernel C types
- Complete access to statics and globals
- Complete support for ANSI-C operators
- Support for strings as first-class citizen
- We'll introduce D features as we need them...

```
#!/usr/sbin/dtrace -s
```

```
probe descriptions
/ predicate /
{
    action statements
}
```



# DTrace – D scripts

```
usenix> cat syscalls_pid.d

#!/usr/sbin/dtrace -s

dtrace:::BEGIN
{
    vtotal = 0;
}

syscall:::entry
/pid == $target/
{
    self->vtime = vtimestamp;
}

syscall:::return
/self->vtime/
{
    @vtime[probefunc] = sum(vtimestamp - self->vtime);
    vtotal += (vtimestamp - self->vtime);
    self->vtime = 0;
}

dtrace:::END
{
    normalize(@vtime, vtotal / 100);
    printa(@vtime);
}
```

a complete dtrace script block,  
including probename, a predicate,  
and an action in the probe clause,  
which sets a thread-local variable



# DTrace – Running syscalls\_pid.d

```
usenix> ./syscalls_pid.d -c date
dtrace: script './sc.d' matched 458 probes
Sun Feb 20 17:01:28 PST 2005
dtrace: pid 2471 has exited
CPU      ID                      FUNCTION:NAME
          0                      :END
  0        2
  getpid                         0
  gettimeofday                   0
  sysi86                          1
  close                           1
  getrlimit                       2
  setcontext                       2
  fstat64                         4
  brk                            8
  open                           8
  read                           9
  munmap                         9
  mmap                           11
  write                          15
  ioctl                          24
```

## DTrace time-based probes

- profile – interval time based probes
  - profile-97hz – profile fires on all CPUs
  - tick-97hz – tick fires on 1 CPU
  - Interval can be specified with various suffixes
    - ns, us, ms, s, min (m), hour (h), day (d), hz
  - arg0 – kernel PC
  - arg1 – user PC
- Use arg0 or arg1 in a predicate for user or kernel profile

```
#dtrace -n 'profile-97hz / arg0 != 0 / { action } /* Am I in the kernel? */  
#dtrace -n 'profile-97hz / arg1 != 0 / { action } /* Am I in user mode? */
```



# Using Providers

Using the syscall provider to track bytes passed to write(2)

```
# dtrace -n 'syscall::write:entry { trace(arg2) }'  
dtrace: description 'write:entry' matched 2 probes  


| CPU | ID   | FUNCTION:NAME |     |
|-----|------|---------------|-----|
| 0   | 1026 | write:entry   | 1   |
| 1   | 1026 | write:entry   | 53  |
| 1   | 9290 | write:entry   | 2   |
| 1   | 1026 | write:entry   | 25  |
| 1   | 9290 | write:entry   | 17  |
| 1   | 1026 | write:entry   | 2   |
| 1   | 9290 | write:entry   | 2   |
| 1   | 1026 | write:entry   | 450 |
| 1   | 9290 | write:entry   | 450 |


```

Using the fbt provider to instrument the kernel ufs\_write() function, and track the filename in the probe action

```
# dtrace -n 'fbt:ufs:ufs_write:entry { printf("%s\n",stringof(args[0]->v_path)); }'  
dtrace: description 'ufs_write:entry' matched 1 probe  


| CPU | ID    | FUNCTION:NAME                                  |  |
|-----|-------|------------------------------------------------|--|
| 13  | 16779 | ufs_write:entry /etc/svc/repository.db-journal |  |
| 13  | 16779 | ufs_write:entry /etc/svc/repository.db-journal |  |
| 13  | 16779 | ufs_write:entry /etc/svc/repository.db-journal |  |
| 13  | 16779 | ufs_write:entry /etc/svc/repository.db-journal |  |
| 13  | 16779 | ufs_write:entry /etc/svc/repository.db-journal |  |
| 13  | 16779 | ufs_write:entry /etc/svc/repository.db-journal |  |
| 13  | 16779 | ufs_write:entry /etc/svc/repository.db-journal |  |
| 13  | 16779 | ufs_write:entry /etc/svc/repository.db-journal |  |
| 13  | 16779 | ufs_write:entry /etc/svc/repository.db-journal |  |


```



# DTrace PID Provider

```
# dtrace -n 'pid221:libc::entry'
dtrace: description 'pid221:libc::entry' matched 2474 probes
CPU      ID          FUNCTION:NAME
 0    41705        set_parking_flag:entry
 0    41762        setup_schedctl:entry
 0    42128        __schedctl:entry
 0    41752        queue_lock:entry
 0    41749        spin_lock_set:entry
 0    41765        no_preempt:entry
 0    41753        queue_unlock:entry
 0    41750        spin_lock_clear:entry
 0    41766        preempt:entry
 0    41791        mutex_held:entry
 0    42160        gettimeofdayofday:entry
 0    41807        _cond_timedwait:entry
 0    41508        abstime_to_relttime:entry
 0    42145        __clock_gettime:entry
 0    41803        cond_wait_common:entry
 0    41800        cond_wait_queue:entry
 0    41799        cond_sleep_queue:entry
 0    41949        _save_nv_regs:entry
 0    41752        queue_lock:entry
 0    41749        spin_lock_set:entry
 0    41765        no_preempt:entry
```

Using the PID provider to instrument  
all the function entry points the process  
calls in libc



# Aggregations

- When trying to understand suboptimal performance, one often looks for patterns that point to bottlenecks
- When looking for patterns, one often doesn't want to study each datum – one wishes to aggregate the data and look for larger trends
- Traditionally, one has had to use conventional tools (e.g. awk(1), perl(1)) to post-process reams of data
- DTrace supports aggregation of data as a first class operation

## Aggregations, cont.

- An aggregation is the result of an aggregating function keyed by an arbitrary tuple
- For example, to count all system calls on a system by system call name:

```
dtrace -n 'syscall:::entry \
{ @syscalls[probefunc] = count(); }'
```

- By default, aggregation results are printed when dtrace(1M) exits



## Aggregations, cont.

- Aggregations need not be named
- Aggregations can be keyed by more than one expression
- For example, to count all ioctl system calls by both executable name and file descriptor:

```
dtrace -n 'syscall::ioctl:entry \
{ @[execname, arg0] = count(); }'
```

# Aggregations, cont.

- Functions:
  - avg() - the average of specified expressions
  - min() - the minimum of specified expressions
  - max() - the maximum of specified expressions
  - count() - number of times the probe fired
  - sum() - running sum
  - quantize() - power-of-two exponential distribution
  - lquantize() - linear frequency distribution
- For example, distribution of write(2) sizes by executable name:

```
dtrace -n 'syscall::write:entry \
{ @[execname] = quantize(arg2); } '
```



# count() aggregation

- Frequency counting syscalls,

```
# dtrace -n 'syscall:::entry { @num[probefunc] = count(); }'
dtrace: description 'syscall:::entry' matched 233 probes
^C
[...]
      writev                                170
      write                                 257
      read                                  896
      pollsys                             959
      ioctl                               1253
```



# Quantize

- Very cool function, here we quantize writech sizes:
  - Here we see that ls processes usually write between 32 and 127 bytes.
    - Makes sense?

```
# dtrace -n 'sysinfo:::writech { @dist[execname] = quantize(arg0); }'
dtrace: description 'sysinfo:::writech' matched 4 probes
^C
[...]
ls
      value  ----- Distribution ----- count
        4 |
        8 |
       16 |
      32 |ooooooooooooooooooooo
        64 |ooooooooooooooooooooo
      128 |
```

[...]



# ls -l

```
# ls -l /etc
dttotal 793
lrwxrwxrwx  1 root      root          12 Mar 21 03:28 TIMEZONE -> default/init
drwxr-xr-x  4 root      sys           6 Apr 16 06:59 X11
drwxr-xr-x  2 adm       adm           3 Mar 20 09:25 acct
drwxr-xr-x  3 root      root          3 Apr 16 23:11 ak
lrwxrwxrwx  1 root      root          12 Mar 21 03:28 aliases -> mail/aliases
drwxr-xr-x  5 root      sys           5 Feb 20 23:29 amd64
drwxr-xr-x  7 root      bin           18 Mar 20 09:20 apache
drwxr-xr-x  4 root      bin           7 Feb 20 23:12 apache2
drwxr-xr-x  2 root      sys           5 Feb 20 23:27 apoc
-rw-r--r--  1 root      bin          1012 Mar 20 09:33 auto_home
-rw-r--r--  1 root      bin          1066 Mar 20 09:33 auto_master
lrwxrwxrwx  1 root      root          16 Mar 21 03:28 autopush -> ../sbin/autopu
[...]
```

**ls writes one line at a time, each around 80 chars long.**

# Quantize aggregations – write(2) bytes by process

# DTrace Enhancements post S10 FCS

- Multiple aggregations with `printa()`
- Aggregation key sort options
- `(u)func(%pc)`, `(u)mod(%pc)`, `(u)sym(%pc)` dtrace functions
  - Get symbolic name from address
- `ucaller` function
  - Track function callers
- String parsing routines
- `fds[]`
  - array of `fileinfo_t`'s indexed by fd
- Providers
  - `fsinfo`
  - `sysevent`
  - `Xserver`
  - `iscsi`

# Multiple aggregation printa()

Release: 08/07-30

- multiple aggregations in single printa()
- aggregations must have same type signature
- output is effectively joined by key
- 0 printed when no value present for a key
- default behavior is to sort by first aggregation value  
(ties broken by key order)



# Multiple aggregation printa()

```
/* multagg.d */
syscall::write::entry
{
    @wbytes[execname, pid] = sum(arg2);
}

syscall::read::entry
{
    @rbytes[execname, pid] = sum(arg2);
}

END
{
    normalize(@rbytes, 1024);
    normalize(@wbytes, 1024);

    printf("%20s %10s %10s %10s\n", "PROGRAM", "PID",
           "READS", "WRITES");
    printa("%20s %10d %10@d %10@d\n", @rbytes, @wbytes);
}
```

# Multiple aggregation printa()

```
# dtrace -q -s ./multagg.d  
^C
```

PROGRAM	PID	READS	WRITES
dtrace	101605	0	0
nautilus	101606	0	0
battstat-applet-	100854	0	15
gnome-settings-d	100781	0	0
gnome-session	100728	0	0
dsdm	100712	6	0
gnome-terminal	101342	8	15
xscreensaver	100946	9	0
soffice.bin	101325	10	48
init	1	32	0
gnome-panel	100792	39	4
gconfd-2	100767	47	5
nautilus	100794	265	170
wnck-applet	100822	714	262
metacity	100789	726	263
gedit	101607	3986	64
Xorg	100535	12374	0

# Aggregation sorting options

Release: 08/07-30

- aggregations sorted by value by default
- options allow change of behaviour
  - aggsortkey - sort by key order, ties broken by value
  - aggsortrev - reverse sort
  - aggsortpos - position of the aggregation to use as sort primary sort key with multiple aggs
  - aggsortkeypos - position of key to use as primary sort key when with multiple aggs
- Use the above in combination

# Aggregation sorting options

```
/* aggsort.d */
syscall::read:entry
{
    @avg[execname, pid] = avg(arg2);
    @max[execname, pid] = max(arg2);
    @min[execname, pid] = min(arg2);
    @cnt[execname, pid] = count();
}
END
{
    printf("%20s %10s %10s %10s %10s %10s\n", "EXECNAME", "PID",
           "COUNT", "MIN", "MAX", "AVG");
    printa("%20s %10d %@10d %@10d %@10d %@10d\n", @cnt, @min,
          @max, @avg);
}
```

# Aggregation sorting options

```
# dtrace -q -s ./aggsort.d
```

```
^C
```

EXECNAME	PID	COUNT	MIN	MAX	AVG
battstat-applet-	100981	2	32	32	32
gnome-settings-d	100853	3	32	64	53
soffice.bin	101382	10	32	32	32
dsdm	100708	10	32	160	54
xscreensaver	101082	14	32	32	32
gnome-panel	100896	24	12	168	51
firefox-bin	101363	31	1	1024	35
gnome-terminal	101029	40	32	4096	163
nautilus	100906	119	32	480	48
wnck-applet	100961	161	8	128	32
metacity	100893	447	4	96	33
Xorg	100534	926	64	5104	3263

Sort by value of first aggregation (default)

# aggregation sorting options

```
# dtrace -q -s ./aggsort.d -x aggsortrev
```

```
^C
```

EXECNAME	PID	COUNT	MIN	MAX	AVG
Xorg	100534	875	64	5104	3433
metacity	100893	434	4	96	33
wnck-applet	100961	145	8	64	32
xscreensaver	101082	71	4	32	28
gnome-terminal	101029	54	32	4096	125
nautilus	100906	51	32	608	74
firefox-bin	101363	23	1	1	1
dsdm	100708	19	32	32	32
gnome-panel	100896	18	12	168	51
gnome-settings-d	100853	7	32	32	32
soffice.bin	101382	6	32	32	32

Reverse sort using value of first aggregation

# Aggregation sorting options

```
# dtrace -q -s ./aggsort.d -x aggsortkey -x aggsortkeypos=1 -x aggsortrev  
^C
```

EXECNAME	PID	COUNT	MIN	MAX	AVG
soffice.bin	101382	525	4	1440	33
firefox-bin	101363	29	1	1024	36
thunderbird-bin	101337	2	1	1024	512
xscreensaver	101082	11	32	64	34
gnome-terminal	101029	27	32	4096	220
wnck-applet	100961	161	8	96	32
nautilus	100906	79	32	320	40
gnome-panel	100896	26	12	168	49
metacity	100893	196	4	128	34
gnome-settings-d	100853	4	32	64	40
dsdm	100708	23	32	128	40
xorg	100534	885	64	4940	3688

Reverse sort by key in second position

# Aggregation sorting options

```
# dtrace -q -s ./aggsort.d -x aggsortpos=2 -x aggsortrev
```

```
^C
```

EXECNAME	PID	COUNT	MIN	MAX	AVG
xorg	100534	1655	64	5104	3756
gnome-terminal	101029	137	32	4096	69
wnck-applet	100961	453	4	1152	34
xscreensaver	101082	23	32	1024	115
nautilus	100906	43	8	736	69
soffice.bin	101382	637	4	288	30
gnome-panel	100896	122	8	168	39
metacity	100893	421	4	128	34
notification-are	100983	2	64	64	64
mixer_applet2	100985	2	64	64	64
gnome-settings-d	100853	7	32	32	32
dssdm	100708	103	4	32	31
thunderbird-bin	101337	7	4	32	24
firefox-bin	101363	39	1	32	5

Reverse sorted by value of third aggregation

# (u)mod/(u)func/(u)sym

Release: 08/07-23

- Profiling often requires post-processing when using %a/%A to print arg0/arg1 symbolically
- Samples in format [module]'[func]+[offset]
- Want to first get high level view and then drill down
- (u)mod(%pc) - module name
- (u)func(%pc) - function name
- (u)sym(%pc) - symbol name



# (u)mod/(u)func/(u)sym

```
#pragma D option aggsortkey

cpc:::dtlbtm-u-5000
{
    @[execname, umod(arg1)] = count();
}

END
{
    printf("%20s %30s %10s\n", "EXECUTABLE", "MODULE",
           "COUNT");
    printa("%20s %30A %10@d\n", @);
}
```

**Example uses prototype cpc provider to show  
TLB misses on a global basis broken down by  
module.**



# (u)mod/(u)func/(u)sym

```
# ./tlbmissbymod.d  
^C
```

EXECUTABLE	MODULE	COUNT
Xorg	Xorg	17
Xorg	libramdac.so	1
Xorg	libfb.so	1
Xorg	radeon_drv.so	1
Xorg	libc.so.1	3
battstat-applet-	libglib-2.0.so.0.1200.4	1
battstat-applet-	libgobject-2.0.so.0.1200.4	1
gconfd-2	libORBit-2.so.0.1.0	1
gconfd-2	libgconf-2.so.4.1.0	1
metacity	libpangocairo-1.0.so.0.1400.7	1
metacity	libgdk_pixbuf-2.0.so.0.1000.6	1
metacity	libgobject-2.0.so.0.1200.4	1
metacity	libglib-2.0.so.0.1200.4	2
metacity	libgdk-x11-2.0.so.0.1000.6	1
metacity	libc.so.1	2



# ucaller variable

## Release: 08/07-23

```
# dtrace -n 'pid$target::malloc:entry{@[ufunc(ucaller)] = count();}' -p 101384
dtrace: description 'pid$target::malloc:entry' matched 2 probes
^C

sax.uno.so`0xf97b2fe7                                2
sax.uno.so`0xf97b30e7                                2
sax.uno.so`0xf97b3104                                2
sax.uno.so`0xf97b44ba                                2
sax.uno.so`0xf97b44cf                                2
sax.uno.so`0xf97b8a18                                4
libX11.so.4`_XAllocScratch                          5
libX11.so.4`miRegionOp                            6
libX11.so.4`XQueryTree                           8
libX11.so.4`XGetWindowProperty                   13
libX11.so.4`XSetClassHint                         13
libX11.so.4`XGetImage                            179
libuno_sal.so.3`osl_createMutex                  322
libX11.so.4`XCreateRegion                        844
libX11.so.4`XCreateGC                            959
libc.so.1`calloc                                 1074
libglib-2.0.so.0.1200.4`standard_malloc          3533
libCrun.so.1`__lc2n6FI_pv_                      28668
```



# The fds[] variable

## Release: 01/06-16

- array of fileinfo\_t's indexed by integer (fd)
- inlines expanded to accommodate associative arrays for this.
- Definition in /usr/lib/dtrace/io.d
- fileinfo\_t gets new fi\_oflags member



# The fds[] variable

```
#pragma D option quiet
```

```
syscall::write:entry
/fds[arg0].fi_oflags & O_APPEND/
{
    printf("%s appending file %s at offset %d\n",
           execname, fds[arg0].fi.pathname, fds[0].fi_offset);
}
```

```
# ./fds.d
ksh appending file /.sh_history at offset 349345
ksh appending file /.sh_history at offset 349378
```



# Allowing dtrace for non-root users

- Setting dtrace privileges
- Add a line for the user in /etc/user\_attr

```
mauroj::::defaultpriv=basic,dtrace_user,dtrace_kernel,dtrace_proc,  
proc_owner
```

from privileges(5)

**PRIV\_DTRACE\_KERNEL**

Allow DTrace kernel-level tracing.

**PRIV\_DTRACE\_PROC**

Allow DTrace process-level tracing. Allow process-level tracing probes to be placed and enabled in processes to which the user has permissions.

**PRIV\_DTRACE\_USER**

Allow DTrace user-level tracing. Allow use of the syscall and profile DTrace providers to examine processes to which the user has permissions.

# Modular Debugger - mdb(1)

- Solaris 8 mdb(1) replaces adb(1) and crash(1M)
- Allows for examining a live, running system, as well as post-mortem (dump) analysis
- Solaris 9 mdb(1) adds...
  - Extensive support for debugging of processes
  - /etc/crash and adb removed
  - Symbol information via compressed typed data
  - Documentation
- MDB Developers Guide
  - mdb implements a rich API set for writing custom dcmds
  - Provides a framework for kernel code developers to integrate with mdb(1)

# Modular Debugger - mdb(1)

- mdb(1) basics
  - 'd' commands (dcmd)
    - ::dcmds -l for a list
    - expression::dcmd
    - e.g. 0x300acde123::ps
  - walkers
    - ::walkers for a list
    - expression::walk <walker\_name>
    - e.g. ::walk cpu
  - macros
    - !ls /usr/lib/adb for a list
    - expression\$<macro
    - e.g. cpu0\$<cpu

# Modular Debugger – mdb(1)

- Symbols and typed data
  - address::print (for symbol)
  - address::print <type>
  - e.g. cpu0::print cpu\_t
  - cpu\_t::sizeof
- Pipelines
  - expression, dcmd or walk can be piped
  - ::walk <walk\_name> | ::dcmd
  - e.g. ::walk cpu | ::print cpu\_t
  - Link Lists
  - address::list <type> <member>
  - e.g. 0x70002400000::list page\_t p\_vpnext
- Modules
  - Modules in /usr/lib/mdb, /usr/platform/lib/mdb etc
  - mdb can use adb macros
  - Developer Interface - write your own dcmands and walkers



```
> ::cpuinfo
ID ADDR      FLG NRUN BSPL PRI RNRN KRNrn SWITCH THREAD      PROC
 0 0000180c000 1b    0    0  37   no    no t-0  30002ec8ca0 threads
 1 30001b78000 1b    0    0  27   no    no t-0  31122698960 threads
 4 30001b7a000 1b    0    0  59   no    no t-0  30ab913cd00 find
 5 30001c18000 1b    0    0  59   no    no t-0  31132397620 sshd
 8 30001c16000 1b    0    0  37   no    no t-0  3112280f020 threads
 9 30001c0e000 1b    0    0  59   no    no t-1  311227632e0 mdb
12 30001c06000 1b    0    0  -1   no    no t-0  2a100609cc0 (idle)
13 30001c02000 1b    0    0  27   no    no t-1  300132c5900 threads
> 30001b78000::cpuinfo -v
ID ADDR      FLG NRUN BSPL PRI RNRN KRNrn SWITCH THREAD      PROC
 1 30001b78000 1b    0    0  -1   no    no t-3  2a100307cc0 (idle)
|
RUNNING <-+
READY
EXISTS
ENABLE

> 30001b78000::cpuinfo -v
ID ADDR      FLG NRUN BSPL PRI RNRN KRNrn SWITCH THREAD      PROC
 1 30001b78000 1b    0    0  27   no    no t-1  300132c5900 threads
|
RUNNING <-+
READY
EXISTS
ENABLE

> 300132c5900::findstack
stack pointer for thread 300132c5900: 2a1016dd1a1
000002a1016dd2f1 user_rtt+0x20()
```

# **mdb(1) & dtrace(1) – Perfect Together**

```
# mdb -k

Loading modules: [ unix krtld genunix specfs dtrace ufs sd ip sctp usba fcp fctl nca nfs random sppp
lofs crypto ptm logindmux md isp cpc fcip ipc ]

> ufs_read::nm -f ctype
C Type
int (*)(struct vnode *, struct uio *, int, struct cred *, struct caller_context *)
> ::print -t struct vnode
{
    kmutex_t v_lock {
        void * [1] _opaque
    }
    uint_t v_flag
    uint_t v_count
    void *v_data
    struct vfs *v_vfsp
    struct stdata *v_stream
    enum vtype v_type
    dev_t v_rdev
    struct vfs *v_vfsmountedhere
    struct vnodeops *v_op
    struct page *v_pages
    pgcnt_t v_npages
    ...
    char *v_path
}

# dtrace -n 'ufs_read:entry { printf("%s\n",stringof(args[0]->v_path))}'
dtrace: description 'ufs_read:entry' matched 1 probe
CPU      ID          FUNCTION:NAME
1  16777          ufs_read:entry /usr/bin/cut
1  16777          ufs_read:entry /usr/bin/cut
1  16777          ufs_read:entry /usr/bin/cut
1  16777          ufs_read:entry /usr/bin/cut
1  16777          ufs_read:entry /lib/ld.so.1
1  16777          ufs_read:entry /lib/ld.so.1
```

```
# dtrace -n 'ufs_read:entry { printf("%s\n",stringof(args[0]->v_path));}'  
dtrace: description 'ufs_read:entry' matched 1 probe  


| CPU | ID    | FUNCTION:NAME               |
|-----|-------|-----------------------------|
| 1   | 16777 | ufs_read:entry /usr/bin/cut |
| 1   | 16777 | ufs_read:entry /lib/ld.so.1 |
| 1   | 16777 | ufs_read:entry /lib/ld.so.1 |


```



# Kernel Statistics

- Solaris uses a central mechanism for kernel statistics
  - "kstat"
  - Kernel providers
    - raw statistics (c structure)
    - typed data
    - classed statistics
  - Perl and C API
  - kstat(1M) command

```
# kstat -n system_misc
module: unix
name: system_misc
instance: 0
class: misc

avenrun_15min          90
avenrun_1min           86
avenrun_5min           87
boot_time              1020713737
clk_intr               2999968
crtime                 64.1117776
deficit                0
lbolt                  2999968
ncpus                  2
```



# kstat

```
zeeroh> kstat -n e1000g0
module: e1000g
name: e1000g0
    brdcstrcv
    brdcstxmt
    collisions
    crttime
    ierrors
    ifspeed
    ipackets
    ipackets64
    multircv
    multixmt
    norcvbuf
    noxmtbuf
    obytes
    obytes64
    oerrors
    opackets
    opackets64
    rbytes
    rbytes64
    snaptime
    unknowns
                                instance: 0
                                class:     net
                                31397806
                                104402
                                0
                                142.615878752
                                0
                                1000000000
                                33545701
                                33545701
                                796061
                                0
                                0
                                0
                                169913542
                                169913542
                                0
                                1141181
                                1141181
                                1559052386
                                5854019682
                                15146011.0147421
                                819041
```



# kstat

```
zeeroh> kstat e1000g:0:e1000g0:opackets 1 10
module: e1000g                                instance: 0
name:   e1000g0                               class:    net
       opackets                           1141265

module: e1000g                                instance: 0
name:   e1000g0                               class:    net
       opackets                           1141268

module: e1000g                                instance: 0
name:   e1000g0                               class:    net
       opackets                           1141273

module: e1000g                                instance: 0
name:   e1000g0                               class:    net
       opackets                           1141279
```

# Procfs Tools

- Observability (and control) for active processes through a pseudo file system (/proc)
- Extract interesting bits of information on running processes
- Some commands work on core files as well

**pargs**

**pfiles**

**pflags**

**pstop**

**pcred**

**prun**

**pldd**

**pwait**

**psig**

**ptree**

**pstack**

**ptime**

**pmap**

**preap**



# pflags, pcred, pldd, pkill

```
nv70b> pflags 3645
```

```
3645: ./loader 2 0 /zp/space
      data model = _ILP32 flags = MSACCT|MSFORK
/1:   flags = ASLEEP lwp_wait(0x2,0x80471f0)
/2:   flags = 0
/3:   flags = 0
```

```
nv70b> pcred 3645
```

```
3645: e/r/suid=20821 e/r/sgid=3
```

```
nv70b> pldd 3645
```

```
3645: ./loader 2 0 /zp/space
/lib/libc.so.1
```

```
nv70b> pkill loader
```

```
nv70b>
```

```
[1]+ Terminated ./loader 2 0 /zp/space
```



# psig

```
nv70b> psig 3654
3654: ./loader 2 0 /zp/space
HUP      default
INT      default
QUIT     default
ILL      default
TRAP     default
ABRT    default
EMT      default
FPE      default
KILL     default
BUS      default
SEGV    default
SYS      default
PIPE     default
ALRM    default
TERM    default
USR1    caught  timer          0
USR2    default
CLD      default            NOCLDSTOP
PWR      default
```

... .



# pstack

```
nv70b> pstack 3602
3602: /usr/lib/firefox/firefox-bin -UILocale C -contentLocale C
----- lwp# 1 / thread# 1 -----
d117ec45 pollsys (82d5910, 8, 0, 0)
d1134212 poll (82d5910, 8, ffffffff) + 52
d0c06653 g_main_context_iterate (80c3260, 1, 1, 815f1f0) + 397
d0c06c8c g_main_loop_run (82f45f8) + 1b8
d0964fae gtk_main (8047028, 808578c, 8189718, 8046d88, cd694803, 818ed20) + b2
cdbc4bb4 __1cKnsAppShellDRun6M_I_ (818ed20) + 34
cd694803 __1cMnsAppStartupDRun6M_I_ (8189718) + 2b
08061824 XRE_main (5, 8047098, 8085760) + 25f4
0805a61d main (5, 8047098, 80470b0) + 25
0805a56a _start (5, 8047210, 804722d, 8047237, 8047239, 8047248) + 7a
----- lwp# 2 / thread# 2 -----
d117ec45 pollsys (cdeabc70, 1, 0, 0)
d1134212 poll (cdeabc70, 1, ffffffff) + 52
d0f48bfa __pr_poll_with_poll (80f00a8, 1, ffffffff) + 39a
d0f48dc6 PR_Poll (80f00a8, 1, ffffffff) + 16
ce15657a __1cYnsSocketTransportServiceEPoll6MpI_i_ (80efbc0, cdeabf74) + 11e
ce157118 __1cYnsSocketTransportServiceDRun6M_I_ (80efbc0) + 68c
d103fff4 __1cInsThreadEMain6Fpv_v_ (80f0298) + 74
d0f4ab0d __pt_root (80f2720) + d1
d117d952 __thr_setup (cdd90200) + 52
d117dbb0 __lwp_start (cdd90200, 0, 0, 0, 0, 0, 0)
----- lwp# 3 / thread# 3 -----
d117dc09 __lwp_park (0, cd58de58, 0)
. . .
```



# pfiles

```
nv70b> pfiles 3602
3602: /usr/lib/firefox/firefox-bin -UILocale C -contentLocale C
      Current rlimit: 512 file descriptors
      0: S_IFCHR mode:0666 dev:279,0 ino:6815752 uid:0 gid:3 rdev:13,2
          O_RDONLY|O_LARGEFILE
          /devices/pseudo/mm@0:null
      1: S_IFCHR mode:0666 dev:279,0 ino:6815752 uid:0 gid:3 rdev:13,2
          O_WRONLY|O_APPEND|O_CREAT|O_LARGEFILE
          /devices/pseudo/mm@0:null
      2: S_IFCHR mode:0666 dev:279,0 ino:6815752 uid:0 gid:3 rdev:13,2
          O_WRONLY|O_APPEND|O_CREAT|O_LARGEFILE
          /devices/pseudo/mm@0:null
      3: S_IFDOOR mode:0444 dev:290,0 ino:43 uid:0 gid:0 size:0
          O_RDONLY|O_LARGEFILE FD_CLOEXEC door to nscd[3545]
          /var/run/name_service_door
      4: S_IFIFO mode:0666 dev:290,0 ino:40 uid:0 gid:0 size:0
          O_RDWR|O_NONBLOCK FD_CLOEXEC
      . . .
      19: S_IFSOCK mode:0666 dev:287,0 ino:28594 uid:0 gid:0 size:0
          O_RDWR|O_NONBLOCK FD_CLOEXEC
          SOCK_STREAM
          SO_REUSEADDR,SO_SNDBUF(49152),SO_RCVBUF(49152)
          sockname: AF_INET 127.0.0.1 port: 59686
      . . .
```



# pfiles

```
solaris10> pfiles 26337
26337: /usr/lib/ssh/sshd
  Current rlimit: 256 file descriptors
  0: S_IFCHR mode:0666 dev:270,0 ino:6815752 uid:0 gid:3 rdev:13,2
    O_RDWR|O_LARGEFILE
    /devices/pseudo/mm@0:null
  1: S_IFCHR mode:0666 dev:270,0 ino:6815752 uid:0 gid:3 rdev:13,2
    O_RDWR|O_LARGEFILE
    /devices/pseudo/mm@0:null
  2: S_IFCHR mode:0666 dev:270,0 ino:6815752 uid:0 gid:3 rdev:13,2
    O_RDWR|O_LARGEFILE
    /devices/pseudo/mm@0:null
  3: S_IFDOOR mode:0444 dev:279,0 ino:59 uid:0 gid:0 size:0
    O_RDONLY|O_LARGEFILE FD_CLOEXEC door to nscd[93]
    /var/run/name_service_door
  4: S_IFSOCK mode:0666 dev:276,0 ino:36024 uid:0 gid:0 size:0
    O_RDWR|O_NONBLOCK
      SOCK_STREAM
      SO_REUSEADDR,SO_KEEPALIVE,SO_SNDBUF(49152),SO_RCVBUF(49880)
      sockaddr: AF_INET6 ::ffff:129.154.54.9 port: 22
      peername: AF_INET6 ::ffff:129.150.32.45 port: 52002
  5: S_IFDOOR mode:0644 dev:279,0 ino:55 uid:0 gid:0 size:0
    O_RDONLY FD_CLOEXEC door to keyserv[179]
    /var/run/rpc_door/rpc_100029.1
....
```



## pwdx, pstop, prun

```
nv70b> pwdx 3666
3666: /zp/home/mauroj/Programs
nv70b> pflags 3666
3666: ./loader 2 0 /zp/space
      data model = _ILP32 flags = MSACCT|MSFORK
/1:   flags = ASLEEP lwp_wait(0x2,0x80471f0)
/2:   flags = 0
/3:   flags = 0

nv70b> pstop 3666
nv70b> pflags 3666
3666: ./loader 2 0 /zp/space
      data model = _ILP32 flags = MSACCT|MSFORK
      sigpend = 0x00008000,0x00000000
/1:   flags = STOPPED|ISTOP|ASLEEP lwp_wait(0x2,0x80471f0)
      why = PR_REQUESTED
/2:   flags = STOPPED|ISTOP
      why = PR_REQUESTED
/3:   flags = STOPPED|ISTOP
      why = PR_REQUESTED

nv70b> prun 3666
nv70b> pflags 3666
3666: ./loader 2 0 /zp/space
      data model = _ILP32 flags = MSACCT|MSFORK
/1:   flags = ASLEEP lwp_wait(0x2,0x80471f0)
/2:   flags = 0
/3:   flags = 0
```



# prstat(1)

- top-like utility to monitor running processes
- Sort on various thresholds (cpu time, RSS, etc)
- Enable system-wide microstate accounting
  - Monitor time spent in each microstate
- Solaris 9 - “projects” and “tasks” aware

PID	USERNAME	SIZE	RSS	STATE	PRI	NICE	TIME	CPU	PROCESS/NLWP
2597	ks130310	4280K	2304K	cpu1	0	0	0:01:25	22%	imapd/1
29195	bc21502	4808K	4160K	sleep	59	0	0:05:26	1.9%	imapd/1
3469	tjobson	6304K	5688K	sleep	53	0	0:00:03	1.0%	imapd/1
3988	tja	8480K	7864K	sleep	59	0	0:01:53	0.5%	imapd/1
5173	root	2624K	2200K	sleep	59	0	11:07:17	0.4%	nfsd/18
2528	root	5328K	3240K	sleep	59	0	19:06:20	0.4%	automountd/2
175	root	4152K	3608K	sleep	59	0	5:38:27	0.2%	ypserv/1
4795	snoqueen	5288K	4664K	sleep	59	0	0:00:19	0.2%	imapd/1
3580	mauroj	4888K	4624K	cpu3	49	0	0:00:00	0.2%	prstat/1
1365	bf117072	3448K	2784K	sleep	59	0	0:00:01	0.1%	imapd/1
8002	root	23M	23M	sleep	59	0	2:07:21	0.1%	esd/1
3598	wabbott	3512K	2840K	sleep	59	0	0:00:00	0.1%	imapd/1
25937	pdanner	4872K	4232K	sleep	59	0	0:00:03	0.1%	imapd/1
11130	smalm	5336K	4720K	sleep	59	0	0:00:08	0.1%	imapd/1



# truss(1)

- “trace” the system calls of a process/command
- Extended to support user-level APIs (-u, -U)
- Can also be used for profile-like functions (-D, -E)
- Is thread-aware as of Solaris 9 (pid/lwp\_id)

```
usenix> truss -c -p 2556
^C
      syscall      seconds    calls  errors
read                  .013     1691
pread                 .015     1691
pread64                .056      846
-----
sys totals:          .085     4228      0
usr time:            .014
elapsed:             7.030
usenix> truss -D -p 2556
/2: 0.0304 pread(11, "02\0\0\001\0\0\0\0\n c\0\0"..., 256, 0) = 256
/2: 0.0008 read(8, "1ED0C2 I", 4) = 4
/2: 0.0005 read(8, "@C9 b @FDD4 EC6", 8) = 8
/2: 0.0006 pread(11, "02\0\0\001\0\0\0\0\n c\0\0"..., 256, 0) = 256
/2: 0.0134 pread64(10, "\0\0\0\0\0\0\0\0\0\0\0\0"..., 8192, 0x18C8A000) = 8192
/2: 0.0006 pread(11, "02\0\0\001\0\0\0\0\n c\0\0"..., 256, 0) = 256
/2: 0.0005 read(8, "D6 vE5 @", 4) = 4
/2: 0.0005 read(8, "E4CA9A -01D7AAA1", 8) = 8
/2: 0.0006 pread(11, "02\0\0\001\0\0\0\0\n c\0\0"..., 256, 0) = 256
```



# lockstat(1M)

- Provides for kernel lock statistics (mutex locks, reader/writer locks)
- Also serves as a kernel profiling tool
- Use “-i 971” for the interval to avoid collisions with the clock interrupt, and gather fine-grained data

```
#lockstat -i 971 sleep 300 > lockstat.out
```

```
#lockstat -i 971 -I sleep 300 > lockstatI.out
```



# Examining Kernel Activity - Kernel Profiling

```
# lockstat -kIi997 sleep 10
Profiling interrupt: 10596 events in 5.314 seconds (1994 events/sec)
Count indv cuml rcnt      nsec CPU+PIL          Caller
-----
 5122  48%  48% 1.00    1419 cpu[0]           default_copyout
1292  12%  61% 1.00    1177 cpu[1]           splx
1288  12%  73% 1.00    1118 cpu[1]           idle
  911   9%  81% 1.00    1169 cpu[1]           disp_getwork
  695   7%  88% 1.00    1170 cpu[1]           i_ddi_splhigh
  440   4%  92% 1.00    1163 cpu[1]+11        splx
  414   4%  96% 1.00    1163 cpu[1]+11        i_ddi_splhigh
  254   2%  98% 1.00    1176 cpu[1]+11        disp_getwork
    27   0%  99% 1.00    1349 cpu[0]           uiomove
    27   0%  99% 1.00    1624 cpu[0]           bzero
    24   0%  99% 1.00    1205 cpu[0]           mmrw
    21   0%  99% 1.00    1870 cpu[0]           (usermode)
     9   0%  99% 1.00    1174 cpu[0]           xcopyout
     8   0%  99% 1.00     650 cpu[0]            kt10
     6   0%  99% 1.00    1220 cpu[0]           mutex_enter
     5   0%  99% 1.00    1236 cpu[0]           default_xcopyout
     3   0% 100% 1.00    1383 cpu[0]           write
     3   0% 100% 1.00    1330 cpu[0]           getminor
     3   0% 100% 1.00     333 cpu[0]            ut10
     2   0% 100% 1.00     961 cpu[0]           mmread
     2   0% 100% 1.00    2000 cpu[0]+10        read_rtc
```



# trapstat(1)

- Solaris 9, Solaris 10 (and beyond...)
- Statistics on CPU traps
  - Very processor architecture specific
- “-t” flag details TLB/TSB miss traps
  - Extremely useful for determining if large pages will help performance
    - Solaris 9 Multiple Page Size Support (MPSS)



#trapstat -t		cpu m  itlb-miss %tim		itsb-miss %tim   dtlb-miss %tim		dtsb-miss %tim   %tim	
0	u	360	0.0	0	0.0	324	0.0
0	k	44	0.0	0	0.0	21517	1.1
1	u	2680	0.1	0	0.0	10538	0.5
1	k	111	0.0	0	0.0	11932	0.7
4	u	3617	0.2	2	0.0	28658	1.3
4	k	96	0.0	0	0.0	14462	0.8
5	u	2157	0.1	7	0.0	16055	0.7
5	k	91	0.0	0	0.0	12987	0.7
8	u	1030	0.1	0	0.0	2102	0.1
8	k	124	0.0	1	0.0	11452	0.6
9	u	7739	0.3	15	0.0	112351	4.9
9	k	78	0.0	3	0.0	65578	3.2
12	u	1398	0.1	5	0.0	8603	0.4
12	k	156	0.0	4	0.0	13471	0.7
13	u	303	0.0	0	0.0	346	0.0
13	k	10	0.0	0	0.0	27234	1.4
ttl		19994	0.1	37	0.0	357610	2.1
						5603	0.2
							2.4

# The \*stat Utilities

- mpstat(1)
  - System-wide view of CPU activity
- vmstat(1)
  - Memory statistics
  - Don't forget “vmstat -p” for per-page type statistics
- netstat(1)
  - Network packet rates
  - Use with care – it does induce probe effect
- iostat(1)
  - Disk I/O statistics
  - Rates (IOPS), bandwidth, service times
- sar(1)
  - The kitchen sink



# cputrack(1)

- Gather CPU hardware counters, per process

```
solaris> cputrack -N 20 -c pic0=DC_access,pic1=DC_miss -p 19849
      time lwp      event      pic0      pic1
    1.007   1      tick  34543793  824363
    1.007   2      tick      0        0
    1.007   3      tick 1001797338 5153245
    1.015   4      tick 976864106 5536858
    1.007   5      tick 1002880440 5217810
    1.017   6      tick 948543113 3731144
    2.007   1      tick 15425817 745468
    2.007   2      tick      0        0
    2.014   3      tick 1002035102 5110169
    2.017   4      tick 976879154 5542155
    2.030   5      tick 1018802136 5283137
    2.033   6      tick 1013933228 4072636
.....
solaris> bc -l
824363/34543793
.02386428728310177171
((100-(824363/34543793)))
99.97613571271689822829
```



# cpustat(1)

```
# cpustat -h
```

Usage:

```
cpustat [-c events] [-p period] [-nstD] [interval [count]]
```

-c events specify processor events to be monitored

-n suppress titles

-p period cycle through event list periodically

-s run user soaker thread for system-only events

-t include %tick register

-D enable debug mode

-h print extended usage information

Use cputrack(1) to monitor per-process statistics.

CPU performance counter interface: SPARC64 VI

event specification syntax:

```
[picn=<eventn>[,attr[n][=<val>]]],[picn=<eventn>[,attr[n][=<val>]]],...]
```

event0: cycle\_counts instruction\_counts op\_stv\_wait

load\_store\_instructions branch\_instructions

floating\_instructions impdep2\_instructions

prefetch\_instructions flush\_rs 2iid\_use toq\_rsbr\_phantom

trap\_int\_vector ts\_by\_sxmiss active\_cycle\_count

op\_stv\_wait\_sxmiss eu\_comp\_wait swpf\_fail\_all

sx\_miss\_wait\_pf jbus\_cpi\_count jbus\_reqbus1\_busy

event1: cycle\_counts instruction\_counts instruction\_flow\_counts

iwr\_empty op\_stv\_wait load\_store\_instructions

branch\_instructions floating\_instructions

....



# cpustat(1)

```
# cpustat -c pic0=cycle_counts,pic1=instruction_counts,sys
      time  cpu event      pic0      pic1
 5.011    8  tick 11411962932 5348741180
 5.011   131  tick 11400819057 5249797028
 5.011   130  tick 11400858896 5244266999
 5.011    16  tick 11407356349 5303050712
 5.011    10  tick 11409702171 5344586657
 5.011    18  tick 11407295550 5306963656
 5.011    19  tick 11406349340 5292477138
 5.011     0  tick 11412859729 5222752733
 5.011    17  tick 11408093975 5307307043
 5.011    26  tick 11403459560 5254359643
 5.011   144  tick 11394770612 5325801245
 5.011    24  tick 11403518595 5256957295
 5.011   128  tick 11397600354 5234695931
 5.012     1  tick 11414392475 5284187266
 5.011   137  tick 11397641918 5313760153
 5.011    11  tick 11410206642 5347201300
 5.011    27  tick 11402593843 5285054790
.
.
.
# bc -l
11394446629/5320324508
2.14168263831774526036
```



USE



IMPROVE



EVANGELIZE

# Applying The Tools - Example



# Start with a System View

```
# mpstat 1
CPU minf mjf xcal  intr ithr  csw icsw migr smtx  srw syscl usr sys wt idl
 0   0   0 294  329 227  117   60   12 40597    0 245787  10  90   0   0
 1   11  0   0 141   4   73   41   12 37736    0 244729  11  89   0   0
 2   0   0   0 140   2   64   37   1 34046    0 243383  10  90   0   0
 3   0   0   0 130   0   49   32   2 31666    0 243440  10  90   0   0
CPU minf mjf xcal  intr ithr  csw icsw migr smtx  srw syscl usr sys wt idl
 0   0   0 16  432 230  149   68   25 42514   25 250163  10  90   0   0
 1   0   0 100 122   5  117   55   26 38418    8 247621  10  90   0   0
 2   0   0 129 103   2  124   53   12 34029   12 244908   9  91   0   0
 3   0   0 24  123   0  110   45    6 30893   18 242016  10  90   0   0
```

- What jumps out at us...
  - Processors a fully utilized, 90% sys
    - Question: Where is the kernel spending time?
  - syscalls-per-second are high
    - Question: What are these system calls, and where are they coming from
  - mutex's per second are high
    - Question: Which mutex locks, and why?



# Processor – kernel profile

```
# lockstat -i997 -Ikw sleep 30
```

```
Profiling interrupt: 119780 events in 30.034 seconds (3988 events/sec)
```

Count	indv	cuml	rcnt	nsec	CPU+PIL	Hottest Caller
29912	25%	25%	0.00	5461	cpu[2]	kcopy
29894	25%	50%	0.00	5470	cpu[1]	kcopy
29876	25%	75%	0.00	5401	cpu[3]	kcopy
29752	25%	100%	0.00	5020	cpu[0]	kcopy
119	0%	100%	0.00	1689	cpu[0]+10	dosoftint
71	0%	100%	0.00	1730	cpu[0]+11	sleepq_wakeone_chan
45	0%	100%	0.00	5209	cpu[1]+11	lock_try
39	0%	100%	0.00	4024	cpu[3]+11	lock_set_spl
33	0%	100%	0.00	5156	cpu[2]+11	setbackdq
30	0%	100%	0.00	3790	cpu[3]+2	dosoftint
6	0%	100%	0.00	5600	cpu[1]+5	ddi_io_getb
3	0%	100%	0.00	1072	cpu[0]+2	apic_redistribute_compute

```
# dtrace -n 'profile-997ms / arg0 != 0 / { @ks[stack()]=count() }'
dtrace: description 'profile-997ms' matched 1 probe
^C
```

```
genunix`syscall_mstate+0x1c7
unix`sys_syscall32+0xbd
    1

unix`bzero+0x3
procfs`pr_read_lwpusage_32+0x2f
procfs`prread+0x5d
genunix`fop_read+0x29
genunix`pread+0x217
genunix`pread32+0x26
unix`sys_syscall32+0x101
    1
```



[Continue from previous slide - dtrace stack() aggregation output...]

```
....  
unix`kcopy+0x38  
genunix`copyin_nowatch+0x48  
genunix`copyin_args32+0x45  
genunix`syscall_entry+0xcb  
unix`sys_syscall32+0xe1  
    1  
  
unix`sys_syscall32+0xae  
    1  
  
unix`mutex_exit+0x19  
ufs`rdip+0x368  
ufs`ufs_read+0x1a6  
genunix`fop_read+0x29  
genunix`pread64+0x1d7  
unix`sys_syscall32+0x101  
    2  
  
unix`kcopy+0x2c  
genunix`uiomove+0x17f  
ufs`rdip+0x382  
ufs`ufs_read+0x1a6  
genunix`fop_read+0x29  
genunix`pread64+0x1d7  
unix`sys_syscall32+0x101  
    13
```



# Another Kernel Stack View

```
# lockstat -i997 -Ikws 10 sleep 30  
  
Profiling interrupt: 119800 events in 30.038 seconds (3988 events/sec)
```

---

Count	indv	cuml	rcnt	nsec	CPU+PIL	Hottest Caller
29919	25%	25%	0.00	5403	cpu[2]	kcopy

---

nsec	----- Time Distribution -----	count	Stack
1024	@@@	2	uiomove
2048	@@@@	18	rdip
4096	@@@@@	25	ufs_read
8192	@@@@@@@	29853	fop_read
16384	@@@@@@@	21	pread64
			sys_syscall132

---

Count	indv	cuml	rcnt	nsec	CPU+PIL	Hottest Caller
29918	25%	50%	0.00	5386	cpu[1]	kcopy

---

nsec	----- Time Distribution -----	count	Stack
4096	@@@	38	uiomove
8192	@@@@@	29870	rdip
16384	@@@@@	10	ufs_read
			fop_read
			pread64
			sys_syscall132

---

Count	indv	cuml	rcnt	nsec	CPU+PIL	Hottest Caller
29893	25%	75%	0.00	5283	cpu[3]	kcopy

---

nsec	----- Time Distribution -----	count	Stack
1024	@@@	140	uiomove
2048	@@@	761	rdip
4096	@	1443	ufs_read
8192	@@@@@	27532	fop_read
16384	@@@@@	17	pread64
			sys_syscall132

---



# Who's Doing What...

```
#prstat -Lmc 10 10 > prstat.out  
#cat prstat.out
```

PID	USERNAME	USR	SYS	TRP	TFL	DFL	LCK	SLP	LAT	VCX	ICX	SCL	SIG	PROCESS/LWPID
4448	root	12	44	0.0	0.0	0.0	0.0	0.0	43	0.5	2K	460	.1M	0 prstat/1
4447	root	1.2	11	0.0	0.0	0.0	0.1	0.1	14	73	54	65	.2M	0 filebench/27
4447	root	1.1	10	0.0	0.0	0.0	0.0	0.1	15	74	57	52	.2M	0 filebench/29
4447	root	1.1	10	0.0	0.0	0.0	0.1	0.0	15	74	64	53	.2M	0 filebench/19
4447	root	1.1	10	0.0	0.0	0.0	0.0	0.4	14	74	49	55	.2M	0 filebench/7
4447	root	1.1	10	0.0	0.0	0.0	0.0	0.2	14	74	51	44	.2M	0 filebench/17
4447	root	1.1	9.9	0.0	0.0	0.0	0.0	0.3	14	74	48	57	.2M	0 filebench/14
4447	root	1.1	9.9	0.0	0.0	0.0	0.0	0.3	14	74	42	61	.2M	0 filebench/9
4447	root	1.1	9.8	0.0	0.0	0.0	0.0	0.1	15	74	51	49	.2M	0 filebench/25
4447	root	1.1	9.8	0.0	0.0	0.0	0.0	0.0	15	74	60	38	.2M	0 filebench/4
4447	root	1.1	9.7	0.0	0.0	0.0	0.0	0.2	14	75	25	69	.2M	0 filebench/26
4447	root	1.0	9.7	0.0	0.0	0.0	0.1	0.0	15	75	54	46	.2M	0 filebench/12
4447	root	1.1	9.6	0.0	0.0	0.0	0.0	0.3	14	75	40	46	.2M	0 filebench/21
4447	root	1.1	9.6	0.0	0.0	0.0	0.0	0.1	15	75	39	70	.2M	0 filebench/31
4447	root	1.1	9.6	0.0	0.0	0.0	0.1	0.0	15	75	38	75	.2M	0 filebench/22

Total: 59 processes, 218 lwps, load averages: 9.02, 14.30, 10.36

PID	USERNAME	USR	SYS	TRP	TFL	DFL	LCK	SLP	LAT	VCX	ICX	SCL	SIG	PROCESS/LWPID
4447	root	1.3	12	0.0	0.0	0.0	0.0	0.0	86	43	41	.3M	0	filebench/16
4447	root	1.3	12	0.0	0.0	0.0	0.0	0.0	87	35	46	.3M	0	filebench/14
4447	root	1.3	12	0.0	0.0	0.0	0.0	0.0	87	36	60	.3M	0	filebench/7
4447	root	1.3	12	0.0	0.0	0.0	0.0	0.0	87	27	44	.3M	0	filebench/24
4447	root	1.3	12	0.0	0.0	0.0	0.0	0.0	87	41	61	.3M	0	filebench/3
4447	root	1.3	12	0.0	0.0	0.0	0.0	0.0	87	38	49	.3M	0	filebench/13
4447	root	1.3	12	0.0	0.0	0.0	0.0	0.0	87	14	71	.3M	0	filebench/2
4447	root	1.3	12	0.0	0.0	0.0	0.0	0.0	87	32	57	.3M	0	filebench/19
4447	root	1.3	12	0.0	0.0	0.0	0.0	0.0	87	31	57	.3M	0	filebench/27
4447	root	1.3	12	0.0	0.0	0.0	0.0	0.0	87	34	47	.3M	0	filebench/4
4447	root	1.3	11	0.0	0.0	0.0	0.0	0.0	87	21	74	.3M	0	filebench/26
4447	root	1.2	11	0.0	0.0	0.0	0.0	0.0	87	42	51	.3M	0	filebench/9
4447	root	1.3	11	0.0	0.0	0.0	0.0	0.0	87	16	83	.3M	0	filebench/18
4447	root	1.2	11	0.0	0.0	0.0	0.0	0.0	87	42	47	.3M	0	filebench/33
4447	root	1.2	11	0.0	0.0	0.0	0.0	0.0	87	15	76	.3M	0	filebench/15

Total: 59 processes, 218 lwps, load averages: 12.54, 14.88, 10.59



# System Calls – What & Who

```
# dtrace -n 'syscall:::entry { @sc[probefunc]=count() }'  
dtrace: description 'syscall:::entry' matched 228 probes  
^C
```

fstat	1
mmap	1
schedctl	1
waitsys	1
recvmsg	2
sigaction	2
sysconfig	3
brk	6
pset	9
gtime	16
lwp_park	20
p_online	21
setcontext	29
write	30
nanosleep	32
lwp_sigmask	45
setitimer	54
pollsys	118
ioctl	427
pread64	1583439
pread	3166885
read	3166955

```
# dtrace -n 'syscall:::read:entry { @[execname,pid]=count() }'  
dtrace: description 'syscall:::read:entry' matched 1 probe  
^C
```

sshd	4342	3
Xorg	536	36
filebench	4376	2727656



# smtx – Lock Operations

```
# lockstat sleep 30 > lockstat.locks1
# more lockstat.locks1
```

Adaptive mutex spin: 3486197 events in 30.031 seconds (116088 events/sec)

Count	indv	cuml	rcnt	spin	Lock	Caller
1499963	43%	43%	0.00		84 pr_pidlock	pr_p_lock+0x29
1101187	32%	75%	0.00		24 0xffffffff810cdec0	pr_p_lock+0x50
285012	8%	83%	0.00		27 0xffffffff827a9858	rdip+0x506
212621	6%	89%	0.00		29 0xffffffff827a9858	rdip+0x134
98531	3%	92%	0.00	103	0xffffffff9321d480	releasef+0x55
92486	3%	94%	0.00	19	0xffffffff8d5c4990	ufs_lockfs_end+0x81
89404	3%	97%	0.00	27	0xffffffff8d5c4990	ufs_lockfs_begin+0x9f
83186	2%	99%	0.00	96	0xffffffff9321d480	getf+0x5d
6356	0%	99%	0.00	186	0xffffffff810cdec0	clock+0x4e9
1164	0%	100%	0.00	141	0xffffffff810cdec0	post_syscall+0x352
294	0%	100%	0.00	11	0xffffffff801a4008	segmap_smapadd+0x77
279	0%	100%	0.00	11	0xffffffff801a41d0	segmap_getmapflt+0x275
278	0%	100%	0.00	11	0xffffffff801a48f0	segmap_smapadd+0x77
276	0%	100%	0.00	11	0xffffffff801a5010	segmap_getmapflt+0x275
276	0%	100%	0.00	11	0xffffffff801a4008	segmap_getmapflt+0x275

...

Adaptive mutex block: 3328 events in 30.031 seconds (111 events/sec)

Count	indv	cuml	rcnt	nsec	Lock	Caller
1929	58%	58%	0.00	48944759	pr_pidlock	pr_p_lock+0x29
263	8%	66%	0.00	47017	0xffffffff810cdec0	pr_p_lock+0x50
255	8%	74%	0.00	53392369	0xffffffff9321d480	getf+0x5d
217	7%	80%	0.00	26133	0xffffffff810cdec0	clock+0x4e9
207	6%	86%	0.00	227146	0xffffffff827a9858	rdip+0x134
197	6%	92%	0.00	64467	0xffffffff8d5c4990	ufs_lockfs_begin+0x9f
122	4%	96%	0.00	64664	0xffffffff8d5c4990	ufs_lockfs_end+0x81
112	3%	99%	0.00	164559	0xffffffff827a9858	rdip+0x506



# smtx – Lock Operations (cont)

Spin lock spin: 3491 events in 30.031 seconds (116 events/sec)

Count	indv	cuml	rcnt	spin Lock	Caller
2197	63%	63%	0.00	2151 turnstile_table+0xbd8	disp_lock_enter+0x35
314	9%	72%	0.00	3129 turnstile_table+0xe28	disp_lock_enter+0x35
296	8%	80%	0.00	3162 turnstile_table+0x888	disp_lock_enter+0x35
211	6%	86%	0.00	2032 turnstile_table+0x8a8	disp_lock_enter+0x35
127	4%	90%	0.00	856 turnstile_table+0x9f8	turnstile_interlock+0x171
114	3%	93%	0.00	269 turnstile_table+0x9f8	disp_lock_enter+0x35
44	1%	95%	0.00	90 0xffffffff827f4de0	disp_lock_enter_high+0x13
37	1%	96%	0.00	581 0xffffffff827f4de0	disp_lock_enter+0x35

...

Thread lock spin: 1104 events in 30.031 seconds (37 events/sec)

Count	indv	cuml	rcnt	spin Lock	Caller
487	44%	44%	0.00	1671 turnstile_table+0xbd8	ts_tick+0x26
219	20%	64%	0.00	1510 turnstile_table+0xbd8	turnstile_block+0x387
92	8%	72%	0.00	1941 turnstile_table+0x8a8	ts_tick+0x26
77	7%	79%	0.00	2037 turnstile_table+0xe28	ts_tick+0x26
74	7%	86%	0.00	2296 turnstile_table+0x888	ts_tick+0x26
36	3%	89%	0.00	292 cpu[0]+0xf8	ts_tick+0x26
27	2%	92%	0.00	55 cpu[1]+0xf8	ts_tick+0x26
11	1%	93%	0.00	26 cpu[3]+0xf8	ts_tick+0x26
10	1%	94%	0.00	11 cpu[2]+0xf8	post_syscall+0x556

...



R/W writer blocked by writer: 17 events in 30.031 seconds (1 events/sec)

Count	indv	cuml	rcnt	nsec	Lock	Caller
17	100%	100%	0.00	465308	0xffffffff831f3be0	ufs_getpage+0x369

R/W writer blocked by readers: 55 events in 30.031 seconds (2 events/sec)

Count	indv	cuml	rcnt	nsec	Lock	Caller
55	100%	100%	0.00	1232132	0xffffffff831f3be0	ufs_getpage+0x369

R/W reader blocked by writer: 22 events in 30.031 seconds (1 events/sec)

Count	indv	cuml	rcnt	nsec	Lock	Caller
18	82%	82%	0.00	56339	0xffffffff831f3be0	ufs_getpage+0x369
4	18%	100%	0.00	45162	0xffffffff831f3be0	ufs_putpages+0x176

R/W reader blocked by write wanted: 47 events in 30.031 seconds (2 events/sec)

Count	indv	cuml	rcnt	nsec	Lock	Caller
46	98%	98%	0.00	369379	0xffffffff831f3be0	ufs_getpage+0x369
1	2%	100%	0.00	118455	0xffffffff831f3be0	ufs_putpages+0x176



# Chasing the hot lock caller...

```
# dtrace -n 'pr_p_lock:entry { @s[stack()]=count() }'
dtrace: description 'pr_p_lock:entry' matched 1 probe
^C
    procfs`pr_read_lwpusage_32+0x4f
    procfs`prread+0x5d
    genunix`fop_read+0x29
    genunix`pread+0x217
    genunix`pread32+0x26
    unix`sys_syscall32+0x101
    12266066

# dtrace -n 'pr_p_lock:entry { @s[execname]=count() }'
dtrace: description 'pr_p_lock:entry' matched 1 probe
^C

filebench                                     8439499
# pgrep filebench
4485
# dtrace -n 'pid4485:libc:pread:entry { @us[ustack()]=count() }'
dtrace: description 'pid4485:libc:pread:entry' matched 1 probe
^C
    libc.so.1`pread
    filebench`flowop_endop+0x5b
    filebench`flowoplib_read+0x238
    filebench`flowop_start+0x2b1
    libc.so.1`_thr_setup+0x51
    libc.so.1`_lwp_start
    2084651

    libc.so.1`pread
    filebench`flowop_beginop+0x6a
    filebench`flowoplib_read+0x200
    filebench`flowop_start+0x2b1
    libc.so.1`_thr_setup+0x51
    libc.so.1`_lwp_start
    2084651
```



# Icing on the cake...

```
# dtrace -q -n 'ufs_read:entry { printf("UFS Read: %s\n",stringof(args[0]->v_path)); }'
UFS Read: /ufs/largefile1
^C

#
#
# dtrace -q -n 'ufs_read:entry { @[execname,stringof(args[0]->v_path)]=count() }'
^C

filebench                               /ufs/largefile1
864609
```



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# Session 2 - Memory

# Virtual Memory

- Simple programming model/abstraction
- Fault Isolation
- Security
- Management of Physical Memory
- Sharing of Memory Objects
- Caching



# Solaris Virtual Memory Glossary

Address Space	Linear memory range visible to a program, that the instructions of the program can directly load and store. Each Solaris process has an address space; the Solaris kernel also has its own address space.
Virtual Memory	Illusion of real memory within an address space.
Physical Memory	Real memory (e.g. RAM)
Mapping	A memory relationship between the address space and an object managed by the virtual memory system.
Segment	A co-managed set of similar mappings within an address space.
Text Mapping	The mapping containing the program's instructions and read-only objects.
Data Mapping	The mapping containing the program's initialized data
Heap	A mapping used to contain the program's heap (malloc'd) space
Stack	A mapping used to hold the program's stack
Page	A linear chunk of memory managed by the virtual memory system
VNODE	A file-system independent file object within the Solaris kernel
Backing Store	The storage medium used to hold a page of virtual memory while it is not backed by physical memory
Paging	The action of moving a page to or from its backing store
Swapping	The action of swapping an entire address space to/from the swap device
Swap Space	A storage device used as the backing store for anonymous pages.



# Solaris Virtual Memory Glossary (cont)

Scanning	The action of the virtual memory system takes when looking for memory which can be freed up for use by other subsystems.
Named Pages	Pages which are mappings of an object in the file system.
Anonymous Memory	Pages which do not have a named backing store
Protection	A set of booleans to describe if a program is allowed to read, write or execute instructions within a page or mapping.
ISM	Intimate Shared Memory - A type of System V shared memory optimized for sharing between many processes
DISM	Pageable ISM
NUMA	Non-uniform memory architecture - a term used to describe a machine with differing processor-memory latencies.
Lgroup	A locality group - a grouping of processors and physical memory which share similar memory latencies
MMU	The hardware functional unit in the microprocessor used to dynamically translate virtual addresses into physical addresses.
HAT	The Hardware Address Translation Layer - the Solaris layer which manages the translation of virtual addresses to physical addresses
TTE	Translation Table Entry - The UltraSPARC hardware's table entry which holds the data for virtual to physical translation
TLB	Translation Lookaside Buffer - the hardware's cache of virtual address translations
Page Size	The translation size for each entry in the TLB
TSB	Translation Software Buffer - UltraSPARC's software cache of TTEs, used for lookup when a translation is not found in the TLB



# Solaris Virtual Memory

- Demand Paged, Globally Managed
- Integrated file caching
- Layered to allow virtual memory to describe multiple memory types (Physical memory, frame buffers)
- Layered to allow multiple MMU architectures



USE

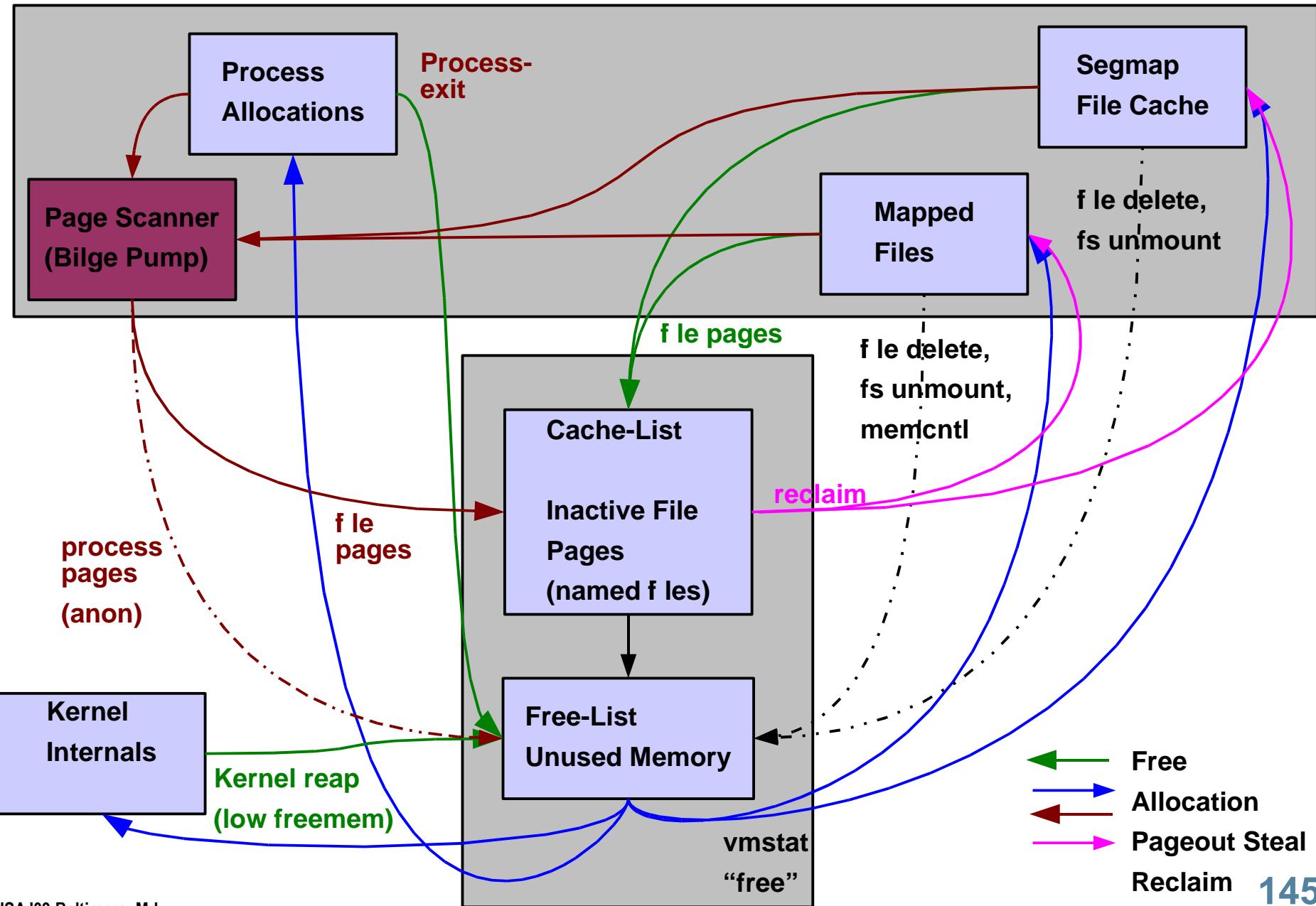


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# Physical Memory Management

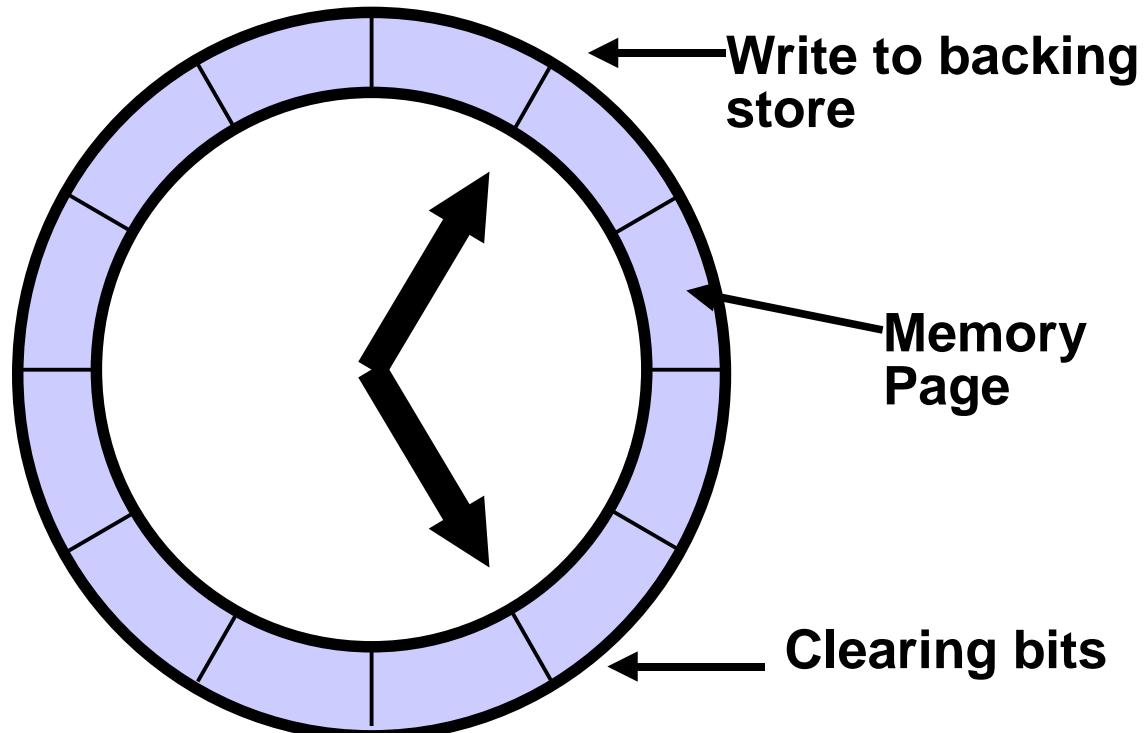


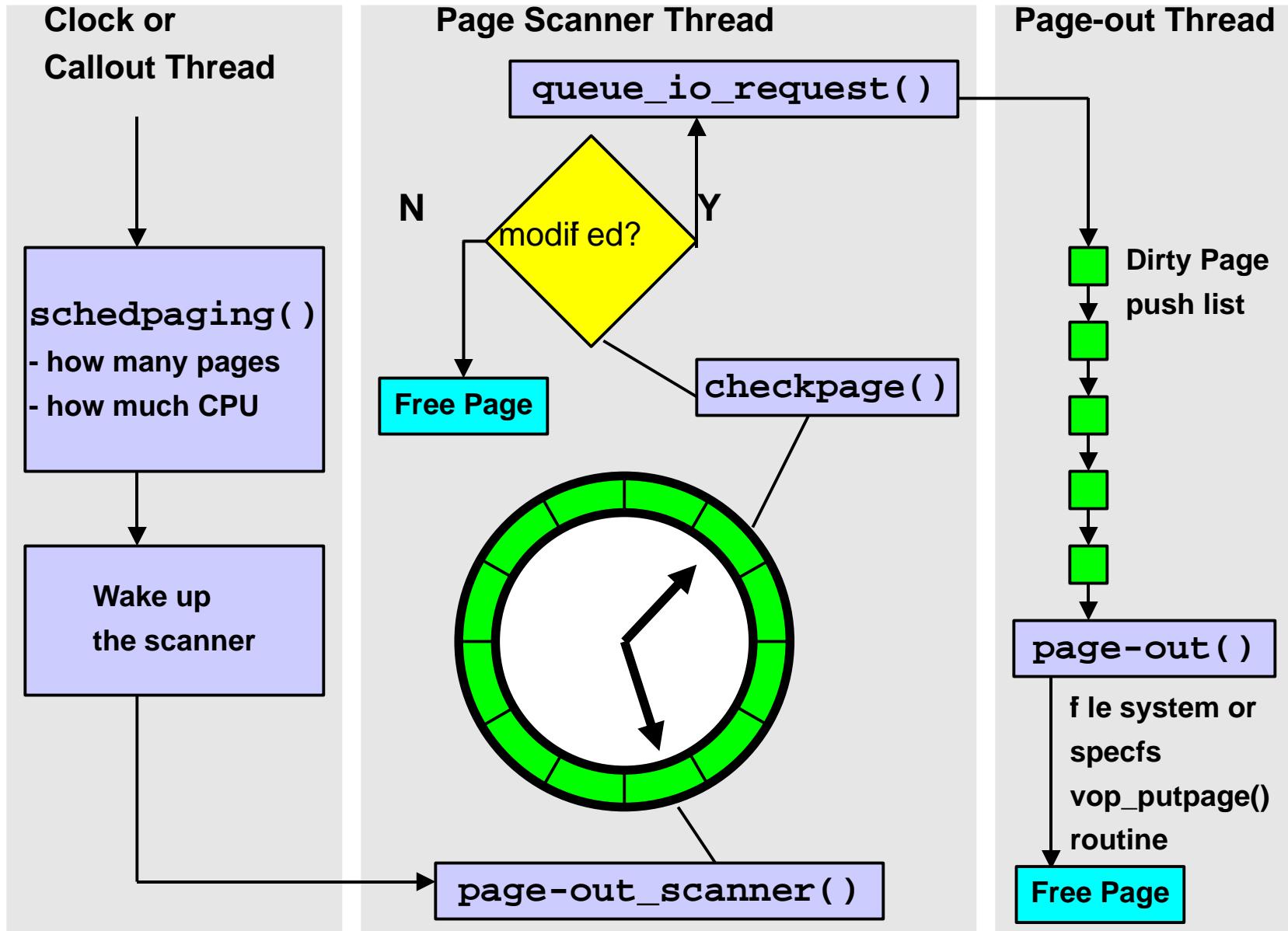
# Page Lists

- Free List
  - does not have a vnode/offset associated
  - put on list at process exit.
  - may be always small (pre Solaris 8)
- Cache List
  - still have a vnode/offset
  - seg\_map free-behind and seg\_vn executables and libraries (for reuse)
  - reclaims are in `vmstat "re"`
- Sum of these two are in `vmstat "free"`

# Page Scanning

- Steals pages when memory is low
- Uses a Least Recently Used process.
- Puts memory out to "backing store"
- Kernel thread does the scanning





# Scanning Algorithm

- Free memory is lower than (lotsfree)
- Starts scanning @ slowscan (pages/sec)
- Scanner Runs:
  - four times / second when memory is short
  - Awoken by page allocator if very low
- Limits:
  - Max # of pages /sec. swap device can handle
  - How much CPU should be used for scanning

$$\text{scanrate} = \left( \frac{\text{lotsfree} - \text{freemem}}{\text{lotsfree}} \times \text{fastscan} \right) + \left( \text{slowscan} \times \frac{\text{freemem}}{\text{lotsfree}} \right)$$

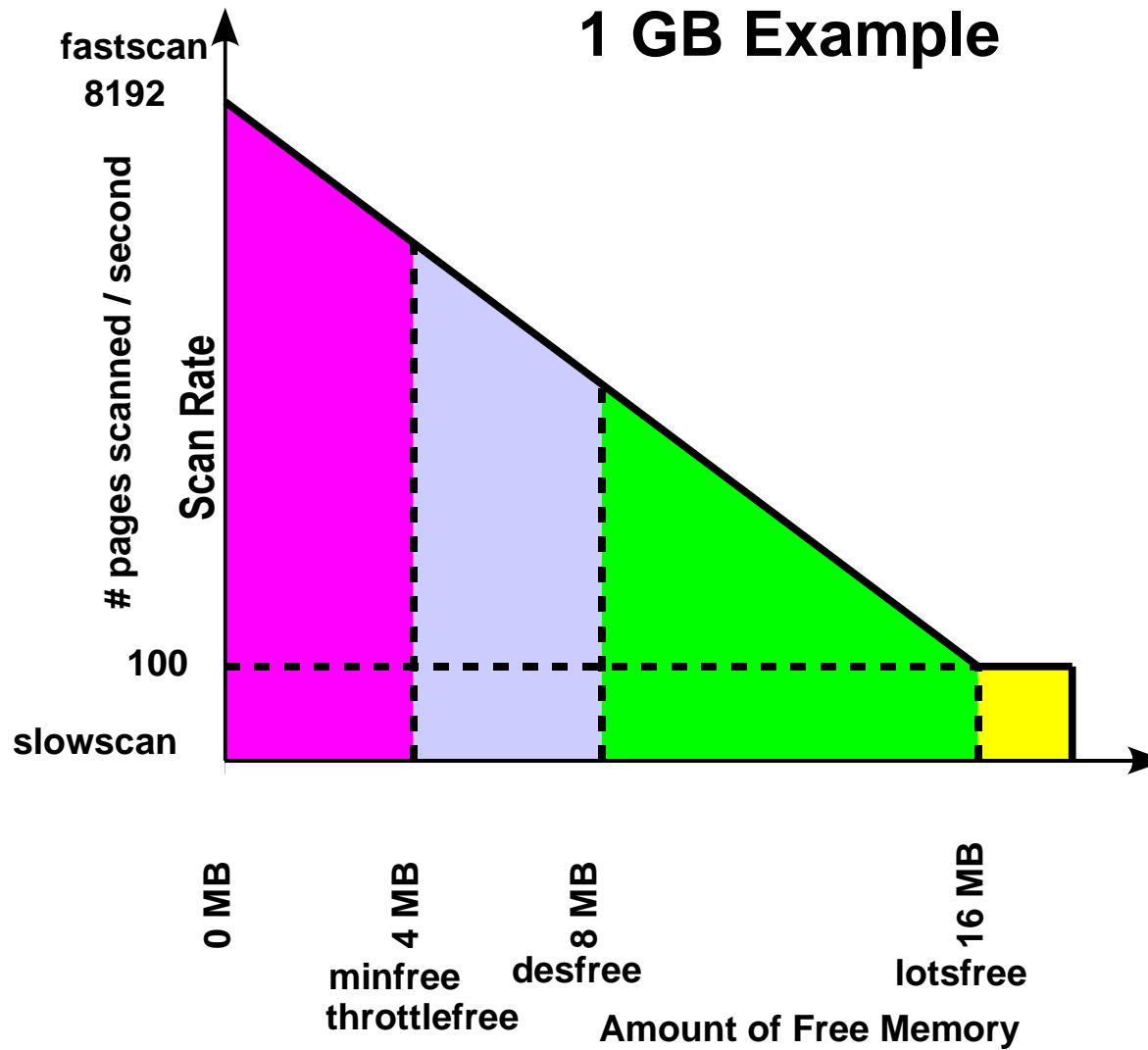


# Scanning Parameters

Parameter	Description	Min	Default ( Solaris 8)
lotsfree	starts stealing anonymous memory pages	512K	1/64 th of memory
desfree	scanner is started at 100 times/second	minfree	1/2 of lotsfree
minfree	start scanning every time a new page is created		1/2 of desfree
throttlefree	page_create routine makes the caller wait until free pages are available		minfree
fastscan	scan rate (pages per second) when free memory = minfree	slowscan	minimum of 64MB/s or 1/2 memory size
slowscan	scan rate (pages per second) when free memory = lotsfree		100
maxpgio	max number of pages per second that the swap device can	~60	60 or 90 pages per spindle
hand-spreadpages	number of pages between the front hand (clearing) and back hand (checking)	1	fastscan
min_percent_cpu	CPU usage when free memory is at lotsfree	4% (~1 clock tick)	of a single CPU

# Scan Rate

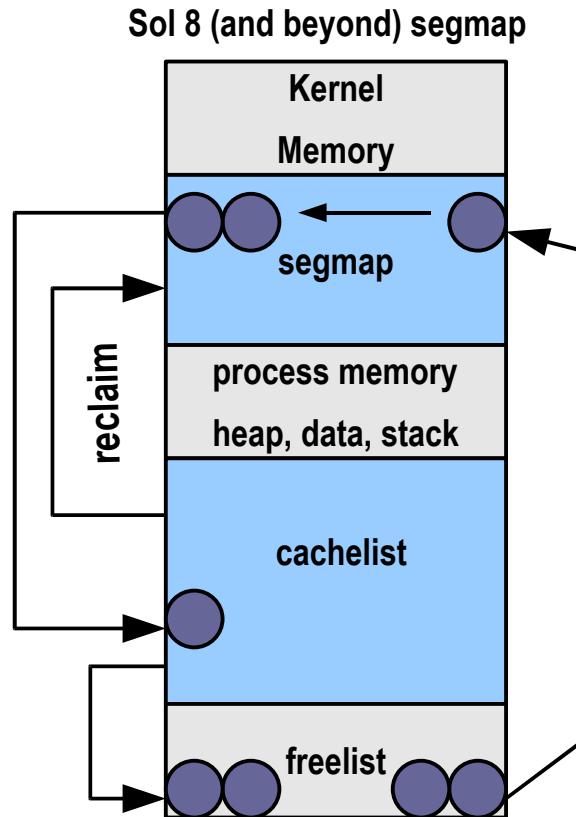
1 GB Example



# The Solaris Page Cache

- Page list is broken into two:
  - Cache List: pages with a valid vnode/offset
  - Free List: pages has no vnode/offset
- Unmapped pages where just released
- Non-dirty pages, not mapped, should be on the "free list"
- Places pages on the "tail" cache/free list
- Free memory = cache + free
- UFS
  - segmap kernel address space segment
  - Starting in Solaris 10 3/05, segkpm integration (SPARC)
- ZFS
  - Uses kernel memory (`kmem_alloc`) for ARC cache

# The Solaris UFS Cache- segmap





# The Solaris Cache

- Now vmstat reports a useful free
- Throw away your old /etc/system pager configuration parameters
  - lotsfree, desfree, minfree
  - fastscan, slowscan
  - priority\_paging, cachefree



# Memory Summary

## Physical Memory:

```
# prtconf  
System Configuration: Sun Microsystems sun4u  
Memory size: 512 Megabytes
```

## Kernel Memory:

```
# sar -k 1 1  
SunOS ian 5.8 Generic_108528-03 sun4u      08/28/01  
13:04:58 sml_mem    alloc   fail  lg_mem    alloc   fail  ovsz_alloc   fail  
13:04:59 10059904 7392775      0 133349376 92888024      0      10346496      0
```

## Free Memory:

```
# vmstat 3 3  
procs      memory          page          disk          faults          cpu  
r b w    swap   free   re   mf pi po fr de sr f0 s0 s1 s6   in   sy   cs us sy id  
0 0 0 478680 204528  0    2  0  0   0  0  0  0  0  1  0 209 1886 724 35  5 61  
0 0 0 415184 123400  0    2  0  0   0  0  0  0  0  0  0 238 825 451  2  1 98  
0 0 0 415200 123416  0    0  0  0   0  0  0  0  0  0  3  0 219 788 427  1  1 98
```



# Memory Summary

```
nv70b> kstat unix:0:system_pages:pp_kernel
module: unix                                instance: 0
name:   system_pages                         class:    pages
        pp_kernel                           53610

nv70b> pagesize
4096
nv70b> bc -l
4096*53610
219586560
nv70b> su
Password:
# mdb -k
Loading modules: [ unix genunix specfs dtrace uppcc pcplusmp scsi_vhci ufs mpt ip
> ::memstat
Page Summary          Pages      MB %Tot
-----  -----
Kernel                50377      196  19%
Anon                  63005      246  24%
Exec and libs         13689       53   5%
Page cache             8871       34   3%
Free (cachelist)     52344      204  20%
Free (freelist)       71696      280  28%

Total                 259982     1015
Physical              259981     1015
```

# vmstat

r = run queue length

b = processes blocked waiting for I/O

w = idle processes that have been swapped at some time

swap = free and unreserved swap in KBytes

free = free memory measured in pages

re = kilobytes reclaimed from cache/free list

mf = minor faults - the page was in memory but was not mapped

pi = kilobytes paged-in from the file system or swap device

po = kilobytes paged-out to the file system or swap device

fr = kilobytes that have been destroyed or freed

de = kilobytes freed after writes

sr = pages scanned / second = disk I/Os per second for disk 0-3

in = interrupts / second

sy = system calls / second

cs = context switches / second

us = user cpu time

sy = kernel cpu time

id = idle + wait cpu time

		memory				page				disk				faults			cpu			
procs		swap	free	re	mf	pi	po	fr	de	sr	f0	s0	s1	s2	in	sy	cs	us	sy	id
0	0	0	46580232	337472	18	194	30	0	0	0	0	0	0	0	5862	81260	28143	19	7	74
0	0	0	45311368	336280	32	249	48	0	0	0	0	0	0	0	6047	93562	29039	21	10	69
0	0	0	46579816	337048	12	216	60	0	0	0	0	10	0	7	5742	100944	27032	20	7	73
0	0	0	46580128	337176	3	111	3	0	0	0	0	0	0	0	5569	93338	26204	21	6	73



# vmstat -p

swap = free and unreserved swap in KBytes

free = free memory measured in pages

re = kilobytes reclaimed from cache/free list

mf = minor faults - the page was in memory but was not mapped

fr = kilobytes that have been destroyed or freed

de = kilobytes freed after writes

sr = kilobytes scanned / second

executable pages: kilobytes in - out - freed

anonymous pages: kilobytes in - out - freed

file system pages:  
kilobytes in - out - freed

# vmstat -p 5 5

memory		page					executable			anonymous			filesystem		
swap	free	re	mf	fr	de	sr	epi	epo	epf	api	apo	apf	fpi	fpo	fpf
46715224	891296	24	350	0	0	0	0	0	0	4	0	0	27	0	0
46304792	897312	151	761	25	0	0	17	0	0	1	0	0	280	25	25
45886168	899808	118	339	1	0	0	3	0	0	1	0	0	641	1	1
46723376	899440	29	197	0	0	0	0	0	0	40	0	0	60	0	0

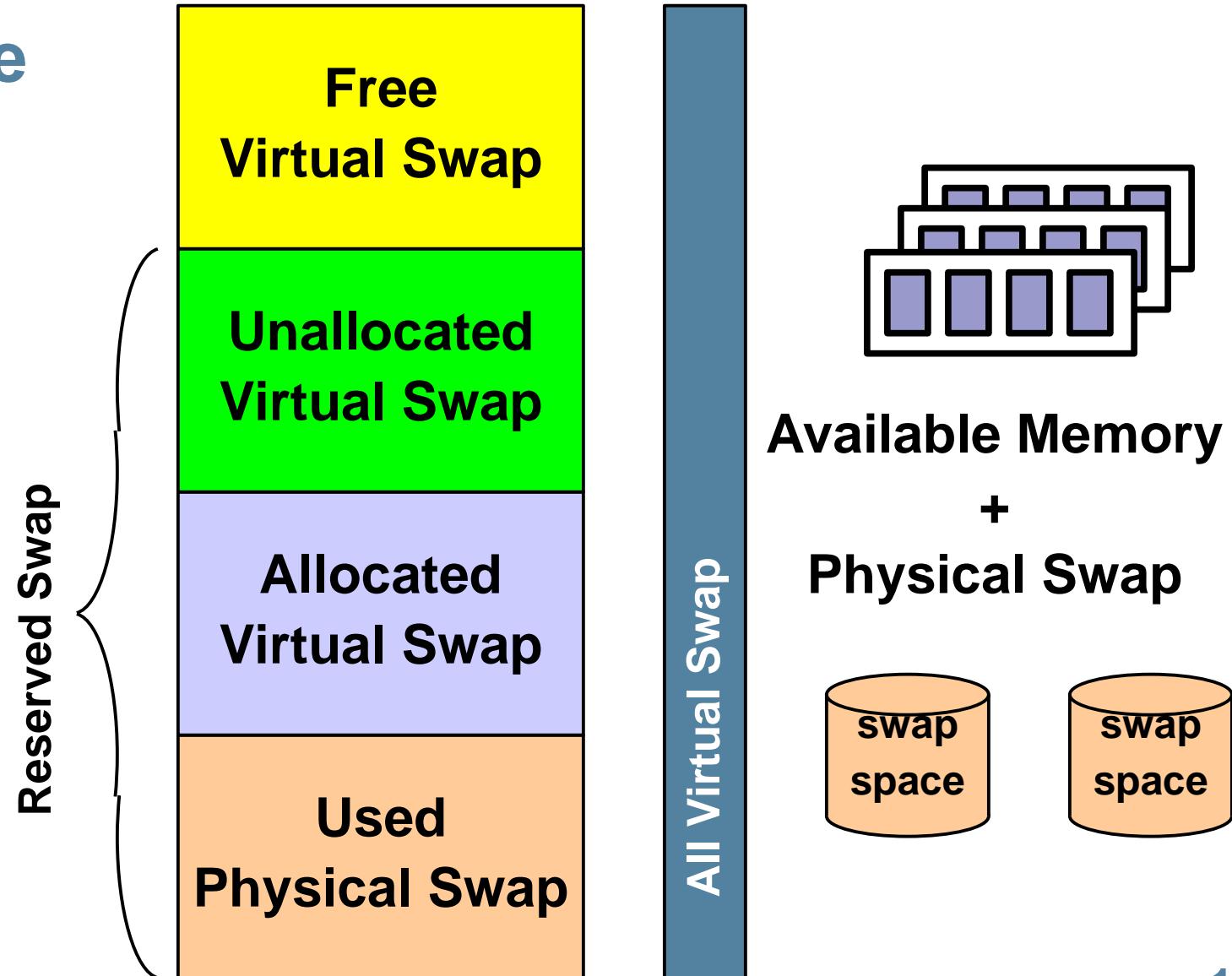
# Swapping

- Scheduler/Dispatcher:
  - Dramatically affects process performance
  - Used when demand paging is not enough
- Soft swapping:
  - Avg. freemem below desfree for 30 sec.
  - Look for inactive processes, at least `maxslp`
- Hard swapping:
  - Run queue  $\geq 2$  (waiting for CPU)
  - Avg. freemem below desfree for 30 sec.
  - Excessive paging, (`pageout + pagein`)  $>$  `maxpgio`
  - Aggressive; unload kernel mods & free cache

# Swap space states

- Reserved:
  - Virtual space is reserved for the segment
  - Represents the virtual size being created
- Allocated:
  - Virtual space is allocated when the first physical page is assigned
  - A swapfs vnode / offset are assigned
- Swapped out:
  - When a shortage occurs
  - Page is swapped out by the scanner, migrated to swap storage

# Swap Space





# Swap Usage

- Virtual Swap:
  - reserved: unallocated + allocated
  - available = bytes

```
# swap -s
total: 175224k bytes unallocated + 24464k allocated = 199688k reserved, 416336k
available
```

- Physical Swap:
  - space available for physical page-outs
  - free = blocks (512 bytes)

```
# swap -l
swapfile          dev  swaplo blocks   free
/dev/dsk/c0t1d0s1  32,9    16 524864 524864
```

- Ensure both are non-zero
  - swap -s "available"
  - swap -l "free"

# A Quick Guide to Analyzing Memory

- Quick Memory Health Check
  - Check free memory and scanning with vmstat
  - Check memory usage with ::memstat in mdb
- Paging Activity
  - Use vmstat -p to check if there are anonymous page-ins
- Attribution
  - Use DTrace to see which processes/files are causing paging
- Time based analysis
  - Use DTrace to estimate the impact of paging on application performance
- Process Memory Usage
  - Use pmap to inspect process memory usage and sharing
- MMU/Page Size Performance
  - Use trapstat to observe time spent in TLB misses



# Memory Kstats – via kstat(1m)

```
sol8# kstat -n system_pages
module: unix
name: system_pages
availrmem
crttime
desfree
desscan
econtig
fastscan
freemem
kernelbase
lotsfree
minfree
nalloc
nalloc_calls
nfree
nfree_calls
nscan
pagesfree
pageslocked
pagestotal
physmem
pp_kernel
slowscan
snaptime
instance: 0
class: pages
343567
0
4001
25
4278190080
256068
248309
3556769792
8002
2000
11957763
9981
11856636
6689
0
248309
168569
512136
522272
64102
100
6573953.83957897
```



# Memory Kstats – via kstat Perl API

```
%{$now} = %{$kstats->{0}{system_pages}};  
print "$now->{pagesfree}\n";  
  
sol8# wget http://www.solarisinternals.com/si/downloads/prtmem.pl  
sol8# prtmem.pl 10  
prtmem started on 04/01/2005 15:46:13 on devnull, sample interval 5  
seconds  
15:46:18          Total       Kernel      Delta      Free      Delta  
15:46:23          2040        250          0         972        -12  
15:46:28          2040        250          0         968          -3  
15:46:33          2040        250          0         968          0  
                                0           0           0         970          1
```



# Checking Paging Activity

- Good Paging
  - Plenty of memory free
  - Only file system page-in/page-outs (vmstat: fpi, fpo > 0)

```
%sol8# vmstat -p 3
      memory          page        executable       anonymous    filesystem
 swap   free   re   mf   fr   de   sr   epi   epo   epf   api   apo   apf   fpi   fpo   fpf
 1512488 837792 160 20 12   0   0     0     0     0     0     0     0     0     12    12    12
 1715812 985116  7  82   0   0   0     0     0     0     0     0     0     0     45    0     0
 1715784 983984  0   2   0   0   0     0     0     0     0     0     0     0     53    0     0
 1715780 987644  0   0   0   0   0     0     0     0     0     0     0     0     33    0     0
```



# Checking Paging Activity

- Bad Paging
  - Non zero Scan rate (vmstat: sr >0)
  - Low free memory (vmstat: free < 1/16<sup>th</sup> physical)
  - Anonymous page-in/page-outs (vmstat: api, apo > 0)

```
sol8# vmstat -p 3
      memory          page        executable      anonymous      filesystem
 swap   free   re   mf   fr   de   sr   epi   epo   epf   api   apo   apf   fpi   fpo   fpf
 2276000 1589424 2128 19969 1 0 0    0    0    0    0    0    0    0    1    1
 1087652 388768 12 129675 13879 0 85590 0 0 12    0 3238 3238 10 9391 10630
 608036 51464 20 8853 37303 0 65871 38    0 781 12 19934 19930 95 16548 16591
 94448 8000 17 23674 30169 0 238522 16 0 810 23 28739 28804 56 547 556
```

# Using prstat to estimate paging slow-downs

- Microstates show breakdown of elapsed time
  - prstat -m
  - USR through LAT columns summed show 100% of wallclock execution time for target thread/process
  - DFL shows time spent waiting in major faults in anon:

```
sol8$ prstat -mL
 PID USERNAME  USR   SYS   TRP   TFL   DFL   LCK   SLP   LAT   VCX   ICX   SCL   SIG   PROCESS/LWPID
15625 rmc      0.1   0.7   0.0   0.0    95   0.0   0.9   3.2   1K    726   88    0   filebench/2
15652 rmc      0.1   0.7   0.0   0.0    94   0.0   1.8   3.6   1K    1K    10    0   filebench/2
15635 rmc      0.1   0.7   0.0   0.0    96   0.0   0.5   3.2   1K    1K    8     0   filebench/2
15626 rmc      0.1   0.6   0.0   0.0    95   0.0   1.4   2.6   1K    813   10    0   filebench/2
15712 rmc      0.1   0.5   0.0   0.0    47   0.0   49    3.8   1K    831   104   0   filebench/2
15628 rmc      0.1   0.5   0.0   0.0    96   0.0   0.0   3.1   1K    735   4     0   filebench/2
15725 rmc      0.0   0.4   0.0   0.0    92   0.0   1.7   5.7   996   736   8     0   filebench/2
15719 rmc      0.0   0.4   0.0   0.0    40   40    17   2.9   1K    708   107   0   filebench/2
15614 rmc      0.0   0.3   0.0   0.0    92   0.0   4.7   2.4   874   576   40    0   filebench/2
```

# Using DTrace for memory Analysis

- The “vminfo” provider has probes at all the places memory statistics are gathered.
- Everything visible via vmstat -p and kstat are defined as probes
  - arg0: the value by which the statistic is to be incremented. For most probes, this argument is always 1, but for some it may take other values; these probes are noted in Table 5-4.
  - arg1: a pointer to the current value of the statistic to be incremented. This value is a 64 bit quantity that is incremented by the value in arg0. Dereferencing this pointer allows consumers to determine the current count of the statistic corresponding to the probe.



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# Using DTrace for Memory Analysis

- For example, if you should see the following paging activity with vmstat, indicating page-in from the swap device, you could drill down to investigate.

```
sol8# vmstat -p 3
      memory          page        executable        anonymous        filesystem
    swap   free   re   mf   fr   de   sr   epi   epo   epf   api   apo   apf   fpi   fpo   fpf
1512488 837792 160 20 12 0 0 0 0 0 0 8102 0 0 12 12 12 12
1715812 985116 7 82 0 0 0 0 0 0 0 7501 0 0 45 12 0 0
1715784 983984 0 2 0 0 0 0 0 0 0 1231 0 0 53 0 0 0
1715780 987644 0 0 0 0 0 0 0 0 0 2451 0 0 33 0 0 0
```

```
sol10$ dtrace -n anonpgin '{@[execname] = count()}'
```

```
dtrace: description anonpgin matched 1 probe
```

svc.startd	1
sshd	2
ssh	3
dtrace	6
vmstat	28
filebench	913



# dtrace pagefaults

```
# cat pf.d
#!/usr/sbin/dtrace -s

#pragma D option quiet

fbt:unix:pagefault:entry
{
    @st[execname] = count();
    self->pfst[execname] = timestamp
}
fbt:unix:pagefault:return
/ self->pfst[execname] /
{
    @pft[execname] = sum(timestamp - self->pfst[execname]);
    self->pfst[execname] = 0;
}
tick-10s
{
    printf("Pagefault counts by execname ...\\n");
    printa(@st);

    printf("\\nPagefault times (in nano's) by execname...\\n");
    printa(@pft);

    clear(@st);
    clear(@pft);
}
```

tracking pagefault entry  
and returns for counts  
and times



# dtrace pagefaults

```
# ./pf.d
Pagefault counts by execname ...
```

dtrace	93
java	1257
kstat	1588

```
Pagefault times (in nano's) by execname...
```

dtrace	798535
kstat	17576367
java	85760822

```
Pagefault counts by execname ...
```

dtrace	2
java	1272
kstat	1588

```
Pagefault times (in nano's) by execname...
```

dtrace	80192
kstat	18227212
java	75422709

```
^C
```



# Large Memory

- Large Memory in Perspective
- 64-bit Solaris
- 64-bit Hardware
- Solaris enhancements for Large Memory
- Large Memory Databases
- Configuring Solaris for Large Memory
- Using larger page sizes

## 64-bit Solaris

- LP64 Data Model
- 32-bit or 64-bit kernel, with 32-bit & 64-bit application support
  - 64-bit kernel only on SPARC
    - 32-bit apps no problem
    - Solaris 10 64-bit on AMD64 and Intel
- Comprehensive 32-bit application compatibility

# Why 64-bit for large memory?

- Extends the existing programming model to large memory
  - Beyond 4GB limit imposed by 32 bits
- Existing POSIX APIs extend to large data types (e.g. file offsets. file handle limits eliminated)
- Simple transition of existing source to 64-bits

# Developer Perspective

- Virtually unlimited address space
  - Data objects, files, large hardware devices can be mapped into virtual address space
  - 64-bit data types, parameter passing
  - Caching can be implemented in application, yielding much higher performance
- Small Overhead
- 64-bit on AMD64
  - Native 64-bit integer arithmetic
  - 16 general purpose registers (instead of 8)
  - optimized function call interface – register based arg passing
  - other instruction set optimizations



# Large Memory Configs

## Configuring Solaris

- fsflush uses too much CPU on Solaris 8
  - Set “autoup” in /etc/system
  - Symptom is one CPU using 100%sys
- Corrective Action
  - Default is 30s, recommend setting larger
  - e.g. 10x nGB of memory



# Large Dump Performance

- Configure “kernel only”
  - dumpadm(1m)
- Estimate dump as 20% of memory size
- Configure separate dump device
  - Reliable dumps
  - Asynchronous saves during boot (savecore)
- Configure a fast dump device
  - If possible, a HW RAID stripe dump device

# Databases

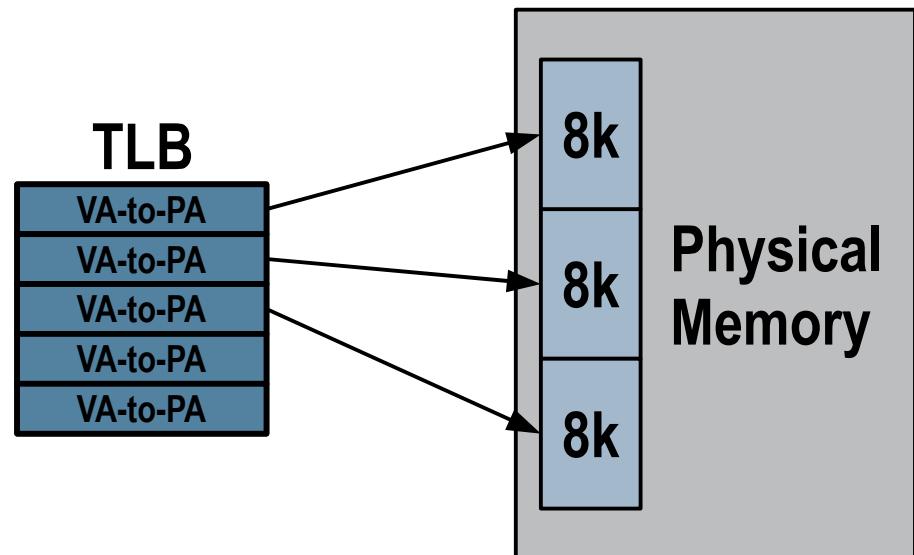
- Exploit memory to reduce/eliminate I/O!
- Eliminating I/O is the easiest way to tune it...
- Increase cache hit rates:
  - 95% means 1 out 20 accesses result in I/O
    - For every 1000 IOs, 50 are going to disk
  - 99% means 1 out of 100
    - For every 1000 IOs, 10 are going to disk
  - **That's a 5X (500%) reduction in physical disk IOs!**
- Use memory for caching
- Write-mostly I/O pattern results
  - Reads satisfied from cache

# Multiple Pagesize Support (MPSS) aka Large Pages

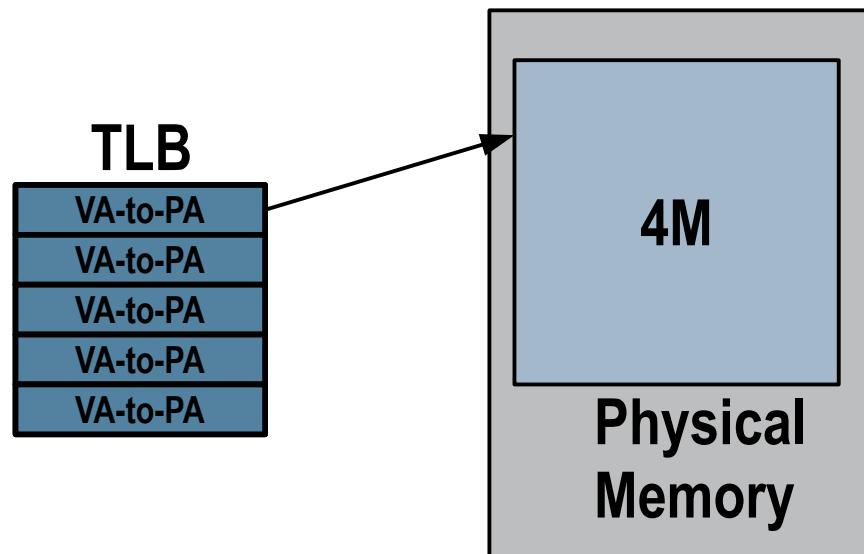
- Leverage hardware MMU support for multiple page sizes
- Supported page sizes will vary across different processors
  - `pagesize(1)`
- Functionality has been an ongoing effort, evolving over time
- Intended to improve performance through more efficient use of hardware TLB
- Be aware of cache effects of large pages (page coloring)
- For DR-capable systems, an interesting dynamic between kernel cage and large pages
  - cage-on: good for LP, may be not good for performance
  - cage-off: more memory fragmentation, not good for LP, but sometimes helps performance

# Why Large Pages?

address references from running threads →  
512 8k pages for a 4MB segment, versus one 4MB page



address references from running threads →



# Large Pages – A Brief History

- Solaris 2.6 – Solaris 8
  - SPARC: 4MB pages for ISM
  - SPARC: 4MB pages for initial kernel text and data segments
- Solaris 9
  - SPARC: 8k, 64k, 512k, 4M for user process anon, heap and stack via ppgsz(1), memcntl(2), mpss.so
  - SPARC: 4M for ISM / DISM
- Solaris 10 1/05
  - SPARC: Same as above
  - AMD64: 4k, 2M pages - same constraints as Solaris 9 SPARC
- Solaris 10 1/06 (Update 1)
  - SPARC: Added MPSS for regular file mappings (VMPSS) – enabled by default, 8k & 4M for sun4u, 8k, 64, 4M for sun4v
  - SPARC: Added Large Pages Out-Of-The-Box (LPOOB) for user process anon, stack and heap
  - SPARC: KPR integrated
  - AMD64: 2M for text can be enabled via /etc/system



## Large Pages – A Brief History (continued)

- Solaris 10 6/06 (Update 2)
  - SPARC: Large page support for kernel heap
  - SPARC: sun4v 8k, 64k, 512k, 4M, **32M, 256M**
- Solaris 10 11/06 (Update 3)
  - SPARCV9 (OPL): 8k, 64k, 512k, 4M, **32M, 256M**
- Solaris 10 8/07 (Update 4)
  - SPARC: MPSS Extended to MAP\_SHARED anon mappings and non-ISM/DISM SysV Shared Segments. 8k and 4M defaults on all sun4u
  - SPARC: LPOOB changed to use smaller page sizes on N1
  - SPARC: Now have large page support for all types of user segments
  - AMD64: 4k, 2M
  - X64:

# Do I need Large Pages?

- Is the application memory intensive?
- How large is the address space?
- How much time is being wasted in MMU traps?
  - MMU traps are not visible with %usr/%sys
  - MMU traps are counted in the current context
  - e.g. User-bound process reports as %usr



# Trapstat Introduction

```
sol9# trapstat -t 1 111
cpu m| itlb-miss %tim itsb-miss %tim | dtlb-miss %tim dtsb-miss %tim | %tim
-----+-----+-----+-----+-----+-----+-----+-----+
 0 u|      1  0.0      0  0.0 | 2171237 45.7      0  0.0 | 45.7
 0 k|      2  0.0      0  0.0 |      3751  0.1      7  0.0 |  0.1
=====+=====+=====+=====+=====+=====+=====+=====
 ttl |      3  0.0      0  0.0 | 2192238 46.2      7  0.0 | 46.2
```

- This application *might* run almost 2x faster!

# Observing MMU traps

```
sol9# trapstat -T 1 111
cpu m size| itlb-miss %tim itsb-miss %tim | dtlb-miss %tim dtsb-miss %tim | %tim
-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+
 0 u 8k|      30 0.0          0 0.0 |      2170236 46.1          0 0.0 | 46.1
 0 u 64k|      0 0.0          0 0.0 |          0 0.0          0 0.0 | 0.0
 0 u 512k|      0 0.0          0 0.0 |          0 0.0          0 0.0 | 0.0
 0 u 4m|      0 0.0          0 0.0 |          0 0.0          0 0.0 | 0.0
-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+
 0 k 8k|      1 0.0          0 0.0 |      4174 0.1          10 0.0 | 0.1
 0 k 64k|      0 0.0          0 0.0 |          0 0.0          0 0.0 | 0.0
 0 k 512k|      0 0.0          0 0.0 |          0 0.0          0 0.0 | 0.0
 0 k 4m|      0 0.0          0 0.0 |          0 0.0          0 0.0 | 0.0
=====+=====+=====+=====+=====+=====+=====+=====+=====+=====+=====+=====+
    ttl |      31 0.0          0 0.0 |      2174410 46.2          10 0.0 | 46.2
sol9# trapstat -t 1 111
cpu m| itlb-miss %tim itsb-miss %tim | dtlb-miss %tim dtsb-miss %tim | %tim
-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+
 0 u|      1 0.0          0 0.0 |      2171237 45.7          0 0.0 | 45.7
 0 k|      2 0.0          0 0.0 |      3751 0.1           7 0.0 | 0.1
=====+=====+=====+=====+=====+=====+=====+=====+=====+=====+=====+=====+
    ttl |      3 0.0          0 0.0 |      2192238 46.2           7 0.0 | 46.2
```



# Setting Page Sizes

- Solution: `ppgsz(1)`, or `mpss.so.1`
  - Sets page size preference
  - Doesn't persist across `exec()`
  - Beginning with Solaris 10 1/06, Large Pages Out Of the Box (LPOOB) is enabled, so you don't need to do this...
    - You really want to be at Solaris 10 Update 4...

```
sol9# ppgsz -o heap=4M ./testprog
sol9# LD_PRELOAD=$LD_PRELOAD:mpss.so.1
sol9# export LD_PRELOAD=$LD_PRELOAD:mpss.so.1
sol9# export MPSSHEAP=4M
sol9# ./testprog
MPSSHEAP=size
MPSSSTACK=size
MPSSHEAP and MPSSSTACK specify the preferred page
sizes for the heap and stack, respectively. The speci-
fied page size(s) are applied to all created
processes.
MPSSCFGFILE=config-file
config-file is a text file which contains one or more
mpss configuration entries of the form:
exec-name:heap-size:stack-size
```

# Checking Allocated Page Sizes

```
Sol9# pmap -sx `pgrep testprog`
```

```
2953: ./testprog
```

Address	Kbytes	RSS	Anon	Locked	Pgsz	Mode	Mapped File
00010000	8	8	-	-	8K	r-x--	dev:277,83 ino:114875
00020000	8	8	8	-	8K	rwx--	dev:277,83 ino:114875
00022000	3960	3960	3960	-	8K	rwx--	[ heap ]
00400000	131072	131072	131072	-	4M	rwx--	[ heap ]
FF280000	120	120	-	-	8K	r-x--	<a href="#">libc.so.1</a>
FF340000	8	8	8	-	8K	rwx--	libc.so.1
FF390000	8	8	-	-	8K	r-x--	libc_psr.so.1
FF3A0000	8	8	-	-	8K	r-x--	libdl.so.1
FF3B0000	8	8	8	-	8K	rwx--	[ anon ]
FF3C0000	152	152	-	-	8K	r-x--	ld.so.1
FF3F6000	8	8	8	-	8K	rwx--	ld.so.1
FFBF0000	24	24	24	-	8K	rwx--	[ stack ]
<hr/>							
total Kb	135968	135944	135112	-			

# TLB traps eliminated

```
sol9# trapstat -T 1 111
cpu m size| itlb-miss %tim itsb-miss %tim | dtlb-miss %tim dtsb-miss %tim | %tim
-----+-----+-----+-----+-----+-----+-----+-----+-----+
 0 u 8k|      30  0.0      0  0.0 |      36  0.1      0  0.0 |  0.1
 0 u 64k|      0  0.0      0  0.0 |      0  0.0      0  0.0 |  0.0
 0 u 512k|      0  0.0      0  0.0 |      0  0.0      0  0.0 |  0.0
 0 u 4m|      0  0.0      0  0.0 |      0  0.0      0  0.0 |  0.0
-----+-----+-----+-----+-----+-----+-----+-----+-----+
 0 k 8k|      1  0.0      0  0.0 |  4174  0.1     10  0.0 |  0.1
 0 k 64k|      0  0.0      0  0.0 |      0  0.0      0  0.0 |  0.0
 0 k 512k|      0  0.0      0  0.0 |      0  0.0      0  0.0 |  0.0
 0 k 4m|      0  0.0      0  0.0 |      0  0.0      0  0.0 |  0.0
-----+-----+-----+-----+-----+-----+-----+-----+-----+
    ttl |      31  0.0      0  0.0 |  4200  0.2     10  0.0 |  0.2
```



USE



IMPROVE



EVANGELIZE

# Address Spaces: A Deeper Dive



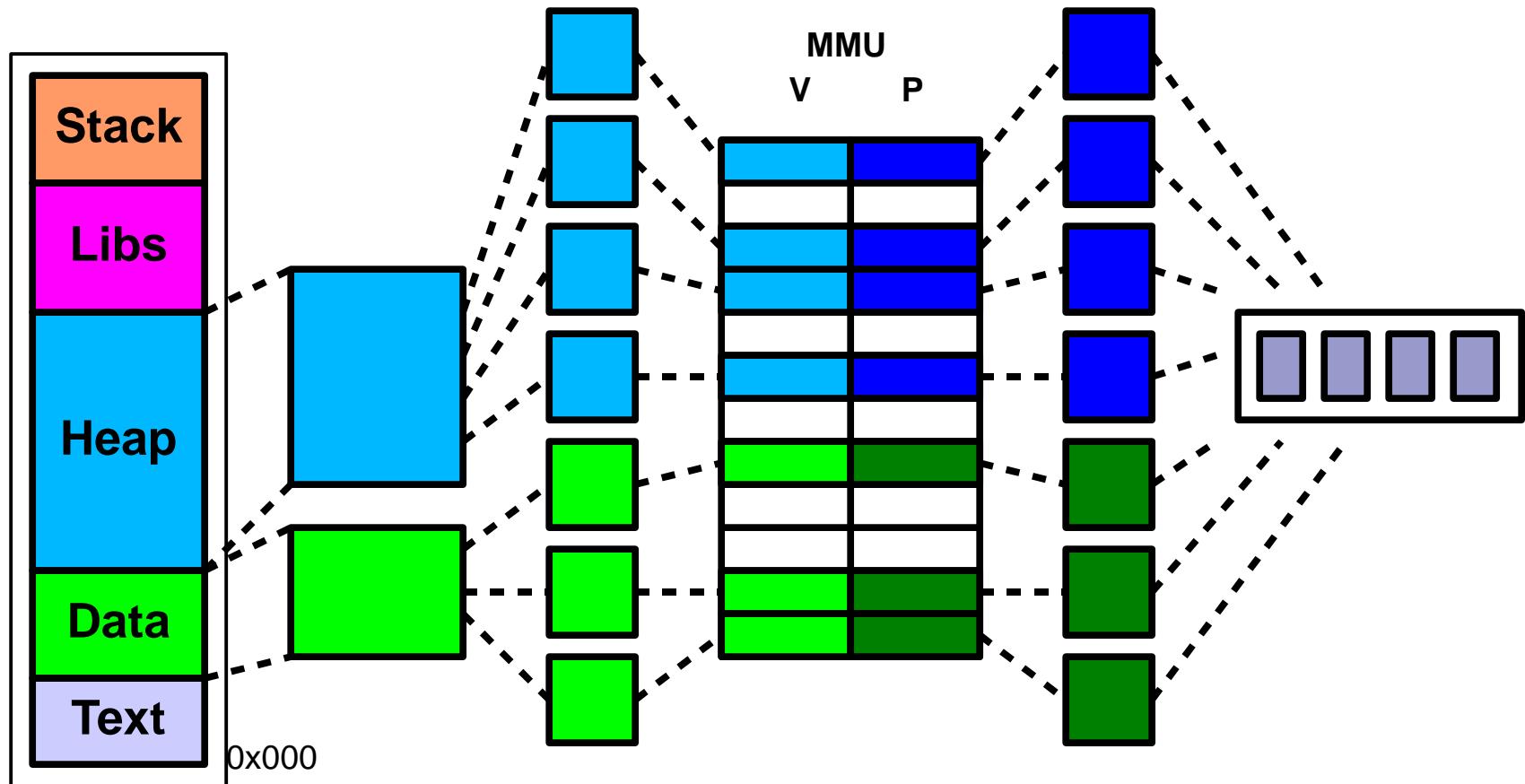
# Example Program

```
#include <sys/types.h>
const char * const_str = "My const string";
char * global_str = "My global string";
int    global_int = 42;
int
main(int argc, char * argv[])
{
    int local_int = 123;
    char * s;
    int i;
    char command[1024];

    global_int = 5;
    s = (char *)malloc(14000);
    s[0] = 'a';
    s[100] = 'b';
    s[8192] = 'c';

}
```

# Virtual to Physical





# Address Space

- Process Address Space
  - Process Text and Data
  - Stack (anon memory) and Libraries
  - Heap (anon memory)
- Kernel Address Space
  - Kernel Text and Data
  - Kernel Map Space (data structs, caches)
  - 32-bit Kernel map (64-bit Kernels only)
  - Trap table
  - Critical virtual memory data structures
  - Mapping File System Cache
    - ARC for ZFS mapped to kernel heap



# Address Space

0xFFBEC000

32-bit sun4u

0xFF3DC000

**Stack**

0xFFFFFFF.7FFFC000

**Libraries**

0xFFFFFFF.7F7F0000

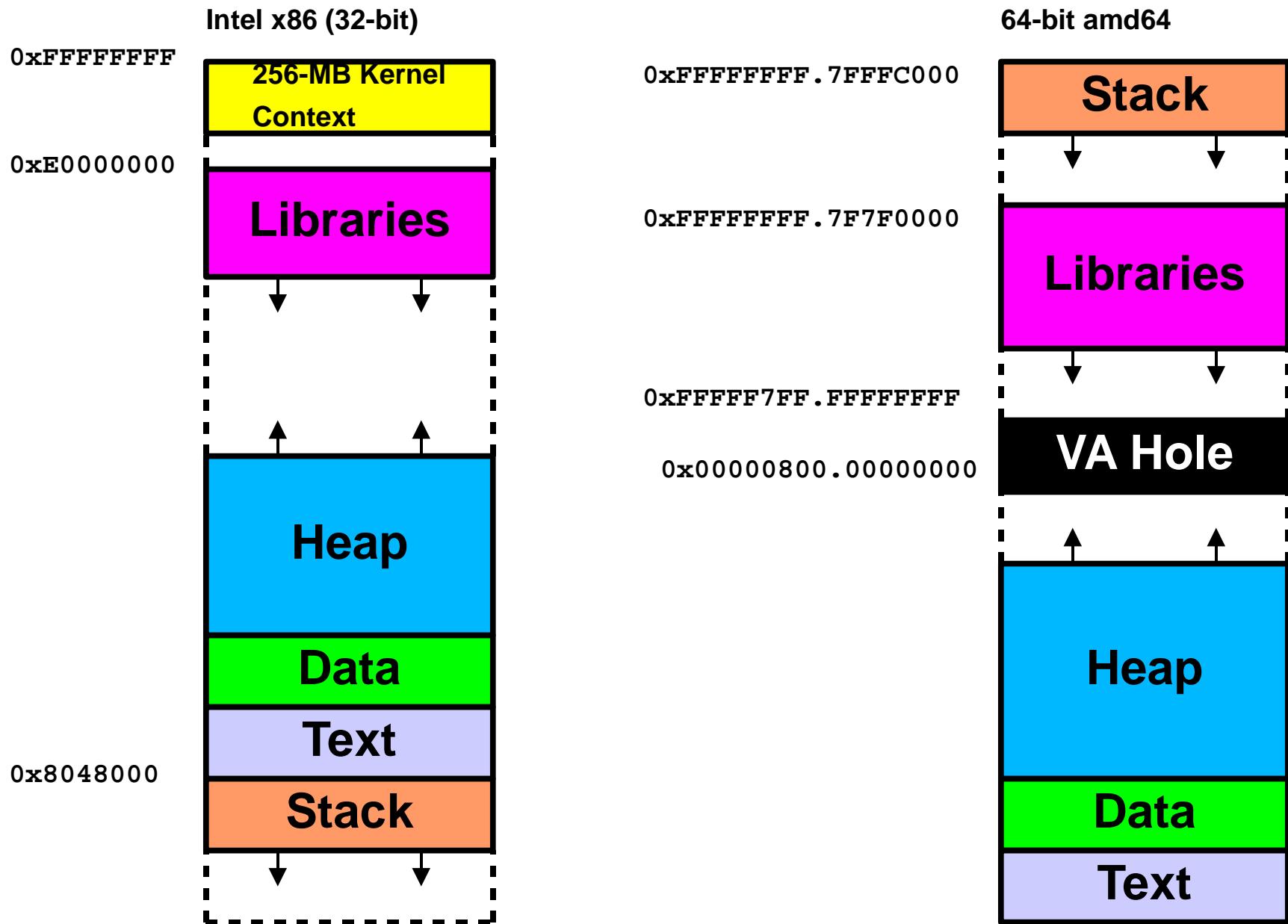
**Heap**

0xFFFF7FF.FFFFFFFF

**Data**

0x00000800.00000000

**Text****Stack****Libraries****VA Hole****Heap****Data****Text**





# pmap -x

```
Sol18# /usr/proc/bin/pmap -x $$
```

```
18084: csh
```

Address	Kbytes	Resident	Shared	Private	Permissions	Mapped File
00010000	144	144	136	8	read/exec	csh
00044000	16	16	-	16	read/write/exec	csh
00048000	120	104	-	104	read/write/exec	[ heap ]
FF200000	672	624	600	24	read/exec	libc.so.1
FF2B8000	24	24	-	24	read/write/exec	libc.so.1
FF2BE000	8	8	-	8	read/write/exec	libc.so.1
FF300000	16	16	8	8	read/exec	libc_psr.so.1
FF320000	8	8	-	8	read/exec	
libbmapmalloc.so.1						
FF332000	8	8	-	8	read/write/exec	
libbmapmalloc.so.1						
FF340000	8	8	-	8	read/write/exec	[ anon ]
FF350000	168	112	88	24	read/exec	libcurses.so.1
FF38A000	32	32	-	32	read/write/exec	libcurses.so.1
FF392000	8	8	-	8	read/write/exec	libcurses.so.1
FF3A0000	8	8	-	8	read/exec	libdl.so.1
FF3B0000	136	136	128	8	read/exec	ld.so.1
FF3E2000	8	8	-	8	read/write/exec	ld.so.1
FFBE6000	40	40	-	40	read/write/exec	[ stack ]
-----	-----	-----	-----	-----	-----	
total Kb	1424	1304	960	344		

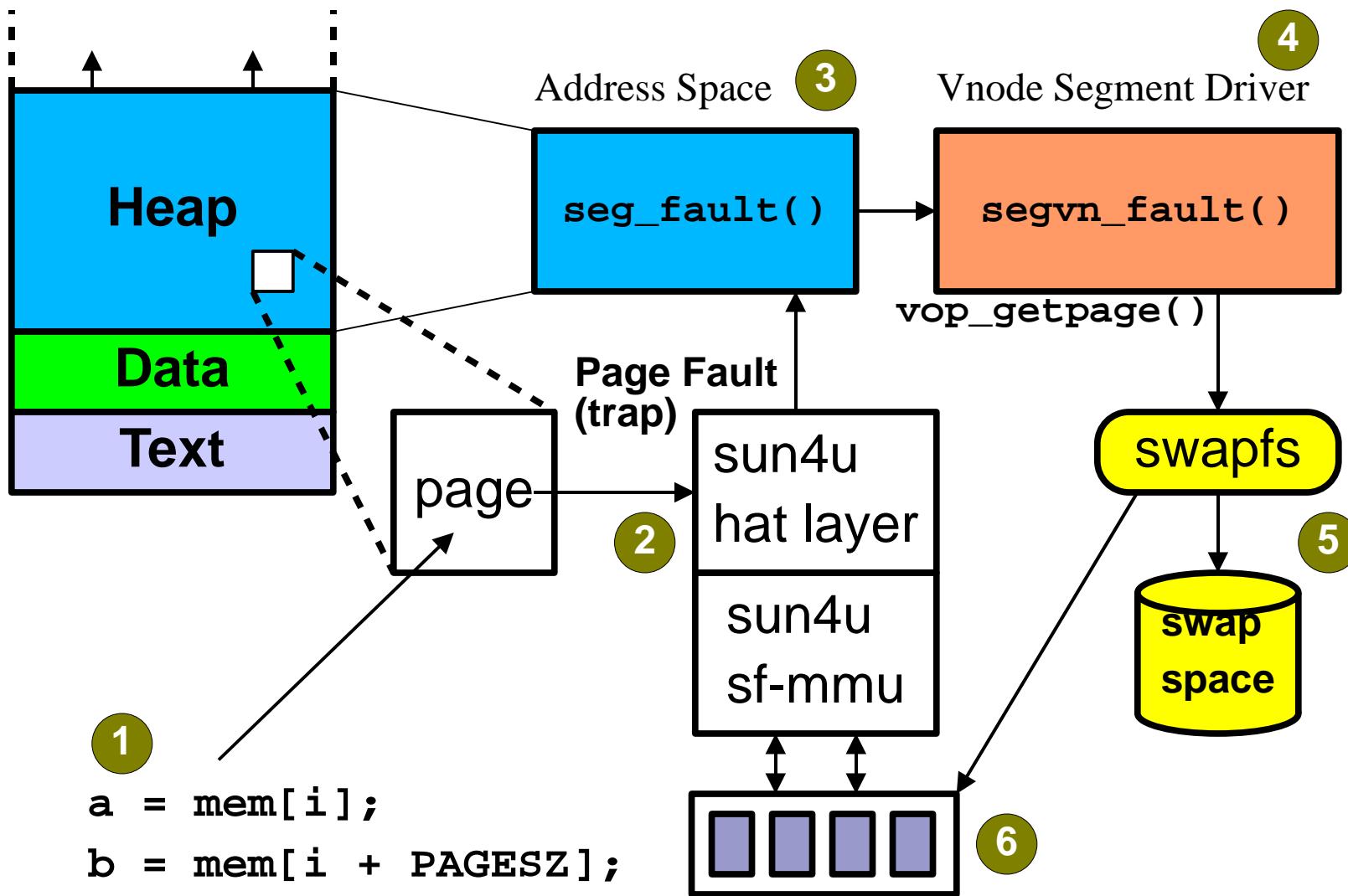
# Address Space Management

- Duplication; `fork()` -> `as_dup()`
- Destruction; `exit()`
- Creation of new segments
- Removal of segments
- Page protection (read, write, executable)
- Page Fault routing
- Page Locking
- Watchpoints

# Page Faults

- MMU-generated exception:
- Major Page Fault:
  - Failed access to VM location, in a segment
  - Page does not exist in physical memory
  - New page is created or copied from swap
  - If addr not in a valid segment (SIG-SEGV)
- Minor Page Fault:
  - Failed access to VM location, in a segment
  - Page is in memory, but no MMU translation
- Page Protection Fault:
  - An access that violates segment protection

## Page Fault Example:





# vmstat -p

swap = free and unreserved swap in KBytes

free = free memory measured in pages

re = kilobytes reclaimed from cache/free list

mf = minor faults - the page was in memory but was not mapped

fr = kilobytes that have been destroyed or freed

de = kilobytes freed after writes

sr = kilobytes scanned / second

executable pages: kilobytes in - out - freed

anonymous pages: kilobytes in - out  
- freed

file system pages:  
kilobytes in - out -  
freed

memory		page					executable			anonymous			filesystem		
swap	free	re	mf	fr	de	sr	epi	epo	epf	api	apo	apf	fpi	fpo	fpf
46715224	891296	24	350	0	0	0	0	0	0	4	0	0	27	0	0
46304792	897312	151	761	25	0	0	17	0	0	1	0	0	280	25	25
45886168	899808	118	339	1	0	0	3	0	0	1	0	0	641	1	1
46723376	899440	29	197	0	0	0	0	0	0	40	0	0	60	0	0



# Examining paging with dtrace VM Provider

- The dtrace VM provider provides a probe for each VM statistic
- We can observe all VM statistics via kstat:

```
$ kstat -n vm
module: cpu
name: vm
anonfree           instance: 0
anonpgin          class:    misc
anonpgout         0
as_fault          3180528
cow_fault         37280
crtime            463.343064
dfree              0
execfree           0
execpgin          442
execpgout         0
fsfree             0
fspgin             2103
fspgout            0
hat_fault          0
kernel_asflt      0
maj_fault          912
```



# Examining paging with dtrace

- Suppose one were to see the following output from vmstat(1M):

```
kthr memory page disk faults cpu
r b w swap    free   mf   pi   po   fr   de   sr   cd   s0s1   s2   in   sy   cs   us   sy   id
0 1 0 1341844 836720 26  311 1644  0   0   0   0 216  0   0   0 797  817  697  9   10  81
0 1 0 1341344 835300 238 934 1576  0   0   0   0 194  0   0   0 750  2795 791  7   14  79
0 1 0 1340764 833668 24   165 1149  0   0   0   0 133  0   0   0 637  813  547  5   4   91
0 1 0 1340420 833024 24   394 1002  0   0   0   0 130  0   0   0 621  2284 653  14  7   79
0 1 0 1340068 831520 14   202 380   0   0   0   0 59   0   0   0 482  5688 1434 25  7   68
```

- The pi column in the above output denotes the number of pages paged in. The vminfo provider makes it easy to learn more about the source of these page-ins:

```
dtrace -n pgin {@[execname] = count()}
dtrace: description "pgin" matched 1 probe
^C
xterm 1
ksh 1
ls 2
lpstat 7
sh 17
soffice 39
javaldx 103
soffice.bin 3065
```



# Examining paging with dtrace

- From the above, we can see that a process associated with the StarOffice Office Suite, soffice.bin, is responsible for most of the page-ins.
- To get a better picture of soffice.bin in terms of VM behavior, we may wish to enable all vminfo probes.
- In the following example, we run dtrace(1M) while launching StarOffice:

```
dtrace -P vminfo/execname == "soffice.bin"/{@[probename] = count()}
dtrace: description vminfo matched 42 probes
^C
pgout 16
anonfree 16
anonpgout 16
ppgout 16
dfree 16
execpgin 80
prot_fault 85
maj_fault 88
pgin 90
ppggin 90
cow_fault 859
zfod 1619
pgfrec 8811
pgrec 8827
as_fault 9495
```



# Examining paging with dtrace

- To further drill down on some of the VM behavior of StarOffice during startup, we could write the following D script:

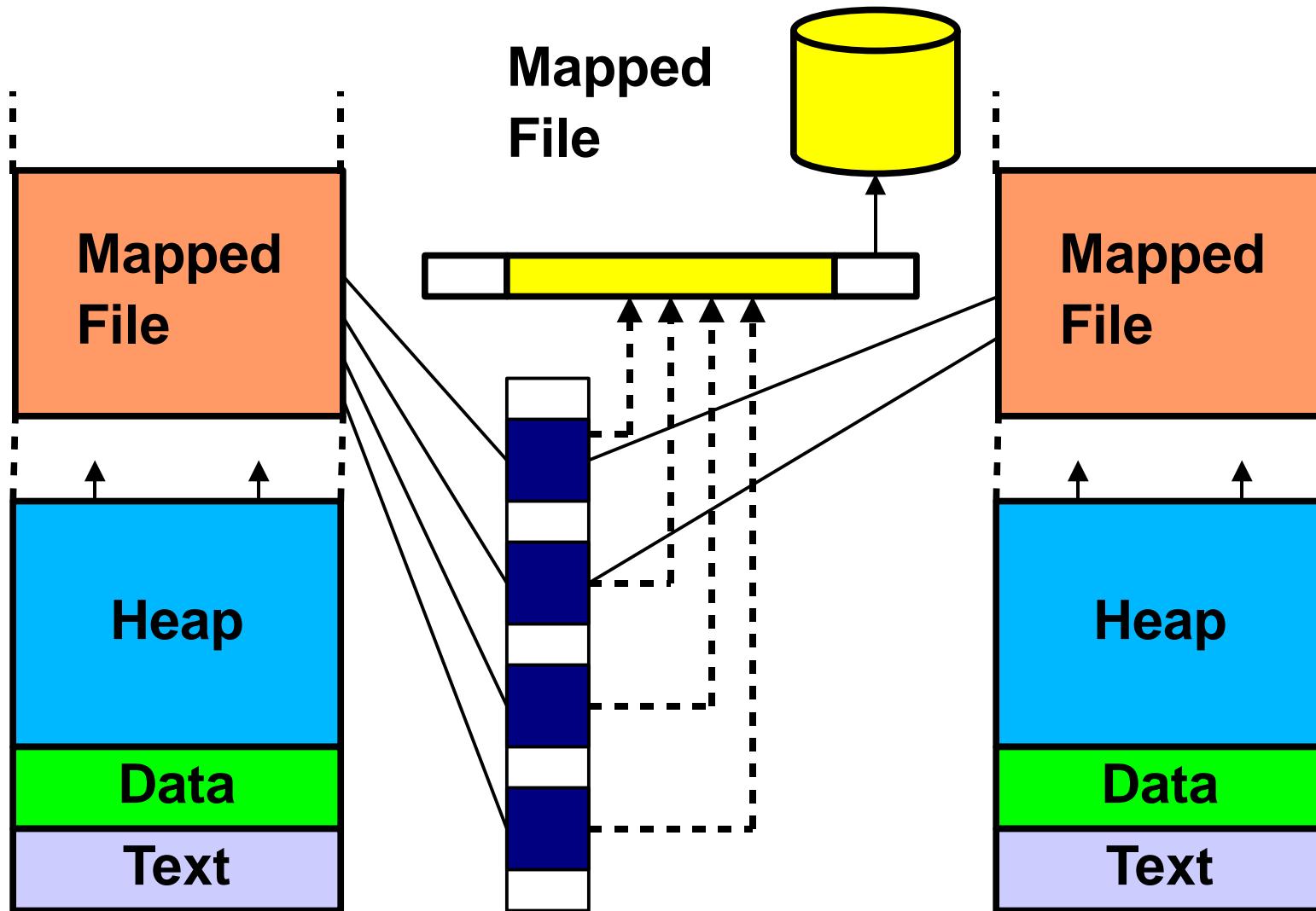
```
vminfo:::maj_fault, vminfo:::zfod, vminfo:::as_fault
/execname == "soffice.bin" && start == 0/
{
    /*
     * This is the first time that a vminfo probe has been hit; record
     * our initial timestamp.
    */
    start = timestamp;
}
vminfo:::maj_fault, vminfo:::zfod,vminfo:::as_fault
/execname == "soffice.bin"/
{
    /*
     * Aggregate on the probename, and lquantize() the number of seconds
     * since our initial timestamp. (There are 1,000,000,000 nanoseconds
     * in a second.) We assume that the script will be terminated before
     * 60 seconds elapses.
    */
    @[probename] = lquantize((timestamp - start) / 1000000000, 0, 60);
}
```

# Examining paging with dtrace

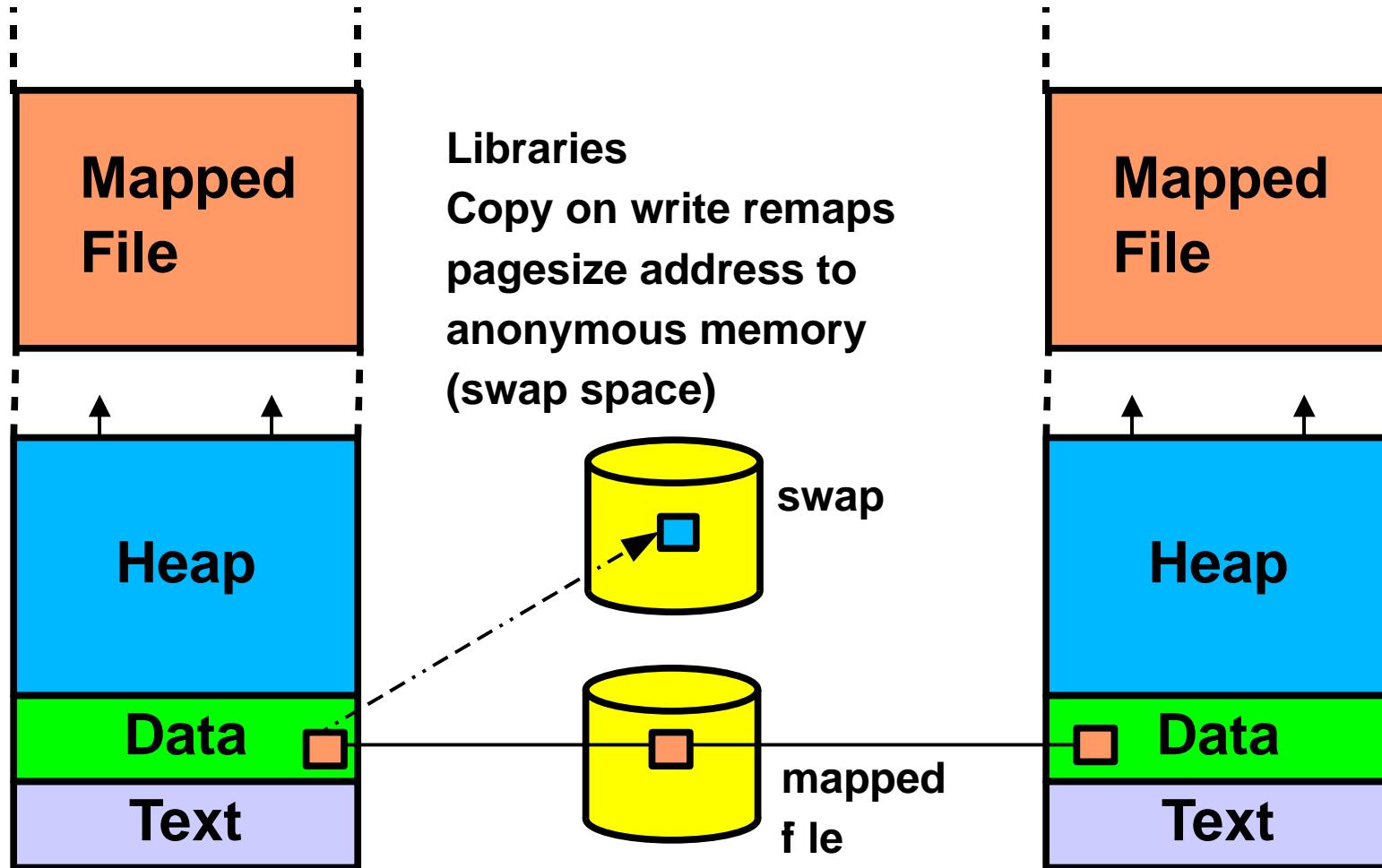
# Examining paging with dtrace

Zfod	value	Distribution	count
< 0			0
0	@@@@@@@		525
1	@@@@@@@@		605
2	@@		208
3	@@@		280
4			4
5			0
6			0
7			0
8			44
9	@@		161
10			2
11			0
12			0
13			4
14			0
15			29
16	@@@@@@@@@@@@@@@		1048
17			24
18			0
19			0
20			1
21			0
22			3
23			0

# Shared Mapped File



# Copy-on-write



# Anonymous Memory

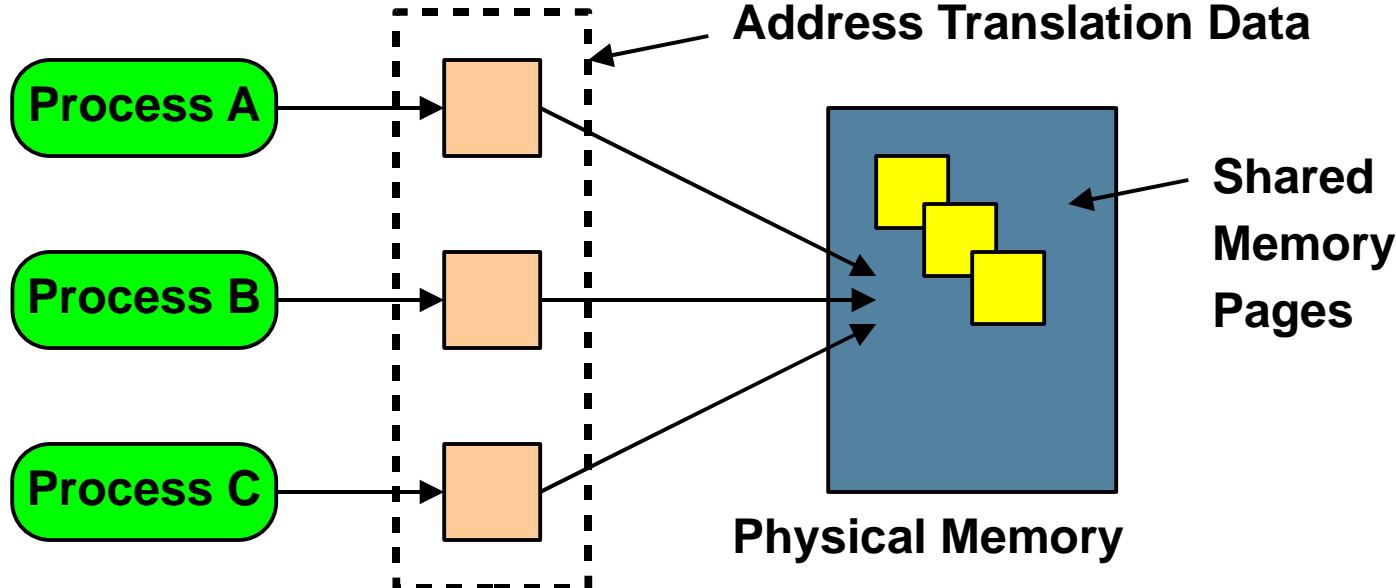
- Pages not "directly" backed by a vnode
- Heap, Stack and Copy-On-Write pages
- Pages are reserved when "requested"
- Pages allocated when "touched"
- Anon layer:
  - creates slot array for pages
  - Slots point to Anon structs
- Swapfs layer:
  - Pseudo file system for anon layer
  - Provides the backing store

# Intimate Shared Memory

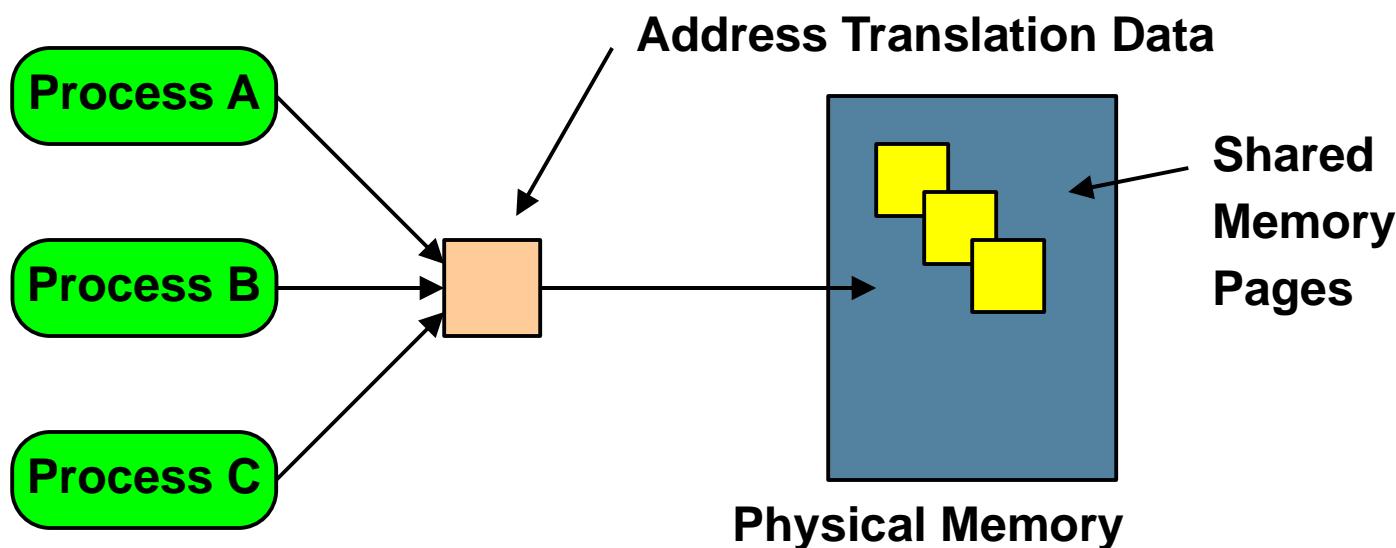
- System V shared memory (ipc) option
- Shared Memory optimization:
  - Additionally share low-level kernel data
  - Reduce redundant mapping info (V-to-P)
- Shared Memory is locked, never paged
  - No swap space is allocated
- Use **SHM\_SHARE\_MMU** flag in **shmat( )**

ISM

non-ISM



ISM





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# Session 3

## Processes, Threads, Scheduling Classes & The Dispatcher



# Process/Threads Glossary

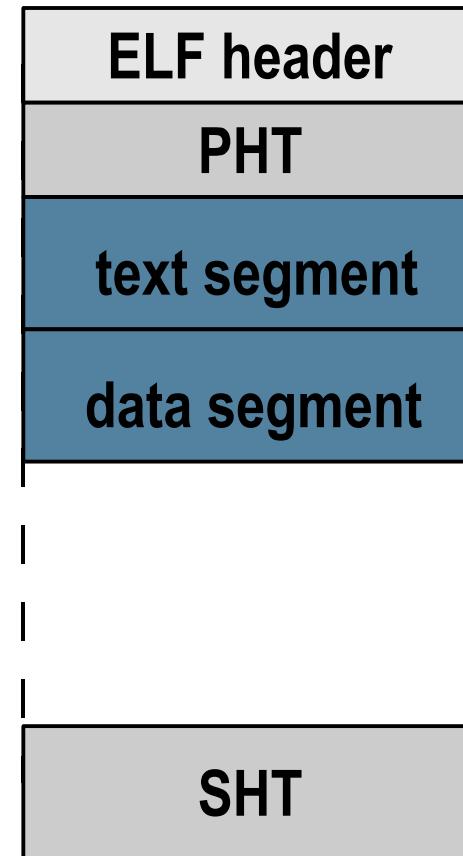
Process	The executable form of a program. An Operating System abstraction that encapsulates the execution context of a program
Thread	An executable entity
User Thread	A thread within the address space of a process
Kernel Thread	A thread in the address space of the kernel
Lightweight Process	LWP – An execution context for a kernel thread
Dispatcher	The kernel subsystem that manages queues of runnable kernel threads
Scheduling Class	Kernel classes that define the scheduling parameters (e.g. priorities) and algorithms used to multiplex threads onto processors
Dispatch Queues	Per-processor sets of queues of runnable threads (run queues)
Sleep Queues	Queues of sleeping threads
Turnstiles	A special implementation of sleep queues that provide priority inheritance.

# Executable Files

- Processes originate as executable programs that are exec'd
- Executable & Linking Format (ELF)
  - Standard executable binary file Application Binary Interface (ABI) format
  - Two standards components
    - Platform independent
    - Platform dependent (SPARC, x86)
  - Defines both the on-disk image format, and the in-memory image
  - ELF files components defined by
    - ELF header
    - Program Header Table (PHT)
    - Section Header Table (SHT)

# Executable & Linking Format (ELF)

- ELF header
  - Roadmap to the file
- PHT
  - Array of Elf\_Phdr structures, each defines a segment for the loader (exec)
- SHT
  - Array of Elf\_Shdr structures, each defines a section for the linker (ld)



# ELF Files

- ELF on-disk object created by the link-editor at the tail-end of the compilation process (although we still call it an a.out by default...)
- ELF objects can be statically linked or dynamically linked
  - Compiler "-B static" flag, default is dynamic
  - Statically linked objects have all references resolved and bound in the binary (libc.a)
  - Dynamically linked objects rely on the run-time linker, ld.so.1, to resolve references to shared objects at run time (libc.so.1)
  - Static linking is discouraged, and not possible for 64-bit binaries



# Examining ELF Files

- Use `elfdump(1)` to decompose ELF files

```
borntorun> elfdump -e /bin/ls
```

## ELF Header

<code>ei_magic:</code>	{ 0x7f, E, L, F }				
<code>ei_class:</code>	<code>ELFCLASS32</code>	<code>ei_data:</code>	<code>ELFDATA2MSB</code>		
<code>e_machine:</code>	<code>EM_SPARC</code>	<code>e_version:</code>	<code>EV_CURRENT</code>		
<code>e_type:</code>	<code>ET_EXEC</code>				
<code>e_flags:</code>	0				
<code>e_entry:</code>	0x10f00	<code>e_ehsize:</code>	52	<code>e_shstrndx:</code>	26
<code>e_shoff:</code>	0x4654	<code>e_shentsize:</code>	40	<code>e_shnum:</code>	27
<code>e_phoff:</code>	0x34	<code>e_phentsize:</code>	32	<code>e_phnum:</code>	6

```
borntorun>
```



# Examining ELF Files

- elfdump -c dumps section headers

```
borntorun> elfdump -c /bin/ls
Section Header[11]: sh_name: .text
  sh_addr: 0x10f00          sh_flags: [ SHF_ALLOC SHF_EXECINSTR ]
  sh_size: 0x2ec4           sh_type: [ SHT_PROGBITS ]
  sh_offset: 0xf00          sh_entsize: 0
  sh_link: 0                sh_info: 0
  sh_addralign: 0x8

Section Header[17]: sh_name: .got
  sh_addr: 0x24000          sh_flags: [ SHF_WRITE SHF_ALLOC ]
  sh_size: 0x4               sh_type: [ SHT_PROGBITS ]
  sh_offset: 0x4000          sh_entsize: 0x4
  sh_link: 0                sh_info: 0
  sh_addralign: 0x2000

Section Header[18]: sh_name: .plt
  sh_addr: 0x24004          sh_flags: [ SHF_WRITE SHF_ALLOC SHF_EXECINSTR ]
  sh_size: 0x28c             sh_type: [ SHT_PROGBITS ]
  sh_offset: 0x4004          sh_entsize: 0xc
  sh_link: 0                sh_info: 0
  sh_addralign: 0x4

Section Header[22]: sh_name: .data
  sh_addr: 0x24380          sh_flags: [ SHF_WRITE SHF_ALLOC ]
  sh_size: 0x154             sh_type: [ SHT_PROGBITS ]
  sh_offset: 0x4380          sh_entsize: 0
  sh_link: 0                sh_info: 0
  sh_addralign: 0x8
```



# Examining ELF Linker Dependencies

- Use `ldd(1)` to invoke the runtime linker (`ld.so`) on a binary file, and `pldd(1)` on a running process

```
borntorun> ldd netstat
 libdhcpagent.so.1 => /usr/lib/libdhcpagent.so.1
 libcmand.so.1 => /usr/lib/libcmand.so.1
 libsocket.so.1 => /usr/lib/libsocket.so.1
 libnsl.so.1 => /usr/lib/libnsl.so.1
 libkstat.so.1 => /usr/lib/libkstat.so.1
 libc.so.1 => /usr/lib/libc.so.1
 libdl.so.1 => /usr/lib/libdl.so.1
 libmp.so.2 => /usr/lib/libmp.so.2
 /usr/platform/SUNW,Ultra-60/lib/libc_psr.so.1
```

```
borntorun> pldd $$  
495:ksh  
/usr/lib/libsocket.so.1  
/usr/lib/libnsl.so.1  
/usr/lib/libc.so.1  
/usr/lib/libdl.so.1  
/usr/lib/libmp.so.2  
/usr/platform/sun4u/lib/libc_psr.so.1  
/usr/lib/locale/en_US.ISO8859-1/en_US.ISO8859-1.so.2  
borntorun>
```



# Runtime Linker Debug

```
solaris> LD_DEBUG=help date
00000:
...
00000: args      display input argument processing (ld only)
00000: audit     display runtime link-audit processing (ld.so.1 only)
00000: basic     provide basic trace information/warnings
00000: bindings  display symbol binding; detail flag shows absolute:relative
                  addresses (ld.so.1 only)
00000: cap       display hardware/software capability processing
00000: detail    provide more information in conjunction with other options
00000: demangle   display C++ symbol names in their demangled form
00000: entry     display entrance criteria descriptors (ld only)
00000: files     display input file processing (files and libraries)
00000: got       display GOT symbol information (ld only)
00000: help      display this help message
00000: libs      display library search paths; detail flag shows actual
                  library lookup (-l) processing
00000: long      display long object names without truncation
00000: map       display map file processing (ld only)
00000: move      display move section processing
00000: reloc     display relocation processing
00000: sections  display input section processing (ld only)
00000: segments  display available output segments and address/offset
                  processing; detail flag shows associated sections (ld only)
00000: statistics display processing statistics (ld only)
00000: strtab    display information about string table compression; detail
                  shows layout of string tables (ld only)
...
...
```

# Runtime Linker Debug - Libs

```
solaris> LD_DEBUG=libs /opt/filebench/bin/filebench
13686:
13686: hardware capabilities - 0x2b [ VIS V8PLUS DIV32 MUL32 ]
...
13686: find object=libc.so.1; searching
13686: search path=/lib (default)
13686: search path=/usr/lib (default)
13686: trying path=/lib/libc.so.1
13686: 1: calling .init (from sorted order): /lib/libc.so.1
13686: 1: calling .init (done): /lib/libc.so.1
13686: 1: transferring control: /opt/filebench/bin/filebench
13686: 1: trying path=/platform/SUNW,Ultra-Enterprise/lib/libc_psr.so.1
...
13686: find object=libm.so.2; searching
13686: search path=/usr/lib/lwp/sparcv9 (RPATH from file /opt/filebench/bin/sparcv9/
filebench)
13686: trying path=/usr/lib/lwp/sparcv9/libm.so.2
13686: search path=/lib/64 (default)
13686: search path=/usr/lib/64 (default)
13686: trying path=/lib/64/libm.so.2
13686:
13686: find object=libl.so.1; searching
13686: search path=/usr/lib/lwp/sparcv9 (RPATH from file /opt/filebench/bin/sparcv9/
filebench)
13686: trying path=/usr/lib/lwp/sparcv9/libl.so.1
13686: search path=/lib/64 (default)
13686: search path=/usr/lib/64 (default)
13686: trying path=/lib/64/libl.so.1
13686: trying path=/usr/lib/64/libl.so.1
```



# Runtime Linker Debug - Bindings

```
solaris> LD_DEBUG=bindings /opt/filebench/bin/filebench
15151:
15151: hardware capabilities - 0x2b [ VIS V8PLUS DIV32 MUL32 ]
15151: configuration file=/var/ld/ld.config: unable to process file
15151: binding file=/opt/filebench/bin/filebench to 0x0 (undefined weak): symbol
`__1cG_CrunMdo_exit_code6F_v_'
15151: binding file=/opt/filebench/bin/filebench to file=/lib/libc.so.1: symbol `__iob'
15151: binding file=/lib/libc.so.1 to 0x0 (undefined weak): symbol `__tnf_probe_notify'
15151: binding file=/lib/libc.so.1 to file=/opt/filebench/bin/filebench: symbol `__end'
15151: binding file=/lib/libc.so.1 to 0x0 (undefined weak): symbol `__ex_unwind'
15151: binding file=/lib/libc.so.1 to file=/lib/libc.so.1: symbol `__fnmatch_C'
15151: binding file=/lib/libc.so.1 to file=/lib/libc.so.1: symbol `__getdate_std'
...
15151: binding file=/opt/filebench/bin/sparcv9/filebench to file=/lib/64/libc.so.1: symbol
`__iob'
15151: binding file=/opt/filebench/bin/sparcv9/filebench to file=/lib/64/libc.so.1: symbol
`optarg'
15151: binding file=/lib/64/libm.so.2 to file=/opt/filebench/bin/sparcv9/filebench: symbol
`free'
15151: binding file=/lib/64/libm.so.2 to file=/lib/64/libm.so.2: symbol `__signgamf'
15151: binding file=/lib/64/libm.so.2 to file=/lib/64/libm.so.2: symbol `__signgaml'
15151: binding file=/lib/64/libm.so.2 to file=/lib/64/libm.so.2: symbol `__xpg6'
...
15151: 1: binding file=/lib/64/libc.so.1 to file=/lib/64/libc.so.1: symbol `__sigemptyset'
15151: 1: binding file=/lib/64/libc.so.1 to file=/lib/64/libc.so.1: symbol `__sigaction'
```



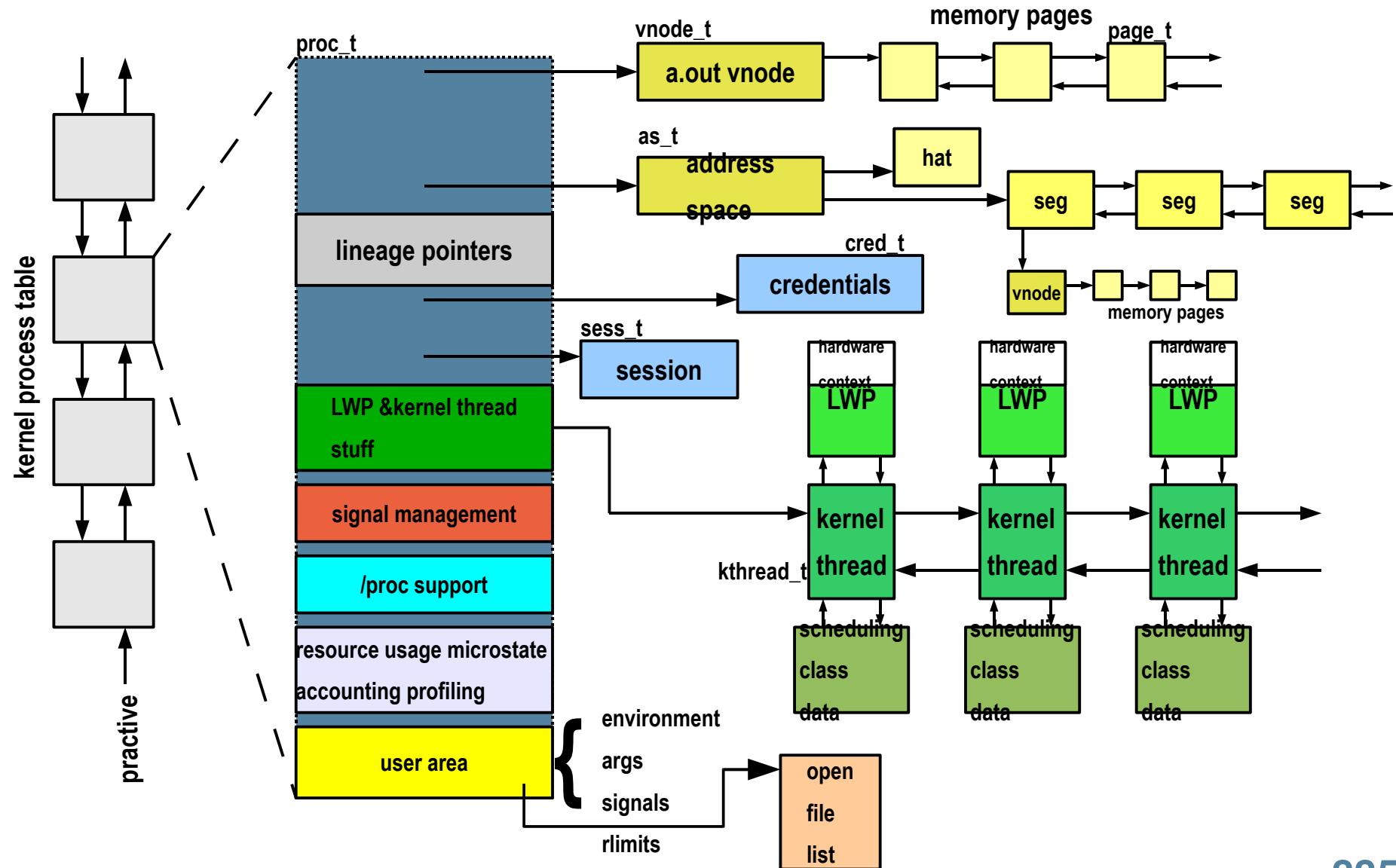
# Runtime Linker – Debug

- Explore the options in **The Linker and Libraries Guide**

# Solaris Process Model

- Solaris implements a multithreaded process model
  - Kernel threads are scheduled/executed
  - LWPs allow for each thread to execute system calls
  - Every kernel thread has an associated LWP
  - A non-threaded process has 1 kernel thread/LWP
  - A threaded process will have multiple kernel threads
  - All the threads in a process share all of the process context
    - Address space
    - Open files
    - Credentials
    - Signal dispositions
  - Each thread has its own stack

# Solaris Process





# Process Structure

```
# mdb -k
Loading modules: [ unix krtld genunix specfs dtrace ufs ip sctp usba fctl nca lofs nfs random
sppp crypto ptm logindmux cpc ]
> ::ps
S      PID      PPID      PGID      SID      UID      FLAGS          ADDR NAME
R        0        0        0        0      0x00000001 ffffffffbc1ce80 sched
R        3        0        0        0      0x00020001 ffffffff880838f8 fsflush
R        2        0        0        0      0x00020001 ffffffff88084520 pageout
R        1        0        0        0      0x42004000 ffffffff88085148 init
R    21344        1    21343    21280    2234 0x42004000 ffffffff95549938 tcpPerfServer
...
> ffffffff95549938::print proc_t
{
    p_exec = 0xffffffff9285dc40
    p_as = 0xffffffff87c776c8
    p_cred = 0xffffffff8fdeb448
    p_lwpcnt = 0x6
    p_zombcnt = 0
    p_tlist = 0xffffffff8826bc20
    .....
    u_ticks = 0x16c6f425
    u_comm = [ "tcpPerfServer" ]
    u_psargs = [ "/export/home/morgan/work/solaris_studio9/bin/tcpPerfServer 9551 9552" ]
    u_argc = 0x3
    u_argv = 0x8047380
    u_envp = 0x8047390
    u_cdir = 0xffffffff8bf3d7c0
    u_saved_rlimit = [
        {
            rlim_cur = 0xfffffffffffffd
            rlim_max = 0xfffffffffffffd
        }
    .....
        fi_nfiles = 0x3f
        fi_list = 0xffffffff8dc44000
        fi_rlist = 0
    }
    p_model = 0x100000
    p_rctls = 0xfffffffffa7cbb4c8
    p_dtrace_probes = 0
    p_dtrace_count = 0
    p_dtrace_helpers = 0
    p_zone = zone0
}
```



# Kernel Process Table

- Linked list of all processes (proc structures)
- kmem\_cache allocator dynamically allocates space needed for new proc structures
  - Up to v.v\_proc

```
borntorun> kstat -n var
module: unix           instance: 0
name:   var            class:    misc
        crttime          61.041156087
        snaptime          113918.894449089
        v_autooup         30
        v_buf              100
        v_bufhwm          20312
        [snip]
        v_maxsyspri       99
        v_maxup             15877
        v_maxupttl          15877
        v_nglobpris         110
        v_pbuf              0
        v_proc              15882
        v_sptmap             0
```

```
# mdb -k
Loading modules: [ unix krtld genunix ... ptm ipc ]
> max_nprocs/D
max_nprocs:
max_nprocs:      15882
>
```



# System-wide Process View - ps(1)

```

F S      UID   PID  PPID  C PRI NI      ADDR      SZ    WCHAN      STIME TTY      TIME CMD
0 S      root   824  386  0 40 20      ?     252      ? Sep 06 console 0:00 /usr/lib/saf/ttymon -g -h
-p mcdoug
0 S      root   823  386  0 40 20      ?     242      ? Sep 06 ?      0:00 /usr/lib/saf/sac -t 300
0 S      nobody 1718  716  0 40 20      ?     834      ? Sep 07 ?      0:35 /usr/apache/bin/httpd
0 S      root   591  374  0 40 20      ?     478      ? Sep 06 ?      0:00
/usr/lib/autofs/automountd
0 S      root   386  374  0 40 20      ?     262      ? Sep 06 ?      0:01 init
1 S      root   374  374  0 0 SY      ?     0      ? Sep 06 ?      0:00 zsched
0 S      daemon  490  374  0 40 20      ?     291      ? Sep 06 ?      0:00 /usr/sbin/rpcbind
0 S      daemon  435  374  0 40 20      ?     450      ? Sep 06 ?      0:00 /usr/lib/crypto/kcfd
0 S      root   603  374  0 40 20      ?     475      ? Sep 06 ?      0:12 /usr/sbin/nscd
0 S      root   580  374  0 40 20      ?     448      ? Sep 06 ?      0:02 /usr/sbin/syslogd
0 S      root   601  374  0 40 20      ?     313      ? Sep 06 ?      0:00 /usr/sbin/cron
0 S      daemon  548  374  0 40 20      ?     319      ? Sep 06 ?      0:00 /usr/lib/nfs/statd
0 S      daemon  550  374  0 40 20      ?     280      ? Sep 06 ?      0:00 /usr/lib/nfs/lockd
0 S      root   611  374  0 40 20      ?     329      ? Sep 06 ?      0:00 /usr/sbin/inetd -s
0 S      root   649  374  0 40 20      ?     152      ? Sep 06 ?      0:00 /usr/lib/utmpd
0 S      nobody  778  716  0 40 20      ?     835      ? Sep 06 ?      0:26 /usr/apache/bin/httpd
0 S      root   678  374  0 40 20      ?     612      ? Sep 06 ?      0:00 /usr/dt/bin/dtlogin
-daemon

```



# System-wide Process View - prstat(1)

PID	USERNAME	SIZE	RSS	STATE	PRI	NICE	TIME	CPU	PROCESS/NLWP
26292	root	5368K	3080K	run	24	0	0:00:00	1.5%	pkginstall/1
26188	rmc	4880K	4512K	cpu0	49	0	0:00:00	0.6%	prstat/1
202	root	3304K	1800K	sleep	59	0	0:00:07	0.3%	nscd/24
23078	root	20M	14M	sleep	59	0	0:00:56	0.2%	lupi_zones/1
23860	root	5104K	2328K	sleep	59	0	0:00:01	0.1%	sshd/1
...									
365	root	4760K	128K	sleep	59	0	0:00:00	0.0%	zoneadmd/4
364	root	4776K	128K	sleep	59	0	0:00:00	0.0%	zoneadmd/4
374	root	0K	0K	sleep	60	-	0:00:00	0.0%	zsched/1
361	root	2016K	8K	sleep	59	0	0:00:00	0.0%	ttymon/1
349	root	8600K	616K	sleep	59	0	0:00:20	0.0%	snmpd/1
386	root	2096K	360K	sleep	59	0	0:00:00	0.0%	init/1
345	root	3160K	496K	sleep	59	0	0:00:00	0.0%	sshd/1
591	root	3824K	184K	sleep	59	0	0:00:00	0.0%	automountd/2
....									
242	root	1896K	8K	sleep	59	0	0:00:00	0.0%	smcboot/1
248	smmsp	4736K	696K	sleep	59	0	0:00:08	0.0%	sendmail/1
245	root	1888K	0K	sleep	59	0	0:00:00	0.0%	smcboot/1
824	root	2016K	8K	sleep	59	0	0:00:00	0.0%	ttymon/1
204	root	2752K	536K	sleep	59	0	0:00:00	0.0%	inetd/1
220	root	1568K	8K	sleep	59	0	0:00:00	0.0%	powerd/3
313	root	2336K	216K	sleep	59	0	0:00:00	0.0%	snmpdx/1
184	root	4312K	872K	sleep	59	0	0:00:01	0.0%	syslogd/13
162	daemon	2240K	16K	sleep	60	-20	0:00:00	0.0%	lockd/2
Total:	126 processes, 311 lwps, load averages: 0.48, 0.48, 0.41								

# The Life Of A Process

- Process creation
  - fork(2) system call creates all processes
    - SIDL state
    - exec(2) overlays newly created process with executable image
- State Transitions
  - Typically runnable (SRUN), running (SONPROC) or sleeping (aka blocked, SSLEEP)
  - Maybe stopped (debugger) SSTOP
- Termination
  - SZOMB state
  - implicit or explicit exit(), signal (kill), fatal error



## Process Creation

- Traditional UNIX fork/exec model
  - fork(2) - replicate the entire process, including all threads
  - fork1(2) - replicate the process, only the calling thread
  - vfork(2) - replicate the process, but do not dup the address space
    - The new child borrows the parent's address space, until exec()

```
main(int argc, char *argv[])
{
    pid_t pid;
    pid = fork();
    if (pid == 0) /* in the child */
        exec();
    else if (pid > 0) /* in the parent */
        wait();
    else
        fork failed
}
```

# fork(2) in Solaris 10

- Solaris 10 unified the process model
  - libthread merged with libc
  - threaded and non-threaded processes look the same
- fork(2) now replicates only the calling thread
  - Previously, fork1(2) needed to be called to do this
  - Linking with -lpthread in previous releases also resulted in fork1(2) behaviour
- forkall(2) added for applications that require a fork to replicate all the threads in the process



# Process create example

## C code calling fork()

```
#include <sys/types.h>
#include <unistd.h>

int main(int argc, char *argv[])
{
    pid_t    ret, cpid, ppid;

    ppid = getpid();
    ret = fork();
    if (ret == -1) {
        perror("fork");
        exit(0);
    } else if (ret == 0) {
        printf("In child...\n");
    } else {
        printf("Child PID: %d\n",ret);
    }
    exit(0);
}
```

## D script to generate kernel trace

```
#!/usr/sbin/dtrace -Fs

syscall::fork1:entry
/ pid == $target /
{
    self->trace = 1;
}

fbt:::
/ self->trace /
{

syscall::fork1:return
/ pid == $target /
{
    self->trace = 0;
    exit(0);
}
```



# Fork Kernel Trace

```
CPU FUNCTION
0  -> fork1
0  <- fork1
0  -> cfork
0    -> secpolicy_basic_fork
0    <- secpolicy_basic_fork
0    -> priv_policy
0    <- priv_policy
0    -> holdlwps
0      -> schedctl_finish_sigblock
0      <- schedctl_finish_sigblock
0      -> pokelwps
0      <- pokelwps
0    <- holdlwps
0    -> flush_user_windows_to_stack
0    -> getproc
0      -> page_mem_avail
0      <- page_mem_avail
0      -> zone_status_get
0      <- zone_status_get
0      -> kmem_cache_alloc
0        -> kmem_cpu_reload
0        <- kmem_cpu_reload
0        <- kmem_cache_alloc
0      -> pid_assign
0        -> kmem_zalloc
0        <- kmem_cache_alloc
0        <- kmem_zalloc
0      -> pid_lookup
0      -> pid_getlocksolt
0    -> crgetruid
0    -> crgetzoneid
0    -> upcount_inc
0      -> rctl_set_dup
0        ...
0          -> project_cpu_shares_set
0          -> project_lwps_set
0          -> project_ntasks_set
0        ...
0      <- rctl_set_dup
```



# Fork Kernel Trace (cont)

```
0      -> as_dup
0          ...
0              <- hat_alloc
0      <- as_alloc
0          -> seg_alloc
0          -> rctl_set_fill_alloc_gp
0      <- rctl_set_dup_ready
0          -> rctl_set_dup
0              ...
0          -> forklwp
0              <- flush_user_windows_to_stack
0          -> save_syscall_args
0          -> lwp_create
0              <- thread_create
0                  -> lwp_stk_init
0                  -> kmem_zalloc
0              <- lwp_create
0          -> init_mstate
0          -> lwp_forkregs
0          -> forkctx
0          -> ts_alloc
0          -> ts_fork
0      <- forklwp
0          -> contract_process_fork
0          -> ts_forkret
0              -> continuelwps
0              -> ts_setrun
0              -> setbackdq
0              -> generic_enq_thread
0      <- ts_forkret
0          -> swtch
0              -> disp
0      <- swtch
0          -> resume
0              -> savectx
0              <- savectx
0                  -> restorectx
0              <- resume
0      <- cfork
0      <= fork1
```



# Watching Forks

## D script for watching fork(2)

```
#!/usr/sbin/dtrace -qs

syscall::forkall:entry
{
    @fall[execname] = count();
}
syscall::forkl:entry
{
    @f1[execname] = count();
}
syscall::vfork:entry
{
    @vf[execname] = count();
}

dtrace:::END
{
    printf("forkall\n");
    printa(@fall);
    printf("forkl\n");
    printa(@f1);
    printf("vfork\n");
    printa(@vf);
}
```

## Example run

```
# ./watchfork.d
^C
forkall
forkl
start-srvr      1
bash            3
4cli            6
vfork
```

# exec(2) – Load a new process image

- Most fork(2) calls are followed by an exec(2)
- exec – execute a new file
- exec overlays the process image with a new process constructed from the binary file passed as an arg to exec(2)
- The exec'd process inherits much of the caller's state:
  - nice value, scheduling class, priority, PID, PPID, GID, task ID, project ID, session membership, real UID & GID, current working directory, resource limits, processor binding, times, etc, ...

# Watching exec(2) with DTrace

- The D script...

```
#pragma D option quiet
proc:::exec
{
    self->parent = execname;
}
proc:::exec-success
/self->parent != NULL/
{
    @[self->parent, execname] = count();
    self->parent = NULL;
}
proc:::exec-failure
/self->parent != NULL/
{
    self->parent = NULL;
}
END
{
    printf("%-20s %-20s %s\n", "WHO", "WHAT", "COUNT");
    printa("%-20s %-20s %@d\n", @);
}
```



# Watching exec(2) with DTrace

- Example output:

```
# dtrace -s ./whoexec.d
^C

WHO        WHAT        COUNT
make.bin   yacc        1
tcsh       make        1
make.bin   spec2map   1
sh         grep        1
lint       lint2       1
sh         lint        1
sh         ln          1
cc         ld          1
make.bin   cc          1
lint       lint1      1
```

# DTrace Toolkit - execsnoop

In this example the command "man gzip" was executed. The output lets us see what the man command is actually doing,

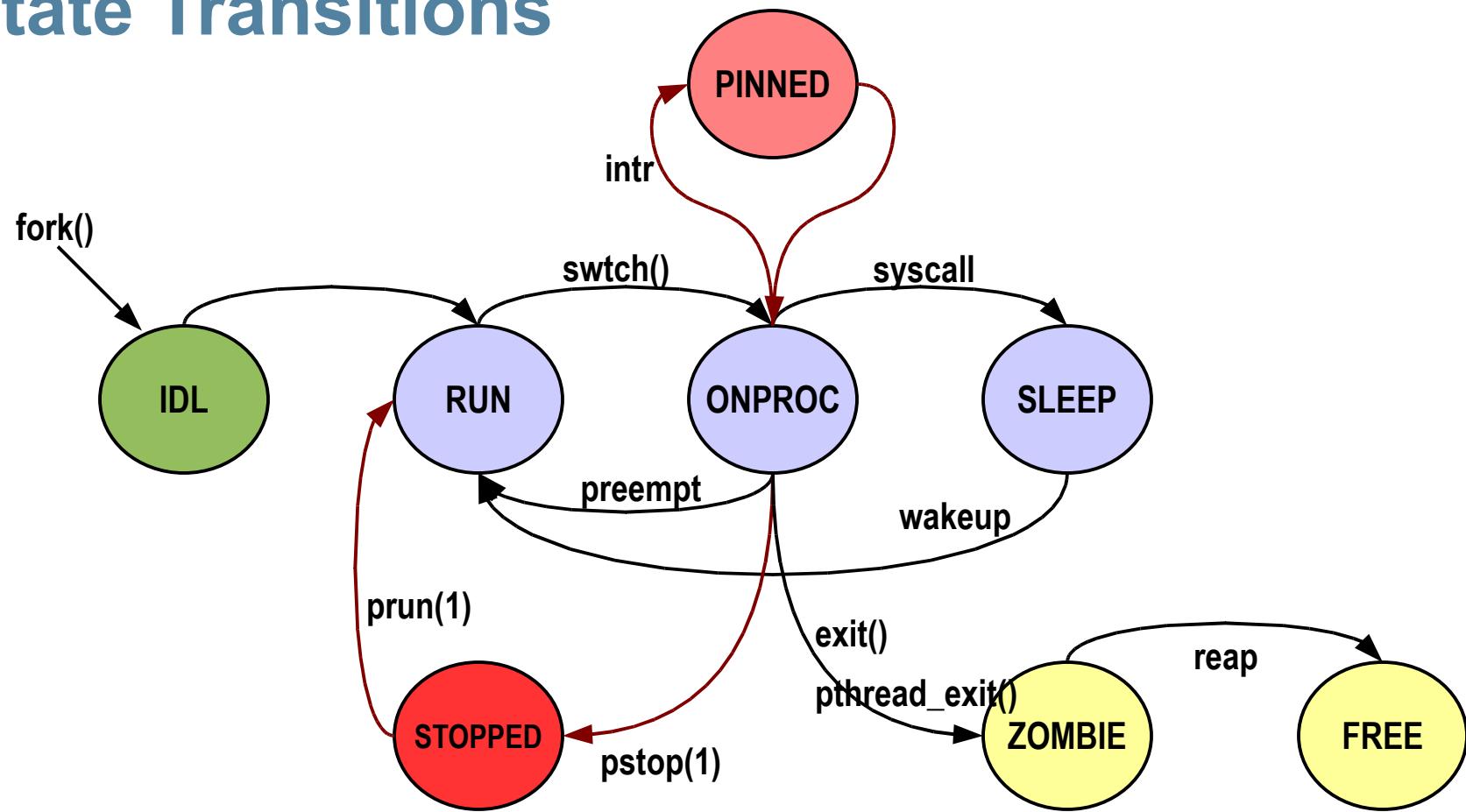
```
# ./execsnoop  
UID PID PPID ARGS  
100 3064 2656 man gzip  
100 3065 3064 sh -c cd /usr/share/man; tbl /usr/share/man/man1/gzip.1 |nroff -u0 -Tlp -man -  
100 3067 3066 tbl /usr/share/man/man1/gzip.1  
100 3068 3066 nroff -u0 -Tlp -man -  
100 3066 3065 col -x  
100 3069 3064 sh -c trap " 1 15; /usr/bin/mv -f /tmp/mpoMaa_f /usr/share/man/cat1/gzip.1 2>  
100 3070 3069 /usr/bin/mv -f /tmp/mpoMaa_f /usr/share/man/cat1/gzip.1  
100 3071 3064 sh -c more -s /tmp/mpoMaa_f  
100 3072 3071 more -s /tmp/mpoMaa_f  
^C
```

# Process / Thread States

- It's really kernel threads that change state
- Kernel thread creation is not flagged as a distinct state
  - Initial state is `TS_RUN`
- Kernel threads are `TS_FREE` when the process, or LWP/kthread, terminates

Process State	Kernel Thread State
SIDL	
SRUN	<code>TS_RUN</code>
SONPROC	<code>TS_ONPROC</code>
SSLEEP	<code>TS_SLEEP</code>
SSTOP	<code>TS_STOPPED</code>
SZOMB	<code>TS_ZOMB</code>
	<code>TS_FREE</code>

# State Transitions





# Watching Process States

PID	USERNAME	SIZE	RSS	STATE	PRI	NICE	TIME	CPU	PROCESS/NLWP
27946	root	4880K	4520K	cpu0	59	0	0:00:00	0.7%	prstat/1
28010	root	4928K	2584K	run	29	0	0:00:00	0.7%	pkginstall/1
23078	root	20M	14M	sleep	59	0	0:00:57	0.3%	lupi_zones/1
25947	root	5160K	2976K	sleep	59	0	0:00:04	0.3%	sshd/1
24866	root	5136K	2136K	sleep	59	0	0:00:01	0.2%	sshd/1
202	root	3304K	1800K	sleep	59	0	0:00:09	0.2%	nscd/24
23001	root	5136K	2176K	sleep	59	0	0:00:04	0.1%	sshd/1
23860	root	5248K	2392K	sleep	59	0	0:00:05	0.1%	sshd/1
25946	rmc	3008K	2184K	sleep	59	0	0:00:02	0.1%	ssh/1
25690	root	1240K	928K	sleep	59	0	0:00:00	0.1%	sh/1
...									
312	root	4912K	24K	sleep	59	0	0:00:00	0.0%	dtlogin/1
250	root	4760K	696K	sleep	59	0	0:00:16	0.0%	sendmail/1
246	root	1888K	0K	sleep	59	0	0:00:00	0.0%	smcboot/1
823	root	1936K	224K	sleep	59	0	0:00:00	0.0%	sac/1
242	root	1896K	8K	sleep	59	0	0:00:00	0.0%	smcboot/1
248	smmmsp	4736K	680K	sleep	59	0	0:00:08	0.0%	sendmail/1
245	root	1888K	0K	sleep	59	0	0:00:00	0.0%	smcboot/1
824	root	2016K	8K	sleep	59	0	0:00:00	0.0%	ttymon/1
204	root	2752K	520K	sleep	59	0	0:00:00	0.0%	inetd/1
220	root	1568K	8K	sleep	59	0	0:00:00	0.0%	powerd/3
313	root	2336K	216K	sleep	59	0	0:00:00	0.0%	snmpd/1

Total: 127 processes, 312 lwps, load averages: 0.62, 0.62, 0.53

# Microstates

- Fine-grained state tracking for processes/threads
  - Off by default in Solaris 8 and Solaris 9
  - On by default in Solaris 10
- Can be enabled per-process via /proc
- prstat -m reports microstates
  - As a percentage of time for the sampling period
    - USR – user mode
    - SYS - kernel mode
    - TRP – trap handling
    - TFL – text page faults
    - DFL – data page faults
    - LCK – user lock wait
    - SLP - sleep
    - LAT – waiting for a processor (sitting on a run queue)



# prstat – process microstates

```
# prstat -m
  PID USERNAME USR SYS TRP TFL DFL LCK SLP LAT VCX ICX SCL SIG PROCESS/NLWP
  739 root     0.3 0.3 0.0 0.0 0.0 0.0 99  0.0 126  3 345  5 Xsun/1
15611 root     0.1 0.3 0.0 0.0 0.0 0.0 100  0.0 23   0 381  0 prstat/1
1125 tlc      0.3 0.0 0.0 0.0 0.0 0.0 100  0.0 29   0 116  0 gnome-panel/1
15553 rmc     0.1 0.2 0.0 0.0 0.0 0.0 100  0.0 24   0 381  0 prstat/1
5591 tlc      0.1 0.0 0.0 0.0 0.0 33   66  0.0 206  0 1K   0 mozilla-bin/6
1121 tlc      0.0 0.0 0.0 0.0 0.0 0.0 100  0.1 50   0 230  0 metacity/1
2107 rmc     0.0 0.0 0.0 0.0 0.0 0.0 100  0.0 25   0 36   0 gnome-terminal/1
  478 root     0.0 0.0 0.0 0.0 0.0 0.0 100  0.0 17   0 14   0 squid/1
  798 root     0.0 0.0 0.0 0.0 0.0 0.0 100  0.0 11   0 23   0 Xsun/1
1145 tlc      0.0 0.0 0.0 0.0 0.0 0.0 100  0.0 25   1 34   0 mixer_applet/1
1141 rmc     0.0 0.0 0.0 0.0 0.0 0.0 100  0.0 25   0 32   0 mixer_applet/1
1119 tlc      0.0 0.0 0.0 0.0 0.0 0.0 100  0.0 5    0 40   0 gnome-smproxy/1
1127 tlc      0.0 0.0 0.0 0.0 0.0 0.0 100  0.0 7    0 29   0 nautilus/3
1105 rmc     0.0 0.0 0.0 0.0 0.0 0.0 100  0.0 7    0 27   0 nautilus/3
  713 root     0.0 0.0 0.0 0.0 0.0 85   15  0.0 2    0 100  0 mibiisa/7
  174 root     0.0 0.0 0.0 0.0 0.0 0.0 100  0.0 5    0 50   5 ipmon/1
 1055 tlc      0.0 0.0 0.0 0.0 0.0 0.0 100  0.0 5    0 30   0 dsdm/1
Total: 163 processes, 275 lwps, load averages: 0.07, 0.07, 0.07
```



# prstat – user summary

```
# prstat -t
NPROC USERNAME   SIZE   RSS  MEMORY    TIME   CPU
 128 root        446M  333M  1.4%  47:14:23  11%
    2 measter     6600K 5016K  0.0%  0:00:07  0.2%
    1 clamb       9152K 8344K  0.0%  0:02:14  0.1%
    2 rmc         7192K 6440K  0.0%  0:00:00  0.1%
    1 bricker     5776K 4952K  0.0%  0:00:20  0.1%
    2 asd          10M   8696K  0.0%  0:00:01  0.1%
    1 fredz        7760K 6944K  0.0%  0:00:05  0.1%
    2 jenks        8576K 6904K  0.0%  0:00:01  0.1%
    1 muffin       15M   14M   0.1%  0:01:26  0.1%
    1 dte          3800K 3016K  0.0%  0:00:04  0.0%
    2 adjg         8672K 7040K  0.0%  0:00:03  0.0%
    3 msw          14M   10M   0.0%  0:00:00  0.0%
    1 welza        4032K 3248K  0.0%  0:00:29  0.0%
    2 kimc         7848K 6344K  0.0%  0:00:25  0.0%
    4 jcmartin    13M   9904K  0.0%  0:00:03  0.0%
    1 rascal        17M   16M   0.1%  0:02:11  0.0%
    1 rab           3288K 2632K  0.0%  0:02:11  0.0%
    1 gjmurphy     3232K 2392K  0.0%  0:00:00  0.0%
    1 Ktheisen      15M   14M   0.1%  0:01:16  0.0%
    1 nagendra     3232K 2400K  0.0%  0:00:00  0.0%
    2 ayong        8320K 6832K  0.0%  0:00:02  0.0%
Total: 711 processes, 902 lwps, load averages: 3.84, 4.30, 4.37
```



# Solaris 9 / 10 ptools

```
/usr/bin/pflags [-r] [pid | core] ...
/usr/bin/pcred [pid | core] ...
/usr/bin/pldd [-F] [pid | core] ...
/usr/bin/psig [-n] pid...
/usr/bin/pstack [-F] [pid | core] ...
/usr/bin/pfiles [-F] pid...
/usr/bin/pwdx [-F] pid...
/usr/bin/pstop pid...
/usr/bin/prun pid...
/usr/bin/pwait [-v] pid...
/usr/bin/ptree [-a] [pid | user] ...
/usr/bin/ptime command [arg...]
/usr/bin/pmap -[xs] [-rs] [pid | core] ...
/usr/bin/pgrep [-flvx] [-n] [-o] [-d delim] [-P ppidlist] [-g pgrp[ist]] [-s sidlist] [-u euidlist] [-U uidlist] [-G gidlist] [-J projidlist] [-t term[list]] [-T taskidlist] [pattern]
/usr/bin/pkill [-signal] [-fvx] [-n | -o] [-P ppidlist] [-g pgrp[ist]] [-s sidlist] [-u euidlist] [-U uidlist] [-G gidlist] [-J projidlist] [-t term[list]] [-T taskidlist] [pattern]
/usr/bin/plimit [-km] pid...
{-cdfnstv} soft,hard... pid...
/usr/bin/ppgsz [-F] -o option[,option] cmd | -p pid...
/usr/bin/prctl [-t [basic | privileged | system]] [-e | -d action]
[-rx] [-n name [-v value]] [-i idtype] [id...]
/usr/bin/preap [-F] pid
/usr/bin/pargs [-aceFx] [pid | core] ...
```

# Tracing

- Trace user signals and system calls - truss
  - Traces by stopping and starting the process
  - Can trace system calls, inline or as a summary
  - Can also trace shared libraries and a.out
- Linker/library interposing/profiling/tracing
  - LD\_ environment variables enable link debugging
  - man ld.so.1
  - using the LD\_PRELOAD env variable
- Trace Normal Formal (TNF)
  - Kernel and Process Tracing
  - Lock Tracing
- Kernel Tracing
  - lockstat, tnf, kgmon

# Process Tracing – Truss



# Process Tracing – System Call Summary

- Counts total cpu seconds per system call and calls

```
# truss -c dd if=500m of=/dev/null bs=16k count=2k
syscall      seconds    calls  errors
_exit        .00        1
read          .34      2048
write         .03      2056
open          .00        4
close          .00        6
brk           .00        2
fstat          .00        3
execve         .00        1
sigaction       .00        2
mmap           .00        7
munmap          .00        2
sysconfig       .00        1
llseek          .00        1
creat64         .00        1
open64          .00        1
-----
sys totals:   .37      4136      0
usr time:     .00
elapsed:      .89
```



# Library Tracing - truss -u

```
# truss -d -u a.out,libc dd if=500m of=/dev/null bs=16k count=2k
Base time stamp: 925932005.2498 [Wed May 5 12:20:05 PDT 1999]
0.0000 execve("/usr/bin/dd", 0xFFBEF68C, 0xFFBEF6A4) argc = 5
0.0073 open("/dev/zero", O_RDONLY)      = 3
0.0077 mmap(0x00000000, 8192, PROT_READ|PROT_WRITE|PROT_EXEC, MAP_PRIVATE, 3, 0) = 0xFF3A0000
0.0094 open("/usr/lib/libc.so.1", O_RDONLY)      = 4
0.0097 fstat(4, 0xFFBEF224)          = 0
0.0100 mmap(0x00000000, 8192, PROT_READ|PROT_EXEC, MAP_PRIVATE, 4, 0) = 0xFF390000
0.0102 mmap(0x00000000, 761856, PROT_READ|PROT_EXEC, MAP_PRIVATE, 4, 0) = 0xFF280000
0.0105 munmap(0xFF324000, 57344)        = 0
0.0107 mmap(0xFF332000, 25284, PROT_READ|PROT_WRITE|PROT_EXEC, MAP_PRIVATE|MAP_FIXED, 4, 663552) = 0xFF332000
0.0113 close(4)                      = 0
0.0116 open("/usr/lib/libdl.so.1", O_RDONLY)      = 4
0.0119 fstat(4, 0xFFBEF224)          = 0
0.0121 mmap(0xFF390000, 8192, PROT_READ|PROT_EXEC, MAP_PRIVATE|MAP_FIXED, 4, 0) = 0xFF390000
0.0124 close(4)                      = 0
0.0127 open("/usr/platform/SUNW,Ultra-2/lib/libc_psr.so.1", O_RDONLY) = 4
0.0131 fstat(4, 0xFFBEF004)          = 0
0.0133 mmap(0x00000000, 8192, PROT_READ|PROT_EXEC, MAP_PRIVATE, 4, 0) = 0xFF380000
0.0135 mmap(0x00000000, 16384, PROT_READ|PROT_EXEC, MAP_PRIVATE, 4, 0) = 0xFF370000
0.0138 close(4)                      = 0
0.2369 close(3)                      = 0
0.2372 munmap(0xFF380000, 8192)        = 0
0.2380 -> libc:atexit(0xff3b9e8c, 0x23400, 0x0, 0x0)
0.2398 <- libc:atexit() = 0
0.2403 -> libc:atexit(0x12ed4, 0xff3b9e8c, 0xff334518, 0xff332018)
0.2419 <- libc:atexit() = 0
0.2424 -> __init(0x0, 0x12ed4, 0xff334518, 0xff332018)
0.2431 <- __init() = 0
0.2436 -> main(0x5, 0xffbef68c, 0xffbef6a4, 0x23400)
0.2443 -> libc:setlocale(0x6, 0x12f14, 0x0, 0x0)
0.2585 <- libc:setlocale() = 0xff31f316
```



# Library Tracing – aptrace(1)

```
sunsys> aptrace ls
ls      -> libc.so.1:atexit(func = 0xff3caa24) = 0x0
ls      -> libc.so.1:atexit(func = 0x13ad4) = 0x0
ls      -> libc.so.1:setlocale(category = 0x6, locale = "") = "/en_US.ISO8859-1/en_"
ls      -> libc.so.1:textdomain(domainname = "SUNW_OST_OSCMD") = "SUNW_OST_OSCMD"
ls      -> libc.so.1:time(tloc = 0x0) = 0x3aee2678
ls      -> libc.so.1:isatty(fildes = 0x1) = 0x1
ls      -> libc.so.1:getopt(argc = 0x1, argv = 0xffbeeff4, optstring =
                           "RaAdC1xmnlogrtucpFbq") = 0xffffffff errno = 0 (Error 0)
ls      -> libc.so.1:getenv(name = "COLUMNS") = "<nil>"
ls      -> libc.so.1:ioctl(0x1, 0x5468, 0x2472a)
ls      -> libc.so.1:malloc(size = 0x100) = 0x25d10
ls      -> libc.so.1:malloc(size = 0x9000) = 0x25e18
ls      -> libc.so.1:lstat64(path = ".", buf = 0xffbeee98) = 0x0
ls      -> libc.so.1:qsort(base = 0x25d10, nel = 0x1, width = 0x4, compar = 0x134bc)
ls      -> libc.so.1:.div(0x50, 0x3, 0x50)
ls      -> libc.so.1:.div(0xffffffff, 0x1a, 0x0)
ls      -> libc.so.1:.mul(0x1, 0x0, 0xffffffff)
ls      -> libc.so.1:.mul(0x1, 0x1, 0x0)
```

# User Threads

- The programming abstraction for creating multithreaded programs
  - Parallelism
  - POSIX and UI thread APIs
    - `thr_create(3THR)`
    - `pthread_create(3THR)`
  - Synchronization
    - Mutex locks, reader/writer locks, semaphores, condition variables
- Solaris 2 originally implemented an MxN threads model (T1)
  - “unbound” threads
- Solaris 8 introduced the 1 level model (T2)
  - `/usr/lib/lwp/libthread.so`
- T2 is the default in Solaris 9 and Solaris 10



# Threads Primer Example:

```
#include <pthread.h>
#include <stdio.h>
mutex_t mem_lock;
void childthread(void *argument)
{
    int i;
    for(i = 1; i <= 100; ++i) {
        print("Child Count - %d\n", i);
    }
    pthread_exit(0);
}
int main(void)
{
    pthread_t thread, thread2;
    int ret;

    if ((pthread_create(&thread, NULL, (void *)childthread, NULL)) < 0) {
        printf ("Thread Creation Failed\n");
        return (1);
    }
    pthread_join(thread,NULL);
    print("Parent is continuing....\n");
    return (0);
}
```



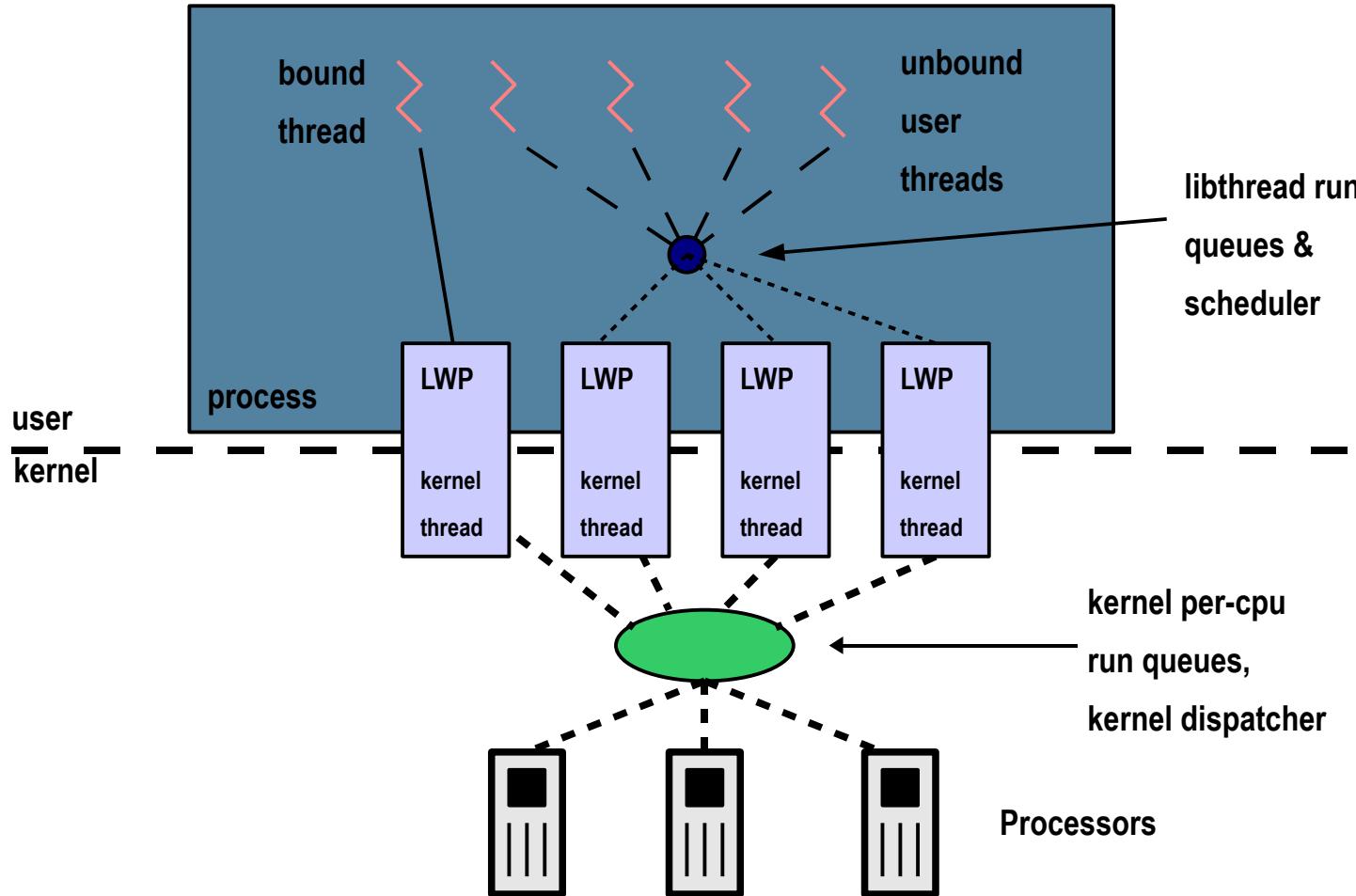
# T1 – Multilevel MxN Model

- /usr/lib/libthread.so.1
- Based on the assumption that kernel threads are expensive, user threads are cheap.
- User threads are virtualized, and may be multiplexed onto one or more kernel threads
  - LWP pool
- User level thread synchronization - threads sleep at user level. (Process private only)
- Concurrency via `set_concurrency()` and bound LWPs

# T1 – Multilevel Model

- Unbound Thread Implementation
  - User Level scheduling
  - Unbound threads switched onto available Iwps
  - Threads switched when blocked on sync object
  - Thread temporary bound when blocked in system call
  - Daemon Iwp to create new Iwps
  - Signal direction handled by Daemon Iwp
  - Reaper thread to manage cleanup
  - Callout Iwp for timers

## T1- Multilevel Model (default in Solaris 8)



# T1 – Multilevel Model

- Pros:

- Fast user thread create and destroy
- Allows many-to-few thread model, to minimize the number of kernel threads and LWPs
- Uses minimal kernel memory
- No system call required for synchronization
- Process Private Synchronization only
- Can have thousands of threads
- Fast context-switching

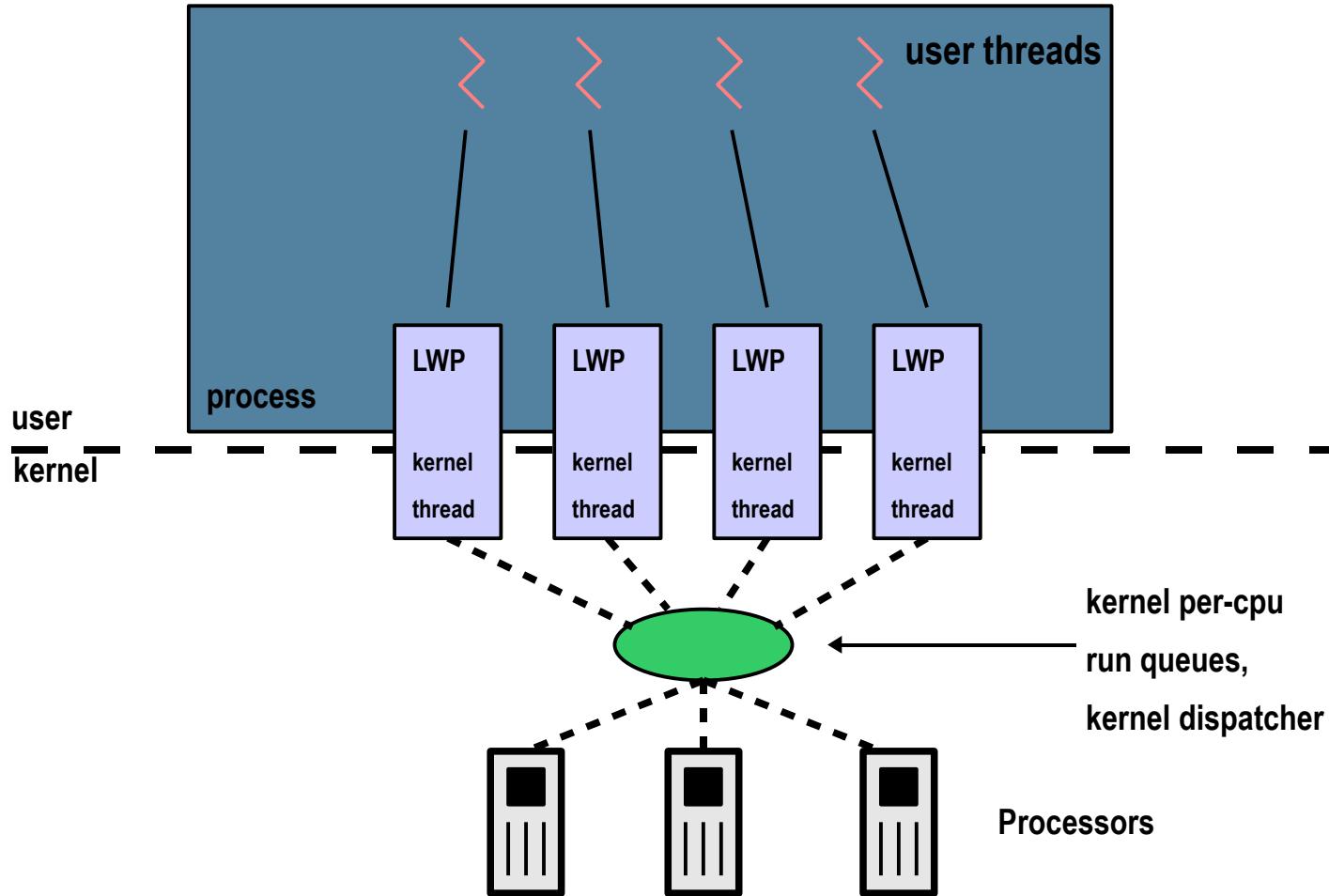
- Cons:

- Complex, and tricky programming model wrt achieving good scalability - need to bind or use `set_concurrency()`
- Signal delivery
- Compute bound threads do not surrender, leading to excessive CPU consumption and potential starving
- Complex to maintain (for Sun)

## T2 – Single Level Threads Model

- All user threads bound to LWPs
  - All bound threads
- Kernel level scheduling
  - No more libthread.so scheduler
- Simplified Implementation
- Uses kernel's synchronization objects
  - Slightly different behaviour LIFO vs. FIFO
  - Allows adaptive lock behaviour
- More expensive thread create/destroy, synchronization
- More responsive scheduling, synchronization

# T2 – Single Level Threads Model



## T2 - Single Level Thread Model

- Scheduling wrt Synchronization (S8U7/S9/S10)
  - Adaptive locks give preference to a thread that is running, potentially at the expense of a thread that is sleeping
  - Threads that rely on fairness of scheduling/CPU could end up ping-ponging, at the expense of another thread which has work to do.
- Default S8U7/S9/S10 Behavior
  - Adaptive Spin
    - 1000 of iterations (spin count) for adaptive mutex locking before giving up and going to sleep.
  - Maximum number of spinners
    - The number of simultaneously spinning threads
    - attempting to do adaptive locking on one mutex is limited to 100.
  - One out of every 16 queuing operations will put a thread at the end of the queue, to prevent starvation.
  - Stack Cache
    - The maximum number of stacks the library retains after threads exit for re-use when more threads are created is 10.



## Thread Semantics Added to pstack, truss

```
# pstack 909/2
909:    dbwr -a dbwr -i 2 -s b0000000 -m /var/tmp/fbencAAAmxaqxb
----- lwp# 2 -----
ceab1809 lwp_park (0, afffdde50, 0)
ceaabf93 cond_wait_queue (ce9f8378, ce9f83a0, afffdde50, 0) + 3b
ceaac33f cond_wait_common (ce9f8378, ce9f83a0, afffdde50) + 1df
ceaac686 _cond_reltimedwait (ce9f8378, ce9f83a0, afffdea0) + 36
ceaac6b4 cond_reltimedwait (ce9f8378, ce9f83a0, afffdea0) + 24
ce9e5902 __aio_waitn (82d1f08, 1000, afffdf2c, afffdf18, 1) + 529
ceaf2a84 aio_waitn64 (82d1f08, 1000, afffdf2c, afffdf18) + 24
08063065 flowoplib_aiowait (b4eb475c, c40f4d54) + 97
08061de1 flowop_start (b4eb475c) + 257
ceab15c0 _thr_setup (ce9a8400) + 50
ceab1780 _lwp_start (ce9a8400, 0, 0, afffdff8, ceab1780, ce9a8400)
pael> truss -p 2975/3
/3: close(5) = 0
/3: open("/space1/3", O_RDWR|O_CREAT, 0666) = 5
/3: lseek(5, 0, SEEK_SET) = 0
/3: write(5, " U U U U U U U U U U U U U U"..., 1056768) = 1056768
/3: lseek(5, 0, SEEK_SET) = 0
/3: read(5, " U U U U U U U U U U U U U U"..., 1056768) = 1056768
/3: close(5) = 0
/3: open("/space1/3", O_RDWR|O_CREAT, 0666) = 5
/3: lseek(5, 0, SEEK_SET) = 0
/3: write(5, " U U U U U U U U U U U U U U"..., 1056768) = 1056768
```



# Thread Microstates

PID	USERNAME	USR	SYS	TRP	TFL	DFL	LCK	SLP	LAT	VCX	ICX	SCL	SIG	PROCESS/LWPID
918	rmc	0.2	0.4	0.0	0.0	0.0	0.0	0.0	99	0.0	27	2	1K	0 prstat/1
919	mauroj	0.1	0.4	0.0	0.0	0.0	0.0	0.0	99	0.1	44	12	1K	0 prstat/1
907	root	0.0	0.1	0.0	0.0	0.0	0.0	0.0	97	3.1	121	2	20	0 filebench/2
913	root	0.1	0.0	0.0	0.0	0.0	0.0	100	0.0	0.0	15	2	420	0 filebench/2
866	root	0.0	0.0	0.0	0.0	0.0	0.0	0.0	96	4.1	44	41	398	0 filebench/2
820	root	0.0	0.0	0.0	0.0	0.0	0.0	0.0	95	5.0	43	42	424	0 filebench/2
814	root	0.0	0.0	0.0	0.0	0.0	0.0	0.0	95	5.0	43	41	424	0 filebench/2
772	root	0.0	0.0	0.0	0.0	0.0	0.0	0.0	96	3.6	46	39	398	0 filebench/2
749	root	0.0	0.0	0.0	0.0	0.0	0.0	0.0	96	3.7	45	41	398	0 filebench/2
744	root	0.0	0.0	0.0	0.0	0.0	0.0	0.0	95	4.7	47	39	398	0 filebench/2
859	root	0.0	0.0	0.0	0.0	0.0	0.0	0.0	95	4.9	44	41	424	0 filebench/2
837	root	0.0	0.0	0.0	0.0	0.0	0.0	0.0	96	4.0	43	43	405	0 filebench/2
[snip]														
787	root	0.0	0.0	0.0	0.0	0.0	0.0	0.0	95	4.5	43	41	424	0 filebench/2
776	root	0.0	0.0	0.0	0.0	0.0	0.0	0.0	95	4.8	43	42	398	0 filebench/2
774	root	0.0	0.0	0.0	0.0	0.0	0.0	0.0	96	4.2	43	40	398	0 filebench/2
756	root	0.0	0.0	0.0	0.0	0.0	0.0	0.0	96	3.8	44	41	398	0 filebench/2
738	root	0.0	0.0	0.0	0.0	0.0	0.0	0.0	96	4.4	43	42	398	0 filebench/2
735	root	0.0	0.0	0.0	0.0	0.0	0.0	0.0	96	3.9	47	39	405	0 filebench/2
734	root	0.0	0.0	0.0	0.0	0.0	0.0	0.0	96	4.3	44	41	398	0 filebench/2
727	root	0.0	0.0	0.0	0.0	0.0	0.0	0.0	96	4.4	43	43	398	0 filebench/2
725	root	0.0	0.0	0.0	0.0	0.0	0.0	0.0	96	4.4	43	43	398	0 filebench/2

Total: 257 processes, 3139 lwps, load averages: 7.71, 2.39, 0.97



# Watching Threads

PID	USERNAME	SIZE	RSS	STATE	PRI	NICE	TIME	CPU	PROCESS/LWPID
29105	root	5400K	3032K	sleep	60	0	0:00:00	1.3%	pkginstall/1
29051	root	5072K	4768K	cpu0	49	0	0:00:00	0.8%	prstat/1
202	root	3304K	1256K	sleep	59	0	0:00:07	0.3%	nscd/23
25947	root	5160K	608K	sleep	59	0	0:00:05	0.2%	sshd/1
23078	root	20M	1880K	sleep	59	0	0:00:58	0.2%	lupi_zones/1
25946	rmc	3008K	624K	sleep	59	0	0:00:02	0.2%	ssh/1
23860	root	5248K	688K	sleep	59	0	0:00:06	0.2%	sshd/1
29100	root	1272K	976K	sleep	59	0	0:00:00	0.1%	mpstat/1
24866	root	5136K	600K	sleep	59	0	0:00:02	0.0%	sshd/1
340	root	2504K	672K	sleep	59	0	0:11:14	0.0%	mibiisa/2
23001	root	5136K	584K	sleep	59	0	0:00:04	0.0%	sshd/1
830	root	2472K	600K	sleep	59	0	0:11:01	0.0%	mibiisa/2
829	root	2488K	648K	sleep	59	0	0:11:01	0.0%	mibiisa/2
1	root	2184K	400K	sleep	59	0	0:00:01	0.0%	init/1
202	root	3304K	1256K	sleep	59	0	0:00:00	0.0%	nscd/13
202	root	3304K	1256K	sleep	59	0	0:00:00	0.0%	nscd/12
202	root	3304K	1256K	sleep	59	0	0:00:00	0.0%	nscd/11
202	root	3304K	1256K	sleep	59	0	0:00:00	0.0%	nscd/10
202	root	3304K	1256K	sleep	59	0	0:00:00	0.0%	nscd/9
202	root	3304K	1256K	sleep	59	0	0:00:00	0.0%	nscd/8
202	root	3304K	1256K	sleep	59	0	0:00:00	0.0%	nscd/7
202	root	3304K	1256K	sleep	59	0	0:00:00	0.0%	nscd/6
202	root	3304K	1256K	sleep	59	0	0:00:00	0.0%	nscd/5
202	root	3304K	1256K	sleep	59	0	0:00:00	0.0%	nscd/4
202	root	3304K	1256K	sleep	59	0	0:00:00	0.0%	nscd/3
202	root	3304K	1256K	sleep	59	0	0:00:00	0.0%	nscd/2
202	root	3304K	1256K	sleep	59	0	0:00:00	0.0%	nscd/1
126	daemon	2360K	8K	sleep	59	0	0:00:00	0.0%	rpcbind/1
814	root	1936K	280K	sleep	59	0	0:00:00	0.0%	sac/1
64	root	2952K	8K	sleep	59	0	0:00:00	0.0%	picld/5
64	root	2952K	8K	sleep	59	0	0:00:00	0.0%	picld/4
64	root	2952K	8K	sleep	59	0	0:00:00	0.0%	picld/3
64	root	2952K	8K	sleep	59	0	0:00:00	0.0%	picld/2
64	root	2952K	8K	sleep	59	0	0:00:00	0.0%	picld/1
61	daemon	3640K	8K	sleep	59	0	0:00:00	0.0%	kcf/3
61	daemon	3640K	8K	sleep	59	0	0:00:00	0.0%	kcf/2
61	daemon	3640K	8K	sleep	59	0	0:00:00	0.0%	kcf/1
55	root	2416K	8K	sleep	59	0	0:00:00	0.0%	syseventd/14
55	root	2416K	8K	sleep	59	0	0:00:00	0.0%	syseventd/13
55	root	2416K	8K	sleep	59	0	0:00:00	0.0%	syseventd/12
55	root	2416K	8K	sleep	59	0	0:00:00	0.0%	syseventd/11

Total: 125 processes, 310 lwp's, load averages: 0.50, 0.38, 0.40



# Who's Creating Threads?

```
# dtrace -n 'thread_create:entry { @[execname]=count()}'  
dtrace: description 'thread_create:entry' matched 1 probe  
^C
```

sh	1
sched	1
do1.6499	2
do1.6494	2
do1.6497	2
do1.6508	2
in.rshd	12
do1.6498	14
do1.6505	16
do1.6495	16
do1.6504	16
do1.6502	16
automountd	17
inetd	19
filebench	34
find	130
csh	177



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# Scheduling Classes & The Kernel Dispatcher

## Solaris Scheduling

- Solaris implements a central dispatcher, with multiple scheduling classes
  - Scheduling classes determine the priority range of the kernel threads on the system-wide (global) scale, and the scheduling algorithms applied
  - Each scheduling class references a dispatch table
    - Values used to determine time quanta and priorities
    - Admin interface to “tune” thread scheduling
  - Solaris provides command line interfaces for
    - Loading new dispatch tables
    - Changing the scheduling class and priority and threads
  - Observability through
    - ps(1)
    - prstat(1)
    - dtrace(1)



# Scheduling Classes

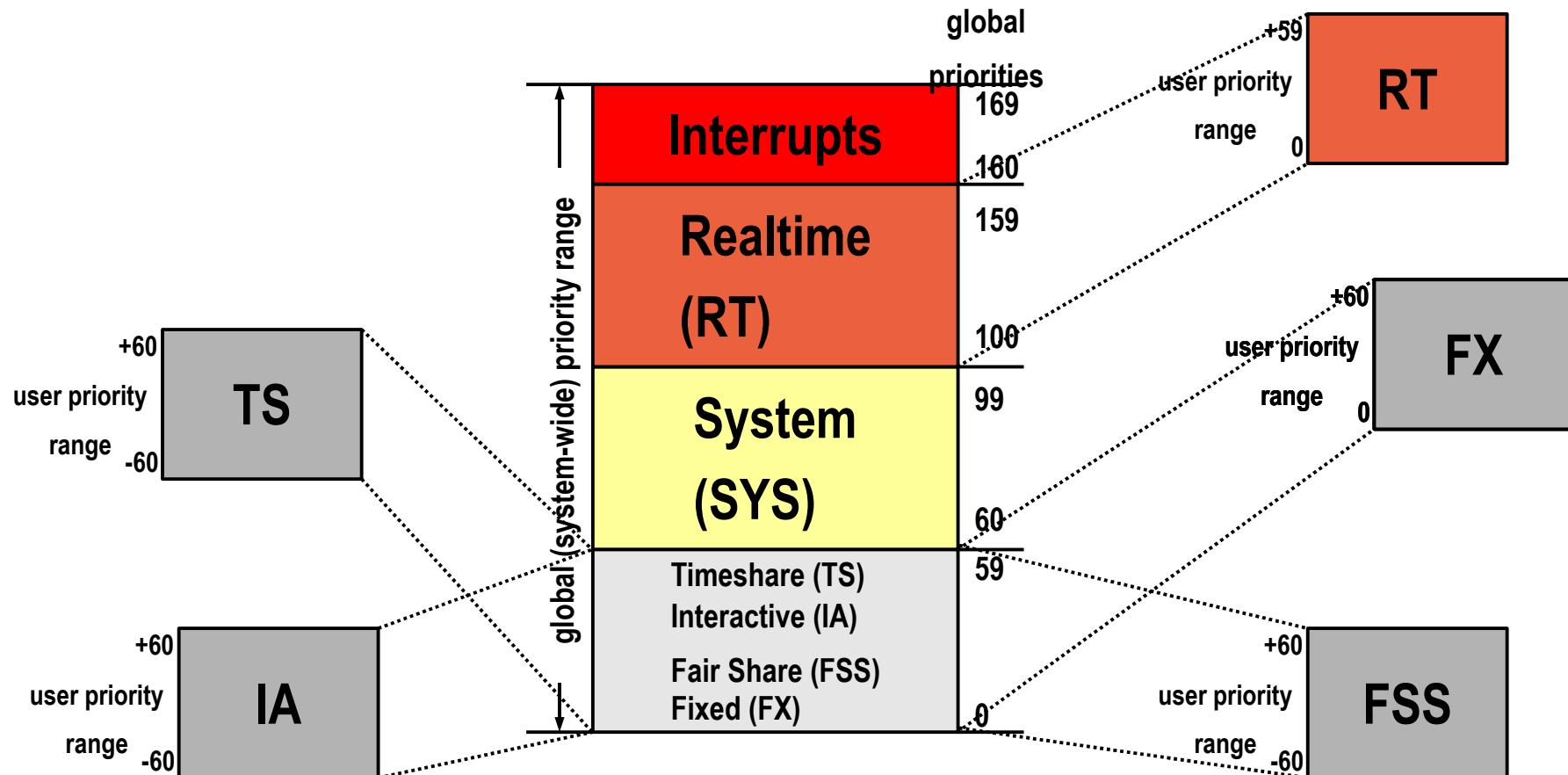
- Traditional Timeshare (TS) class
  - attempt to give every thread a fair shot at execution time
- Interactive (IA) class
  - Desktop only
  - Boost priority of active (current focus) window
  - Same dispatch table as TS
- System (SYS)
  - Only available to the kernel, for OS kernel threads
- Realtime (RT)
  - Highest priority scheduling class
  - Will preempt kernel (SYS) class threads
  - Intended for realtime applications
    - Bounded, consistent scheduling latency



# Scheduling Classes – Solaris 9 & 10

- Fair Share Scheduler (FSS) Class
  - Same priority range as TS/IA class
  - CPU resources are divided into shares
  - Shares are allocated (projects/tasks) by administrator
  - Scheduling decisions made based on shares allocated and used, not dynamic priority changes
- Fixed Priority (FX) Class
  - The kernel will not change the thread's priority
  - A “batch” scheduling class
- Same set of commands for administration and management
  - `dispadmin(1M)`, `priocntl(1)`
  - Resource management framework
    - `rctladm(1M)`, `prctl(1)`

# Scheduling Classes and Priorities





# Scheduling Classes

- Use `dispadmin(1M)` and `priocntl(1)`

```
# dispadmin -l
CONFIGURED CLASSES
=====

SYS  (System Class)
TS   (Time Sharing)
FX   (Fixed Priority)
IA   (Interactive)
FSS  (Fair Share)
RT   (Real Time)
# priocntl -l
CONFIGURED CLASSES
=====

SYS (System Class)

TS (Time Sharing)
    Configured TS User Priority Range: -60 through 60

FX (Fixed priority)
    Configured FX User Priority Range: 0 through 60

IA (Interactive)
    Configured IA User Priority Range: -60 through 60

FSS (Fair Share)
    Configured FSS User Priority Range: -60 through 60

RT (Real Time)
    Maximum Configured RT Priority: 59
#
```



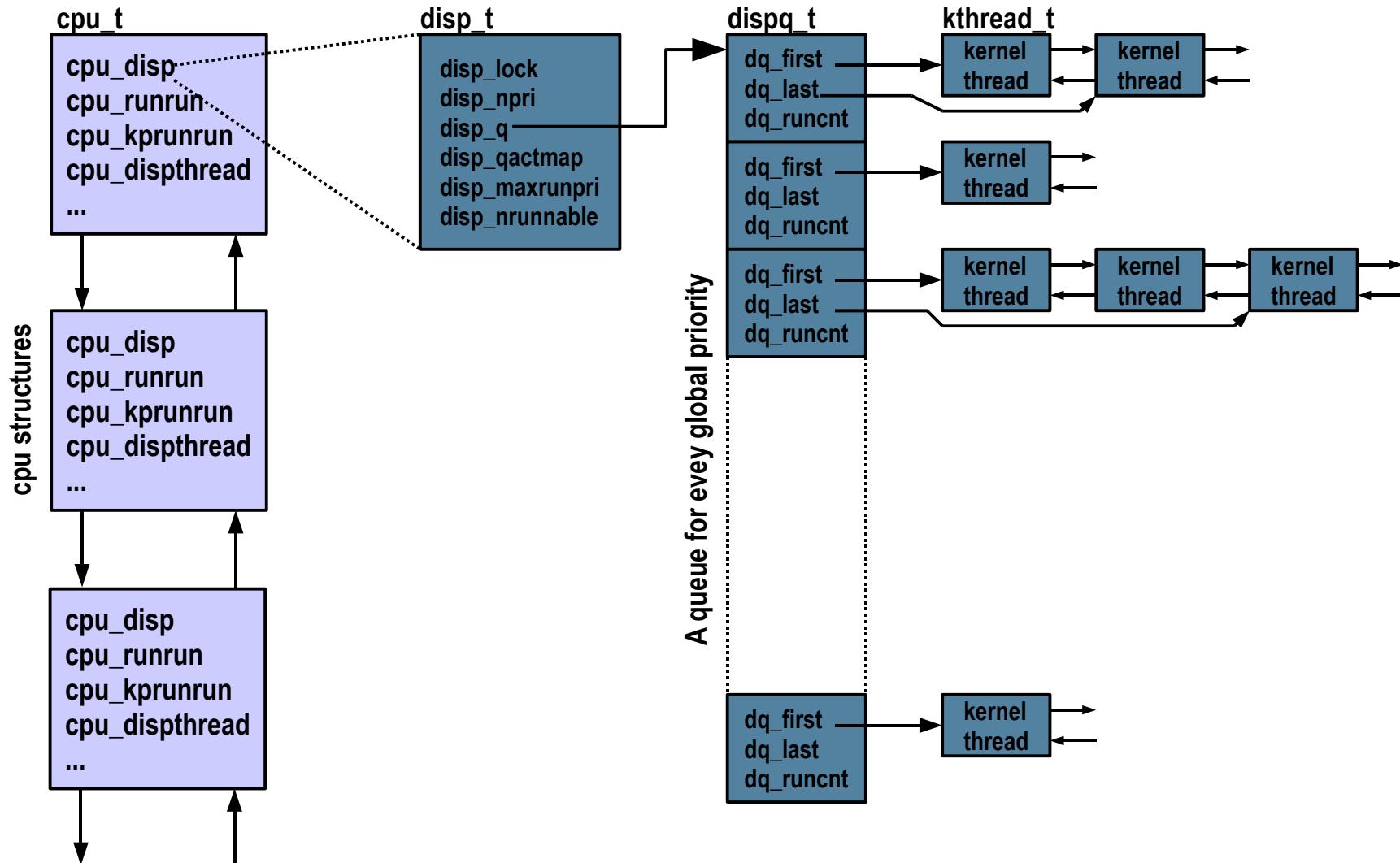
# Scheduling Class & Priority of Threads

```
solaris10> ps -eLc
  PID  LWP  CLS PRI TTY      LTIME CMD
    0    1  SYS  96 ?      0:00  sched
    1    1  TS   59 ?      0:00  init
    2    1  SYS  98 ?      0:00  pageout
    3    1  SYS  60 ?      5:08  fsflush
  402    1  TS   59 ?
  269    1  TS   59 ?
  225    1  TS   59 ?
  225    2  TS   59 ?
  225    4  TS   59 ?
    54    1  TS   59 ?
    54    2  TS   59 ?
    54    3  TS   59 ?
  [snip]
  426    1  IA   59 ?
  343    1  TS   59 ?
  345    1  FX   60 ?
  345    3  FX   60 ?
  350    1  TS   59 ?
  375    1  TS   59 ?
  411    1  IA   59 ?
  412    1  IA   59 ???
  403    1  TS   59 console
  405    1  TS   59 ?
  406    1  IA   59 ?
  410    1  TS   59 ?
  409    1  TS   59 ?
 1040    1  TS   59 ?
 1059    1  TS   49 pts/2
  solaris10>
```

# Dispatch Queues & Dispatch Tables

- Dispatch queues
  - Per-CPU run queues
    - Actually, a queue of queues
    - Ordered by thread priority
    - Queue occupation represented via a bitmap
    - For Realtime threads, a system-wide kernel preempt queue is maintained
      - Realtime threads are placed on this queue, not the per-CPU queues
      - If processor sets are configured, a kernel preempt queue exists for each processor set
  - Dispatch tables
    - Per-scheduling class parameter tables
    - Time quantums and priorities
    - tuneable via `dispadmin(1M)`

# Per-CPU Dispatch Queues





## Timeshare Dispatch Table

- TS and IA class share the same dispatch table
  - *RES*. Defines the granularity of *ts\_quantum*
  - *ts\_quantum*. CPU time for next ONPROC state
  - *ts\_tqexp*. New priority if time quantum expires
  - *ts\_slpret*. New priority when state change from TS\_SLEEP to TS\_RUN
  - *ts\_maxwait*. “waited too long” ticks
  - *ts\_lwait*. New priority if “waited too long”

```
# dispadmin -g -c TS
# Time Sharing Dispatcher Configuration
RES=1000

# ts_quantum          ts_tqexp ts_slpret ts_maxwait      ts_lwait      PRIORITY LEVEL
  200                 0         50          0            50          #
  200                 0         50          0            50          #
  .....
  160                 0         51          0            51          #
  160                 1         51          0            51          #
  .....
  120                10        52          0            52          #
  120                11        52          0            52          #
  .....
  80                  20        53          0            53          #
  80                  21        53          0            53          #
  .....
  40                  30        55          0            55          #
  40                  31        55          0            55          #
  .....
  20                  49        59          32000        59          #       59
```

# RT, FX & FSS Dispatch Tables

- RT
  - Time quantum only
  - For each possible priority
- FX
  - Time quantum only
  - For each possible priority
- FSS
  - Time quantum only
  - Just one, not defined for each priority level
    - Because FSS is share based, not priority based
- SYS
  - No dispatch table
  - Not needed, no rules apply
- INT
  - Not really a scheduling class



# Dispatch Queue Placement

- Queue placement is based a few simple parameters
  - The thread priority
  - Processor binding/Processor set
  - Processor thread last ran on
    - Warm affinity
  - Depth and priority of existing runnable threads
  - Solaris 9 added Memory Placement Optimization (MPO) enabled will keep thread in defined locality group (lgroup)

```
if (thread is bound to CPU-n) && (pri < kpreemptpri)
    CPU-n dispatch queue
if (thread is bound to CPU-n) && (pri >= kpreemptpri)
    CPU-n dispatch queue
if (thread is not bound) && (pri < kpreemptpri)
    place thread on a CPU dispatch queue
if (thread is not bound) && (pri >= kpreemptpri)
    place thread on cp_kp_queue
```

# Thread Preemption

- Two classes of preemption
  - User preemption
    - A higher priority thread became runnable, but it's not a realtime thread
    - Flagged via `cpu_runrun` in CPU structure
    - Next clock tick, you're outta here
  - Kernel preemption
    - A realtime thread became runnable. Even OS kernel threads will get preempted
    - Poke the CPU (cross-call) and preempt the running thread now
  - *Note that threads that use-up their time quantum are evicted via the preempt mechanism*
  - Monitor via “icsw” column in `mpstat(1)`



# Thread Execution

- Run until
  - A preemption occurs
    - Transition from S\_ONPROC to S\_RUN
    - placed back on a run queue
  - A blocking system call is issued
    - e.g. read(2)
    - Transition from S\_ONPROC to S\_SLEEP
    - Placed on a sleep queue
  - Done and exit
    - Clean up
  - Interrupt to the CPU you're running on
    - pinned for interrupt thread to run
    - unpinned to continue



# Context Switching

CPU	minf	mjf	xcal	intr	ithr	csw	icsw	migr	smtx	srw	syscl	usr	sys	wt	idl
0	74	2	998	417	302	450	18	45	114	0	1501	56	7	0	37
1	125	3	797	120	102	1107	16	58	494	0	1631	41	16	0	44
4	209	2	253	114	100	489	12	45	90	0	1877	56	11	0	33
5	503	7	2448	122	100	913	21	53	225	0	2626	32	21	0	48
8	287	3	60	120	100	771	20	35	122	0	1569	50	12	0	38
9	46	1	51	115	99	671	16	20	787	0	846	81	16	0	3
12	127	2	177	117	101	674	14	27	481	0	881	74	12	0	14
13	375	7	658	1325	1302	671	23	49	289	0	1869	48	16	0	37
CPU	minf	mjf	xcal	intr	ithr	csw	icsw	migr	smtx	srw	syscl	usr	sys	wt	idl
0	0	0	733	399	297	548	10	8	653	0	518	80	11	0	9
1	182	4	45	117	100	412	16	34	49	0	904	54	8	0	38
4	156	4	179	108	102	1029	6	46	223	0	1860	15	16	0	70
5	98	1	53	110	100	568	9	19	338	0	741	60	9	0	31
8	47	1	96	111	101	630	6	22	712	0	615	56	13	0	31
9	143	4	127	116	102	1144	11	42	439	0	2443	33	15	0	52
12	318	0	268	111	100	734	9	30	96	0	1455	19	12	0	69
13	39	2	16	938	929	374	8	9	103	0	756	69	6	0	25



```
#!/usr/sbin/dtrace -Zqs
long inv_cnt; /* all involuntary context switches */
long tqe_cnt; /* time quantum expiration count */
long hpp_cnt; /* higher-priority preempt count */
long csw_cnt; /* total number context switches */

dtrace:::BEGIN
{
    inv_cnt = 0; tqe_cnt = 0; hpp_cnt = 0; csw_cnt = 0;

    printf("%-16s %-16s %-16s %-16s\n", "TOTAL CSW", "ALL INV", "TQE_INV", "HPP_INV");
    printf("=====\\n");
}

sysinfo:unix:preempt:inv_swtch
{
    inv_cnt += arg0;
}
sysinfo:unix::pswitch
{
    csw_cnt += arg0;
}

fbt:TS:ts_preempt:entry
/ ((tsproc_t *)args[0]->t_cldata)->ts_timeleft <= 1 /
{
    tqe_cnt++;
}

fbt:TS:ts_preempt:entry
/ ((tsproc_t *)args[0]->t_cldata)->ts_timeleft > 1 /
{
    hpp_cnt++;
}

fbt:RT:rt_preempt:entry
/ ((rtproc_t *)args[0]->t_cldata)->rt_timeleft <= 1 /
{
    tqe_cnt++;
}

fbt:RT:rt_preempt:entry
/ ((rtproc_t *)args[0]->t_cldata)->rt_timeleft > 1 /
{
    hpp_cnt++;
}

tick-1sec
{
    printf("%-16d %-16d %-16d %-16d\\n", csw_cnt, inv_cnt, tqe_cnt, hpp_cnt);

    inv_cnt = 0; tqe_cnt = 0; hpp_cnt = 0; csw_cnt = 0;
}
```



```
solaris10> ./csw.d
```

TOTAL CSW	ALL_INV	TQE_INV	HPP_INV
1544	63	24	40
3667	49	35	14
4163	59	34	26
3760	55	29	26
3839	71	39	32
3931	48	33	15

^C

```
solaris10> ./threads &
```

```
[2] 19913
```

```
solaris10>
```

```
solaris10> ./csw.d
```

TOTAL CSW	ALL_INV	TQE_INV	HPP_INV
3985	1271	125	1149
5681	1842	199	1648
5025	1227	151	1080
9170	520	108	412
4100	390	84	307
2487	174	74	99
1841	113	64	50
6239	170	74	96

^C

1440	155	68	88
------	-----	----	----

# Observability and Performance

- Use `prstat(1)` and `ps(1)` to monitor running processes and threads
- Use `mpstat(1)` to monitor CPU utilization, context switch rates and thread migrations
- Use `dispadmin(1M)` to examine and change dispatch table parameters
- Use `priocntl(1)` to change scheduling classes and priorities
  - `nice(1)` is obsolete (but there for compatibility)
  - User priorities also set via `priocntl(1)`
  - Must be root to use RT class



## Dtrace sched provider probes:

- change-pri – change pri
- dequeue – exit run q
- enqueue – enter run q
- off-cpu – start running
- on-cpu – stop running
- preempt - preempted
- remain-cpu
- schedctl-nopreempt – hint that it is not ok to preempt
- schedctl-preempt – hint that it is ok to preempt
- schedctl-yield - hint to give up runnable state
- sleep – go to sleep
- surrender – preempt from another cpu
- tick – tick-based accounting
- wakeup – wakeup from sleep



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# Processors, Processor Controls & Binding

# Processor Controls

- Processor controls provide for segregation of workload(s) and resources
- Processor status, state, management and control
  - Kernel linked list of CPU structs, one for each CPU
  - Bundled utilities
    - psradm(1)
    - psrinfo(1)
  - Processors can be taken offline
    - Kernel will not schedule threads on an offline CPU
  - The kernel can be instructed not to bind device interrupts to processor(s)
    - Or move them if bindings exist



# Processor Control Commands

- `psrinfo(1M)` - provides information about the processors on the system. Use "-v" for verbose
- `psradm(1M)` - online/offline processors. Pre Sol 7, offline processors still handled interrupts. In Sol 7, you can disable interrupt participation as well
- `psrset(1M)` - creation and management of processor sets
- `pbind(1M)` - original processor bind command. Does not provide exclusive binding
- `processor_bind(2)`, `processor_info(2)`,  
`pset_bind(2)`, `pset_info(2)`, `pset_creat(2)`,  
`p_online(2)`
  - system calls to do things programmatically

# Processor Sets

- Partition CPU resources for segregating workloads, applications and/or interrupt handling
- Dynamic
  - Create, bind, add, remove, etc, without reboots
- Once a set is created, the kernel will only schedule threads onto the set that have been explicitly bound to the set
  - And those threads will only ever be scheduled on CPUs in the set they've been bound to
- Interrupt disabling can be done on a set
  - Dedicate the set, through binding, to running application threads
  - Interrupt segregation can be effective if interrupt load is heavy
    - e.g. high network traffic



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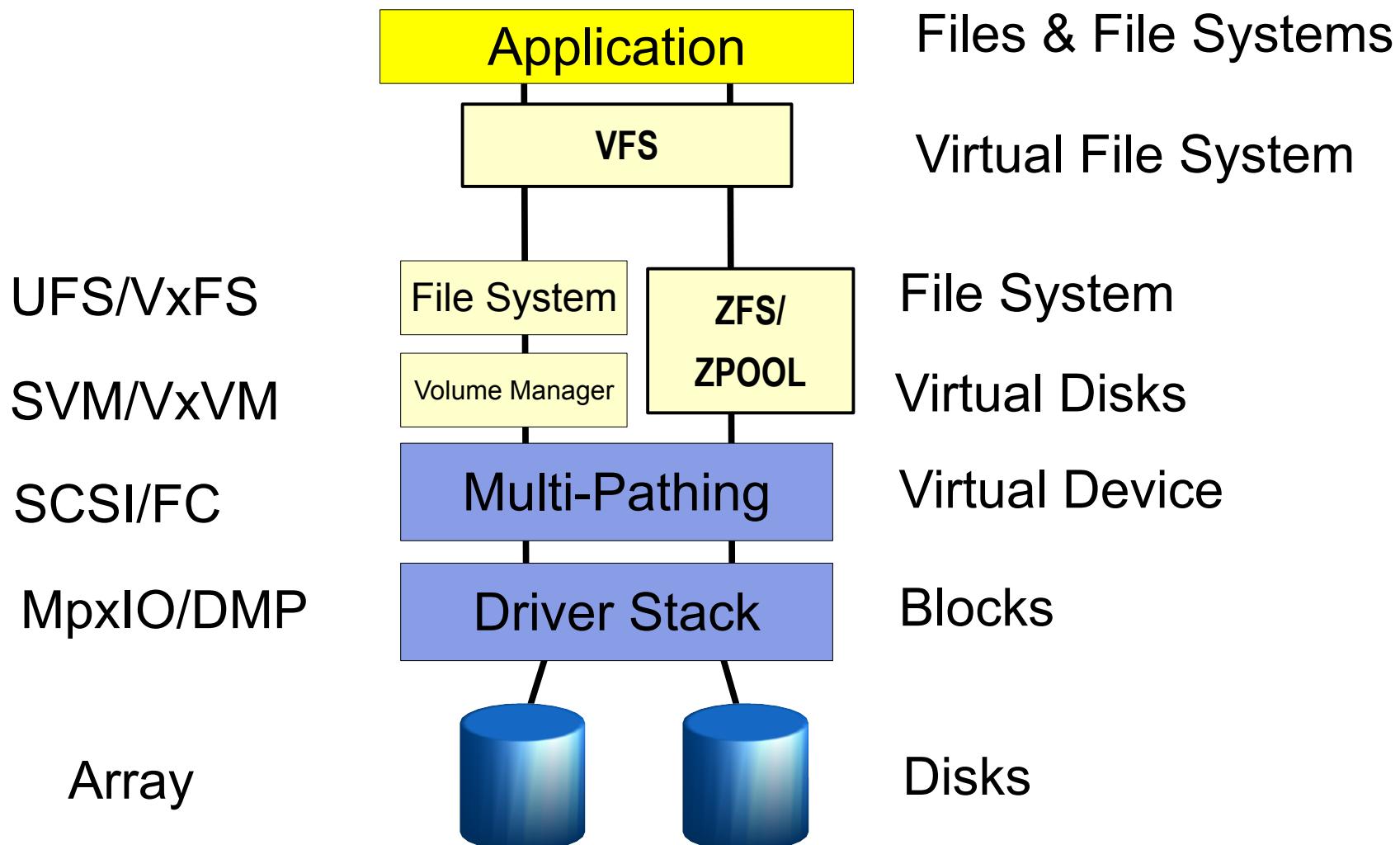
# Session 4

# File Systems &

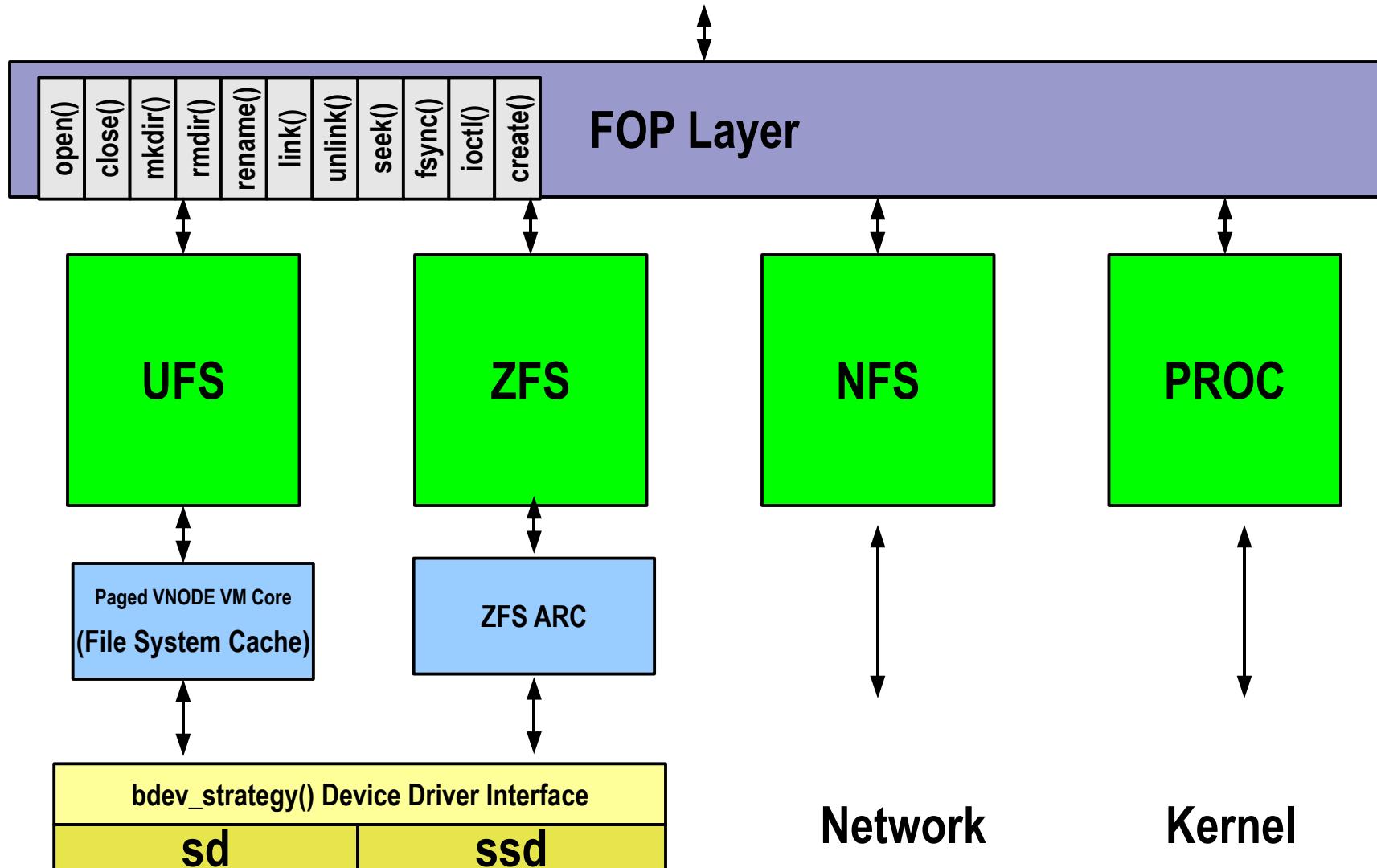
# Disk I/O

# Performance

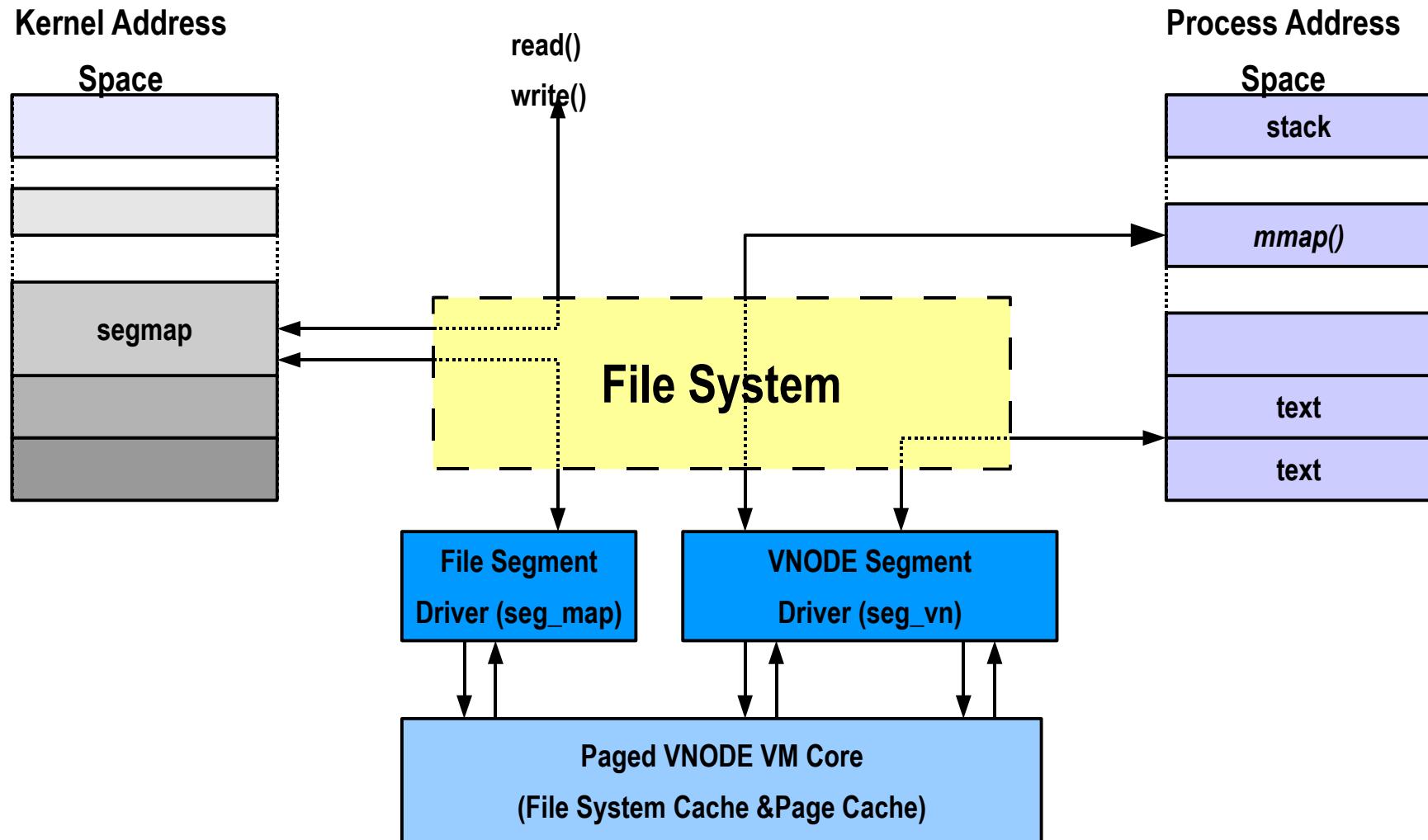
# The Solaris File System/IO Stack



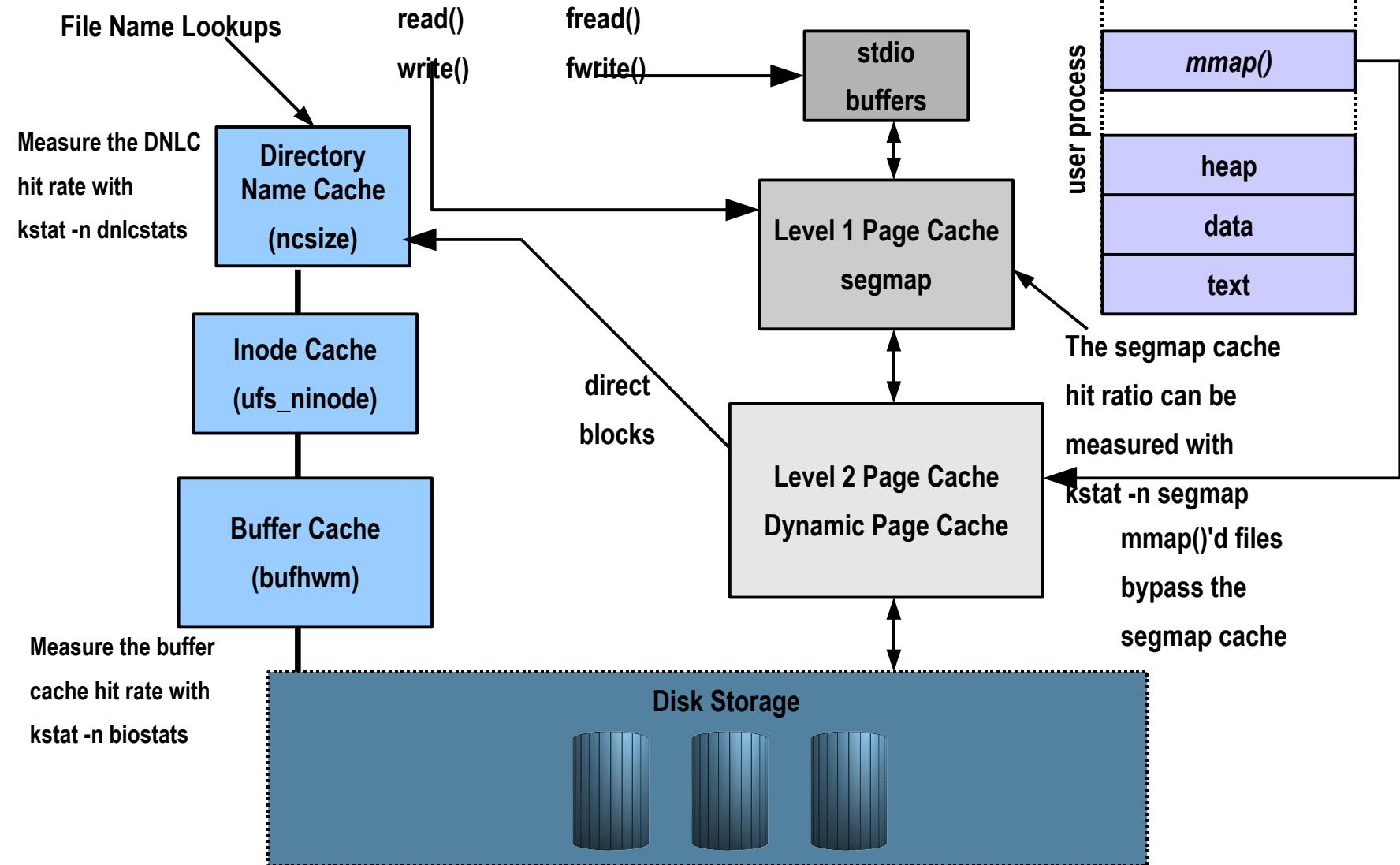
# File System Architecture



# UFS I/O



# UFS Caching



# Filesystem performance

- Attribution
  - How much is my application being slowed by I/O?
  - i.e. How much faster would my app run if I optimized I/O?
- Accountability
  - What is causing I/O device utilization?
  - i.e. What user is causing this disk to be hot?
- Tuning/Optimizing
  - Tuning for sequential, random I/O and/or meta-data intensive applications

# Solaris FS Perf Tools

- iostat: raw disk statistics
- sar -b: meta-data buffer cachestat
- vmstat -s: monitor dnlc
- Filebench: emulate and measure various FS workloads
- DTrace: trace physical I/O – IO provider
- DTrace: fsinfo provider
- DTrace: top for files – logical and physical per file
- DTrace: top for fs – logical and physical per filesystem
- DTrace Toolkit – iosnoop and iotop

## Simple performance model

- Single-threaded processes are simpler to estimate
  - Calculate elapsed vs. waiting for I/O time, express as a percentage
  - i.e. My app spent 80% of its execution time waiting for I/O
  - Inverse is potential speed up – e.g. 80% of time waiting equates to a potential 5x speedup
  - The key is to estimate the time spent waiting



# Estimating wait time

- Elapsed vs. cpu seconds
  - Time <cmd>, estimate wait as real – user - sys
- Etruss
  - Uses microstates to estimate I/O as wait time
  - <http://www.solarisinternals.com>
- Measure explicitly with dtrace
  - Measure and total I/O wait per thread



# Examining IO wait with dtrace

- Measuring on-cpu vs io-wait time:

```
sol10$ ./iowait.d 639
^C

Time breakdown (milliseconds):
<on cpu>                                2478
<I/O wait>                               6326

I/O wait breakdown (milliseconds):
file1                                     236
file2                                     241
file4                                     244
file3                                     264
file5                                     277
file7                                     330

.
```

# Solaris iostat

```
# iostat -xsz
              extended device statistics
      r/s     w/s    kr/s   kw/s  wait  actv wsvc_t asvc_t  %w  %b device
  687.8     0.0 38015.3     0.0   0.0   1.9    0.0    2.7    0 100 c0d0
```



wait

svc

- Wait: number of threads queued for I/O
- Actv: number of threads performing I/O
- wsvc\_t: Average time spend waiting on queue
- asvc\_t: Average time performing I/O
- %w: Only useful if one thread is running on the entire machine – time spent waiting for I/O
- %b: Device utilization – only useful if device can do just 1 I/O at a time (invalid for arrays etc...)



## Thread I/O example

```
sol$ cd labs/disks
sol$ ./lthread
1079: 0.007: Random Read Version 1.8 05/02/17 IO personality successfully loaded
1079: 0.008: Creating/pre-allocating files
1079: 0.238: Waiting for preallocation threads to complete...
1079: 0.238: Re-using file /filebench/bigfile0
1079: 0.347: Starting 1 rand-read instances
1080: 1.353: Starting 1 rand-thread threads
1079: 4.363: Running for 600 seconds...
sol$ iostat -xncz 5
      cpu
us sy wt id
22  3  0 75
                           extended device statistics
      r/s     w/s    kr/s    kw/s wait  actv wsvc_t asvc_t  %w  %b device
  62.7     0.3   501.4     2.7   0.0   0.9     0.0   14.1     0   89 cld0
```



# 64 Thread I/O example

```
sol$ cd labs/disks
sol$ ./64thread
1089: 0.095: Random Read Version 1.8 05/02/17 IO personality successfully loaded
1089: 0.096: Creating/pre-allocating files
1089: 0.279: Waiting for preallocation threads to complete...
1089: 0.279: Re-using file /filebench/bigfile0
1089: 0.385: Starting 1 rand-read instances
1090: 1.389: Starting 64 rand-thread threads
1089: 4.399: Running for 600 seconds...
```

```
sol$ iostat -xncz 5
      cpu
us sy wt id
15 1 0 83
                           extended device statistics
 r/s    w/s    kr/s    kw/s wait  actv wsvc_t asvc_t  %w  %b device
 71.0     0.3   568.0    17.3 61.8   2.0   866.5    28.0 100 100 cld0
```



## Solaris iostat

- New Formatting flags -C, -l, -m, -r, -s, -z, -T
  - -C: report disk statistics by controller
  - -l n: Limit the number of disks to n
  - -m: Display mount points (most useful with -p)
  - -r: Display data n comma separated format
  - -s: Suppress state change messages
  - -z: Suppress entries with all zero values
  - -T d|u Display a timestamp in date (d) or unix time\_t (u)



# Examining Physical IO by file with dtrace

```
#pragma D option quiet

BEGIN
{
    printf("%10s %58s %2s %8s\n", "DEVICE", "FILE", "RW", "Size");
}

io:::start
{
    printf("%10s %58s %2s %8d\n", args[1]->dev_statname,
           args[2]->fi.pathname, args[0]->b_flags & B_READ ? "R" : "W",
           args[0]->b_bcount);
}

# dtrace -s ./iotrace
```

DEVICE	FILE	RW	SIZE
cmdk0	/export/home/rmc/.sh_history	W	4096
cmdk0	/opt/Acrobat4/bin/acroread	R	8192
cmdk0	/opt/Acrobat4/bin/acroread	R	1024
cmdk0	/var/tmp/wscon-:0.0-gLaW9a	W	3072
cmdk0	/opt/Acrobat4/Reader/AcroVersion	R	1024
cmdk0	/opt/Acrobat4/Reader/intelsolaris/bin/acroread	R	8192
cmdk0	/opt/Acrobat4/Reader/intelsolaris/bin/acroread	R	8192
cmdk0	/opt/Acrobat4/Reader/intelsolaris/bin/acroread	R	4096
cmdk0	/opt/Acrobat4/Reader/intelsolaris/bin/acroread	R	8192
cmdk0	/opt/Acrobat4/Reader/intelsolaris/bin/acroread	R	8192



# Physical Trace Example

```
sol8$ cd labs/disks
sol8$ ./64thread
1089: 0.095: Random Read Version 1.8 05/02/17 IO personality successfully loaded
1089: 0.096: Creating/pre-allocating files
1089: 0.279: Waiting for preallocation threads to complete...
1089: 0.279: Re-using file /filebench/bigfile0
1089: 0.385: Starting 1 rand-read instances
1090: 1.389: Starting 64 rand-thread threads
1089: 4.399: Running for 600 seconds...
```

```
sol8$ iotrace.d
```

DEVICE	FILE	RW	Size
cmdk0	/filebench/bigfile0	R	8192
cmdk0	/filebench/bigfile0	R	8192
cmdk0	/filebench/bigfile0	R	8192
cmdk0	/filebench/bigfile0	R	8192
cmdk0	/filebench/bigfile0	R	8192
cmdk0	/filebench/bigfile0	R	8192
cmdk0	/filebench/bigfile0	R	8192
cmdk0	/filebench/bigfile0	R	8192
cmdk0	/filebench/bigfile0	R	8192
cmdk0	/filebench/bigfile0	R	8192
cmdk0	/filebench/bigfile0	R	8192
cmdk0	/filebench/bigfile0	R	8192
cmdk0	/filebench/bigfile0	R	8192
cmdk0	/filebench/bigfile0	R	8192
cmdk0	/filebench/bigfile0	R	8192
cmdk0	/filebench/bigfile0	R	8192



## DTrace Toolkit - iotop

```
# iotop -C
Sampling... Please wait.
2005 Jul 16 00:34:40, load: 1.21, disk_r: 12891 Kb, disk_w: 1087 Kb

UID      PID    PPID  CMD                DEVICE   MAJ MIN D          BYTES
 0        3        0  fsflush            cmdk0    102   4 W          512
 0        3        0  fsflush            cmdk0    102   0 W         11776
 0    27751  20320  tar               cmdk0    102  16 W         23040
 0        3        0  fsflush            cmdk0    102   0 R         73728
 0        0        0  sched             cmdk0    102   0 R        548864
 0        0        0  sched             cmdk0    102   0 W        1078272
 0    27751  20320  tar               cmdk0    102  16 R        1514496
 0    27751  20320  tar               cmdk0    102   3 R        11767808
```



# DTrace Toolkit - iosnoop

```
# iosnoop
```

UID	PID	D	BLOCK	SIZE	COMM	PATHNAME
100	15795	R	3808	8192	tar	/usr/bin/eject
100	15795	R	35904	6144	tar	/usr/bin/eject
100	15795	R	39828	6144	tar	/usr/bin/env
100	15795	R	3872	8192	tar	/usr/bin/expr
100	15795	R	21120	7168	tar	/usr/bin/expr
100	15795	R	43680	6144	tar	/usr/bin/false
100	15795	R	44176	6144	tar	/usr/bin/fdetach
100	15795	R	3920	8192	tar	/usr/bin/fdformat
100	15795	R	3936	8192	tar	/usr/bin/fdformat
100	15795	R	4080	8192	tar	/usr/bin/fdformat
100	15795	R	9680	3072	tar	/usr/bin/fdformat
100	15795	R	4096	8192	tar	/usr/bin/fgrep
100	15795	R	46896	6144	tar	/usr/bin/fgrep
100	15795	R	4112	8192	tar	/usr/bin/file

```
[...]
```

# File system I/O via Virtual Memory

- File system I/O is performed by the VM system
  - Reads are performed by page-in
  - Write are performed by page-out
- Practical Implications
  - Virtual memory caches files, cache is dynamic
  - Minimum I/O size is the page size
  - Read/modify/write may occur on sub page-size writes
- Memory Allocation Policy:
  - File system cache is lower priority than app, kernel etc
  - File system cache grows when there is free memory available
  - File system cache shrinks when there is demand elsewhere.

## File System Reads: A UFS Read

- Application calls `read()`
- Read system call calls `fop_read()`
- FOP layer redirector calls underlying filesystem
- FOP jumps into `ufs_read`
- UFS locates a mapping for the corresponding pages in the file system page cache using `vnode/offset`
- UFS asks segmap for a mapping to the pages
- If the page exists in the fs, data is copied to App.
  - We're done.
- If the page doesn't exist, a Major fault occurs
  - VM system invokes `ufs_getpage()`
  - UFS schedules a page size I/O for the page
  - When I/O is complete, data is copied to App.

# vmstat -p

swap = free and unreserved swap in KBytes

free = free memory measured in pages

re = kilobytes reclaimed from cache/free list

mf = minor faults - the page was in memory but was not mapped

fr = kilobytes that have been destroyed or freed

de = kilobytes freed after writes

sr = kilobytes scanned / second

executable pages: kilobytes in - out - freed

anonymous pages: kilobytes in - out  
- freed

file system pages:  
kilobytes in - out -  
freed

memory		page					executable			anonymous			filesystem		
swap	free	re	mf	fr	de	sr	epi	epo	epf	api	apo	apf	fpi	fpo	fpf
...															
46715224	891296	24	350	0	0	0	0	0	0	4	0	0	27	0	0
46304792	897312	151	761	25	0	0	17	0	0	1	0	0	280	25	25
45886168	899808	118	339	1	0	0	3	0	0	1	0	0	641	1	1
46723376	899440	29	197	0	0	0	0	0	0	40	0	0	60	0	0



# Observing the File System I/O Path

```
sol10# cd labs/fs_paging
sol10# ./fsread
2055: 0.004: Random Read Version 1.8 05/02/17 IO personality successfully loaded
2055: 0.004: Creating/pre-allocating files
2055: 0.008: Waiting for preallocation threads to complete...
2055: 28.949: Pre-allocated file /filebench/bigfile0
2055: 30.417: Starting 1 rand-read instances
2056: 31.425: Starting 1 rand-thread threads
2055: 34.435: Running for 600 seconds...

sol10# vmstat -p 3
      memory          page          executable          anonymous          filesystem
    swap   free   mf   fr   de   sr   epi   epo   epf   api   apo   apf   fpi   fpo   fpf
1057528 523080 22 105   0   0    8     5     0     0     0     0     0     0     63     0     0
776904 197472   0  12   0   0    0     0     0     0     0     0     0     0    559     0     0
776904 195752   0   0   0   0    0     0     0     0     0     0     0     0    555     0     0
776904 194100   0   0   0   0    0     0     0     0     0     0     0     0    573     0     0

sol10# ./pagingflow.d
0  => pread64                                0
0  |  pageio_setup:pgin                      40
0  |  pageio_setup:pgpgin                    42
0  |  pageio_setup:maj_fault                  43
0  |  pageio_setup:fspgin                     45
0  |  bdev_strategy:start                   52
0  |  biodone:done                           11599
0  <= pread64                                11626
```



# Observing File System I/O

```
sol10# cd labs/fs_paging  
sol10# ./fsread  
2055: 0.004: Random Read Version 1.8 05/02/17 IO personality successfully loaded  
2055: 0.004: Creating/pre-allocating files  
2055: 0.008: Waiting for preallocation threads to complete...  
2055: 28.949: Pre-allocated file /filebench/bigfile0  
2055: 30.417: Starting 1 rand-read instances  
2056: 31.425: Starting 1 rand-thread threads  
2055: 34.435: Running for 600 seconds...
```

```
sol10# ./fspaging.d
```

Event	Device	Path	RW	Size
get-page		/filebench/bigfile0		8192
getpage-io	cmdk0	/filebench/bigfile0	R	8192
get-page		/filebench/bigfile0		8192
getpage-io	cmdk0	/filebench/bigfile0	R	8192
get-page		/filebench/bigfile0		8192
getpage-io	cmdk0	/filebench/bigfile0	R	8192
get-page		/filebench/bigfile0		8192



# Observing File System I/O: Sync Writes

```
Sol10# cd labs/fs_paging  
sol10# ./fswritesync  
2276: 0.008: Random Write Version 1.8 05/02/17 IO personality successfully loaded  
2276: 0.009: Creating/pre-allocating files  
2276: 0.464: Waiting for preallocation threads to complete...  
2276: 0.464: Re-using file /filebench/bigfile0  
2276: 0.738: Starting 1 rand-write instances  
2277: 1.742: Starting 1 rand-thread threads  
2276: 4.743: Running for 600 seconds...
```

```
sol10# ./fspaging.d  
Event           Device          Path   RW    Size   Offset  
put-page        cmdk0          /filebench/bigfile0  8192  
putpage-io     cmdk0          /filebench/bigfile0  8192 18702224  
other-io        cmdk0          <none>      W    512   69219  
put-page        cmdk0          /filebench/bigfile0  8192  
putpage-io     cmdk0          /filebench/bigfile0  8192 11562912  
other-io        cmdk0          <none>      W    512   69220  
put-page        cmdk0          /filebench/bigfile0  8192  
putpage-io     cmdk0          /filebench/bigfile0  8192 10847040  
other-io        cmdk0          <none>      W    512   69221  
put-page        cmdk0          /filebench/bigfile0  8192  
putpage-io     cmdk0          /filebench/bigfile0  8192 22170752  
other-io        cmdk0          <none>      W    512   69222  
put-page        cmdk0          /filebench/bigfile0  8192  
putpage-io     cmdk0          /filebench/bigfile0  8192 25189616  
other-io        cmdk0          <none>      W    512   69223  
put-page        cmdk0          /filebench/bigfile0  8192
```



# fsinfo(1)

```
# fsstat ufs 1
```

new	name	name	attr	attr	lookup	rddir	read	read	write	write	
file	remov	chng	get	set	ops	ops	ops	bytes	ops	bytes	
96.0K	91.7K	48	10.6M	46.0K	40.6M	429K	362K	57.2G	8.46M	67.3G	ufs
0	0	0	0	0	0	0	0	0	3	352	ufs
0	0	0	0	0	0	0	0	0	2	176	ufs
0	0	0	0	0	0	0	0	0	2	176	ufs
0	0	0	0	0	0	0	0	0	2	176	ufs
0	0	0	5	0	4	0	1	2.47K	3	240	ufs
0	0	0	0	0	0	0	0	0	2	176	ufs
0	0	0	0	0	9	0	0	0	2	176	ufs
0	0	0	58	0	505	4	0	0	2	176	ufs



# Memory Mapped I/O

- Application maps file into process with `mmap()`
- Application references memory mapping
- If the page exists in the cache, we're done.
- If the page doesn't exist, a Major fault occurs
  - VM system invokes `ufs_getpage()`
  - UFS schedules a page size I/O for the page
  - When I/O is complete, data is copied to App.



USE



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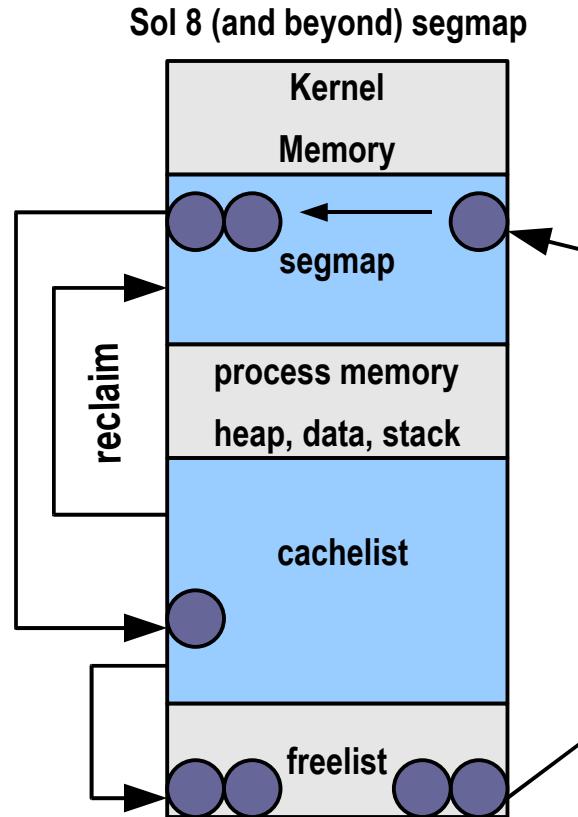
EVANGELIZE

# Optimizing Random I/O File System Performance

# Random I/O

- Attempt to cache as much as possible
  - The best I/O is the one you don't have to do
  - Eliminate physical I/O
  - Add more RAM to expand caches
  - Cache at the highest level
    - Cache in app if we can
    - In Oracle if possible
- Match common I/O size to FS block size
  - e.g. Write 2k on 8k FS = Read 8k, Write 8k

# The Solaris UFS Cache



# Tuning segmap (UFS L1 cache)

- By default, on SPARC, segmap is sized at 12% of physical memory
  - Effectively sets the minimum amount of file system cache on the system by caching in segmap over and above the dynamically-sized cachelist
- On Solaris 8/9
  - If the system memory is used primarily as a cache, cross calls (mpstat xcall) can be reduced by increasing the size of segmap via the system parameter segmap\_percent (12 by default)
  - segmap\_percent = 100 is like Solaris 7 without priority paging, and will cause a paging storm
  - Must keep segmap\_percent at a reasonable value to prevent paging pressure on applications e.g. 50%
  - segkpm in Solaris 10 and OpenSolaris
- On Solaris 10 on X64, segmap is 64MB by default
  - Tune with segmapsize in /etc/system or eeprom
    - set segmapsize = 1073741824 (1 GB)
  - On 32-bit X64, max segmapsize is 128MB



# Tuning segmap\_percent

- There are kstat statistics for segmap hit rates
  - Estimate hit rate as  $(\text{get\_reclaim} + \text{get\_use}) / \text{getmap}$

```
# kstat -n segmap
module: unix
name: segmap
instance: 0
class: vm

crttime          17.299814595
fault            17361
faulta           0
free             0
free_dirty       0
free_notfree     0
get_nofree       0
get_reclaim      67404
get_reuse        0
get_unused       0
get_use          83
getmap           71177
pagecreate       757
rel_abort        0
rel_async        3073
rel_dontneed    3072
rel_free         616
rel_write        2904
release          67658
snaptime         583596.778903492
```

# UFS Access times

- Access times are updated when file is accessed or modified
  - e.g. A web server reading files will storm the disk with atime writes!
- Options allow atimes to be eliminated or deferred
  - dfratime: defer atime write until write
  - noatime: do not update access times, great for web servers and databases

# Asynchronous I/O

- An API for single-threaded process to launch multiple outstanding I/Os
  - Multi-threaded programs could just have multiple threads
  - Oracle databases use this extensively
  - See `aio_read()`, `aio_write()` etc...
- Slightly different variants for RAW disk vs file system
  - UFS, NFS etc: libaio creates lwp's to handle requests via standard `pread/pwrite` system calls
  - RAW disk: I/Os are passed into kernel via `kaio()`, and then managed via task queues in the kernel
    - Moderately faster than user-level LWP emulation



# Key UFS Features

- Direct I/O
  - Solaris 2.6+
- Logging
  - Solaris 7+
- Async I/O
  - Oracle 7.x, -> 8.1.5 - Yes
  - 8.1.7, 9i - New Option
- Concurrent Write Direct I/O
  - Solaris 8, 2/01

# Database big rules...

- Always put re-do logs on Direct I/O
- Cache as much as possible in the SGA
- Use 64-Bit RDBMS
- Always use Asynch I/O
- Use Solaris 8 Concurrent Direct I/O
- Place as many tables as possible on Direct I/O, assuming SGA sized correct
- Place write-intensive tables on Direct I/O

# Sequential I/O

- Disk performance fundamentals
  - Disk seek latency will dominate for random I/O
    - ~5ms per seek
    - A typical disk will do ~200 I/Os per second random I/O
    - $200 \times 8k = 1.6\text{MB/s}$
    - Seekless transfers are typically capable of ~50MB/s
      - Requires I/O sizes of 64k+
  - Optimizing for sequential I/O
    - Maximizing I/O sizes
    - Eliminating seeks
    - Minimizing OS copies



## Sequential I/O – Looking at disks via iostat

- Use iostat to determine average I/O size
  - I/O size = kbytes/s divided by I/Os per second
- What is the I/O size in our example?
  - $38015 / 687 = 56k$
  - Too small for best sequential performance

```
# iostat -xnz
                         extended device statistics
      r/s      w/s      kr/s      kw/s wait  actv wsvc_t asvc_t  %w  %b device
 687.8      0.0  38015.3      0.0    0.0   1.9      0.0     2.7      0 100 c0d0
```

# Sequential I/O – Maximizing I/O Sizes

- Application
  - Ensure application is issuing large writes
    - 1MB is a good starting point
    - truss or dtrace app
- File System
  - Ensure file system groups I/Os and does read ahead
  - A well tuned fs will group small app I/Os into large Physical I/Os
  - e.g. UFS cluster size
- IO Framework
  - Ensure large I/O's can pass though
  - System param *maxphys* set largest I/O size
- Volume Manager
  - *md\_maxphys* for SVM, or equiv for Veritas
- SCSI or ATA drivers often set defaults to upper layers



# Sequential on UFS

- Sequential mode is detected by 2 adjacent operations
  - e.g read 8k, read8k
- UFS uses “clusters” to group reads/write
  - UFS “maxcontig” param, units are 8k
  - Maxcontig becomes the I/O size for sequential
  - Cluster size defaults to 1MB on Sun FCAL
    - 56k on x86, 128k on SCSI
    - Auto-detected from SCSI driver's default
      - Set by default at newfs time (can be overridden)
  - e.g. Set cluster to 1MB for optimal sequential perf...
  - Check size with “mkfs -m”, set with “tunefs -a”

```
# mkfs -m /dev/dsk/c0d0s0
mkfs -F ufs -o nsect=63,ntrack=32,bsize=8192,fragsize=1024,cgszie=49,free=1,rps=60,
nbpi=8143,opt=t,apc=0,gap=0,nrpos=8,maxcontig=7,mtb=n /dev/dsk/c0d0s0 14680512

# tunefs -a 128 /dev/rdsk/...
```

# Sequential on UFS

- Cluster Read
  - When sequential detected, read ahead entire cluster
  - Subsequent reads will hit in cache
  - Sequential blocks will not pollute cache by default
    - i.e. Sequential reads will be freed sooner
    - Sequential reads go to head of cachelist by default
    - Set system param *cache\_read\_ahead*=1 if all reads should be cached
- Cluster Write
  - When sequential detected, writes are deferred until cluster is full



# UFS write throttle

- UFS will block when there are too many pending dirty pages
  - Application writes by default go to memory, and are written asynchronously
  - Throttle blocks to prevent filling memory with async. Writes
- Solaris 8 Defaults
  - Block when 384k of unwritten cache
    - Set *ufs\_HW=<bytes>*
  - Resume when 256k of unwritten cache
    - Set *ufs\_IW=<bytes>*
- Solaris 9+ Defaults
  - Block when >16MB of unwritten cache
  - Resume when <8MB of unwritten cache



# Direct I/O

- Introduced in Solaris 2.6
- Bypasses page cache
  - Zero copy: DMA from controller to user buffer
- Eliminate any paging interaction
  - No 8k block size I/O restriction
  - I/Os can be any multiple of 512 bytes
  - Avoids write breakup of O\_SYNC writes
- But
  - No caching! Avoid unless application caches
  - No read ahead – application must do it's own
- Works on multiple file systems
  - UFS, NFS, VxFS, QFS



# Direct I/O

- Enabling direct I/O
  - Direct I/O is a global setting, per file or filesystem
  - Mount option

```
# mount -o forcedirectio /dev/dsk... /mnt
```

- Library call

```
directio(fd, DIRECTIO_ON | DIRECTIO_OFF)
```

- Some applications can call directio(3c)
  - e.g. Oracle – see later slides



# Enabling Direct I/O

- Monitoring Direct I/O via directiostat
  - See <http://www.solarisinternals.com/tools>

```
# directiostat 3
lreads lwrites  preads pwrites      Krd      Kwr holdrds nflush
    0      0      0      0      0      0      0      0
    0      0      0      0      0      0      0      0
    0      0      0      0      0      0      0      0
```

lreads = logical reads to the UFS via directio

lwrites = logical writes to the UFS via directio

preads = physical reads to media

pwrites = physical writes to media

Krd = kilobytes read

Kwr = kilobytes written

nflush = number of cached pages flushed

holdrds = number of times the read was a "hole" in the file.

# Using Direct I/O

- Enable per-mount point is the simplest option
- Remember, it's a system-wide setting
- Use sparingly, only applications which don't want caching will benefit
  - It disables caching, read ahead, write behind
  - e.g. Databases that have their own cache
  - e.g. Streaming high bandwidth in/out
- Check the side effects
  - Even though some applications can benefit, it may have side affects for others using the same files
    - e.g. Broken backup utils doing small I/O's will hurt due to lack of prefetch



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# ZFS

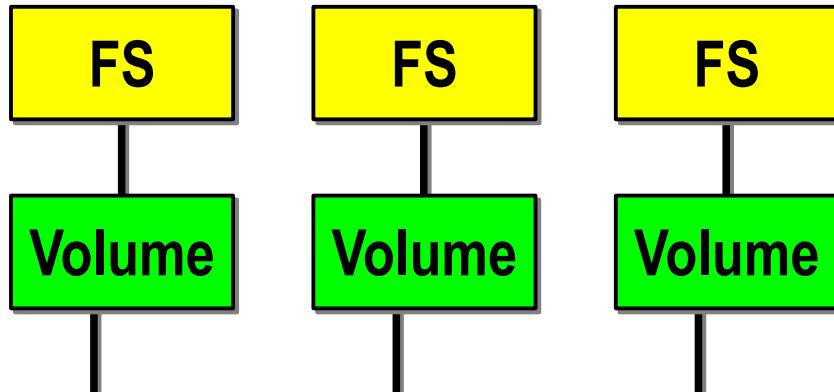
# ZFS

- Started from scratch with today's problems in mind
- Pooled Storage
  - Do for storage what VM does for RAM
- End-to-End Data integrity
  - Block-level checksum
  - Self-correcting when redundant data available
  - No more silent data corruption
- Transaction Model
  - COW updates – no changes to on-disk data
  - FS on-disk integrity maintained
  - Many opportunities for performance optimizations (IO scheduler and transaction reordering)
- Massive Scale – 128 bits

# FS/Volume Model vs. Pooled Storage

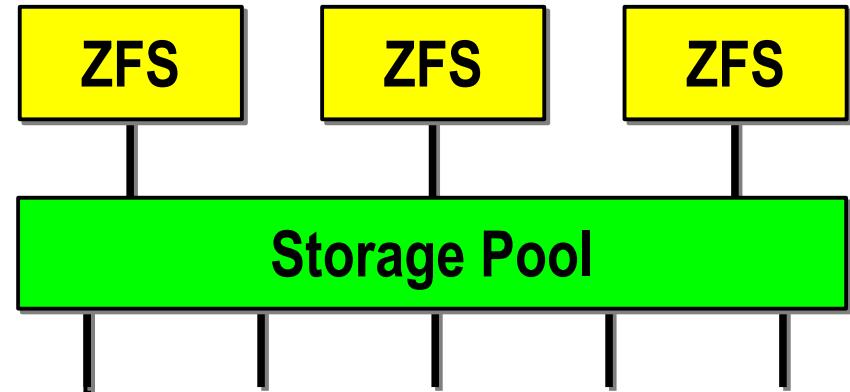
## Traditional Volumes

- Abstraction: virtual disk
- Partition/volume for each FS
- Grow/shrink by hand
- Each FS has limited bandwidth
- Storage is fragmented, stranded



## ZFS Pooled Storage

- Abstraction: malloc/free
- No partitions to manage
- Grow/shrink automatically
- All bandwidth always available
- All storage in the pool is shared



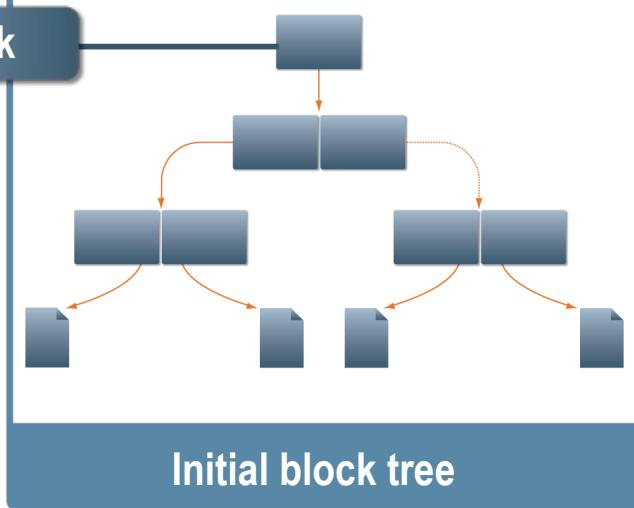


## ZFS Data Integrity Model

- Copy-on-write, transactional design
- Everything is checksummed
- RAID-Z/Mirroring protection
- Ditto Blocks
- Disk Scrubbing
- Write Failure Handling

# Copy-on-Write and Transactional

Uber-block



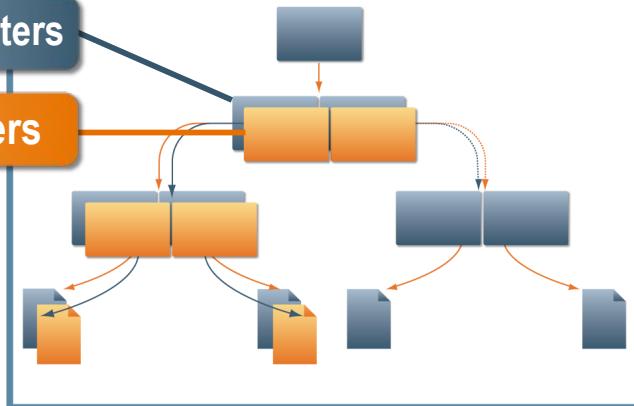
Original Data

New Data

Writes a copy of some changes

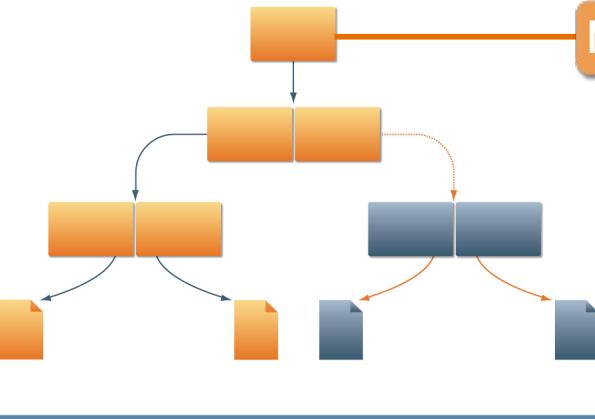
Original Pointers

New Pointers



New Uber-block

Rewrites the Uber-block

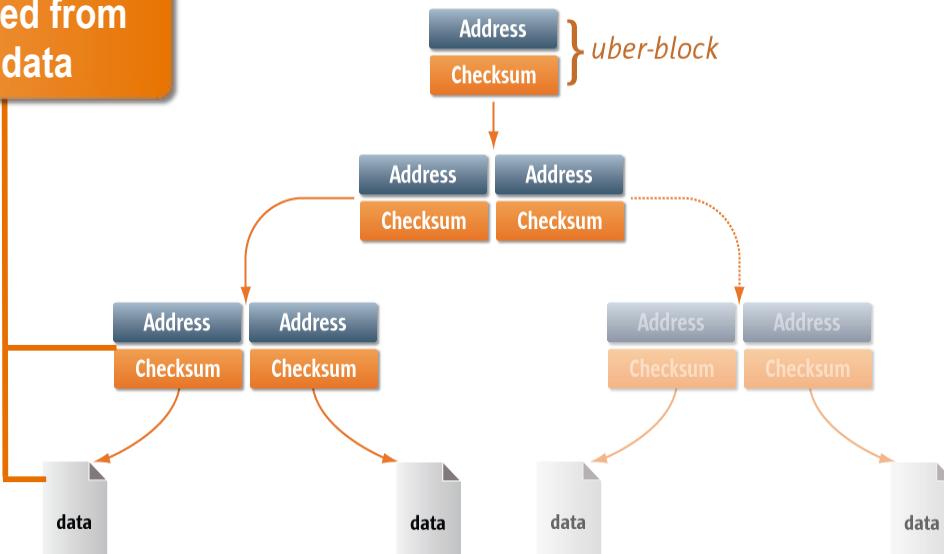


# Measurements at CERN

- Wrote a simple application to write/verify 1GB file
  - Write 1MB, sleep 1 second, etc. until 1GB has been written
  - Read 1MB, verify, sleep 1 second, etc.
- Ran on 3000 rack servers with HW RAID card
- After 3 weeks, found 152 instances of silent data corruption
  - Previously thought “everything was fine”
- HW RAID only detected “noisy” data errors
- Need end-to-end verification to catch silent data corruption

# End-to-End Checksums

Checksums are separated from the data



Entire I/O path is self-validating (*uber-block*)

## Prevents:

- > Silent data corruption
- > Panics from corrupted metadata
- > Phantom writes
- > Misdirected reads and writes
- > DMA parity errors
- > Errors from driver bugs
- > Accidental overwrites

# Disk Scrubbing

- Uses checksums to verify the integrity of all the data
- Traverses metadata to read every copy of every block
- Finds latent errors while they're still correctable
- It's like ECC memory scrubbing – but for disks
- Provides fast and reliable re-silvering of mirrors



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Copy-on-Write  
Design  
Multiple Block Sizes  
Pipelined I/O  
Dynamic Striping  
Intelligent Pre-Fetch  
Separate Intent Log  
Dedicated cache  
device

## Architected for Speed

# Variable Block Size

- No single block size is optimal for everything
  - Large blocks: less metadata, higher bandwidth
  - Small blocks: more space-efficient for small objects
  - Record-structured files (e.g. databases) have natural granularity;  
filesystem must match it to avoid read/modify/write
- Why not arbitrary extents?
  - Extents don't COW or checksum nicely (too big)
  - Large blocks suffice to run disks at platter speed
- Per-object granularity
  - A 37k file consumes 37k – no wasted space
- Enables transparent block-based compression

# ZFS Intent Log (ZIL)

- Satisfies synchronous writes semantics
  - O\_SYNC/O\_DSYNC
- Blocks are allocated from the main pool
  - Guaranteed to be written to stable storage before system call returns
- Examples:
  - Database often utilize synchronous writes to ensure transactions are on stable storage
  - NFS and other applications can issue fsync() to commit prior to writes

# Separate Intent Log (slog)

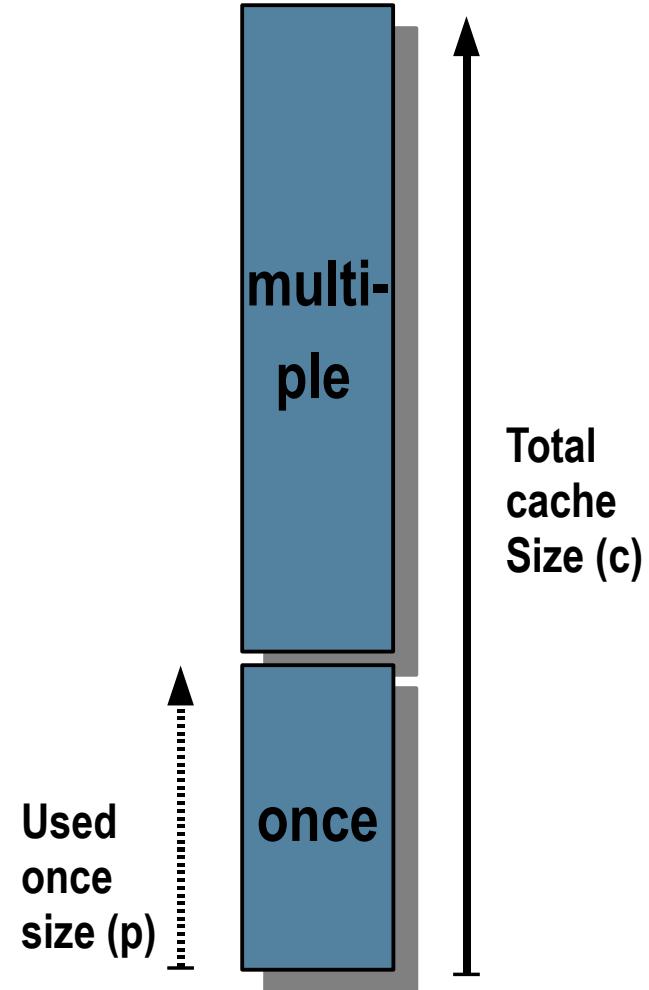
- Leverages high speed devices for dedicated intent log processing
  - Low latency devices such as SSDs (aka Logzilla)
- Can be mirrored and striped
- Blocks are allocated from dedicated log device
  - Failure reverts back to general pool

**Example: Create a pool with a dedicated log device**

```
# zpool create tank mirror c0d0 c1d0 log c2d0
```

# Adaptive Replacement Cache (ARC)

- Scan-resistant LRU (least recently used)
- Cache size divided into two:
  - Used once
  - Used multiple times
- Automatically adjust to memory pressure and workload
  - Data which is not being referenced is evicted
  - Ratio of once/multiple adjust dynamically based on workload



# L2ARC – cache device



- Provides a level of caching between main memory and disk
  - Utilizes specialized read-biased SSDs to extend the cache (aka “Readzilla”)
- Asynchronously populates the cache
  - Moves blocks from the ARC to L2ARC cache device

**Example: Create a pool with a cache device**

```
# zpool create tank mirror c0d0 c1d0 cache c2d0
```

# Typical way to Improve Performance

- Buy lots of RAM
  - Cache as much as possible
  - Use DRAM to compensate for slower disks
- Use lots of spindles
  - Spread the load across as many devices as possible
  - Use the outer most cylinders of the disk (make sure the disks don't seek)
  - Use NVRAM
- Throw \$\$\$ at the problem



# How to get terrible performance

- Run against storage array that flush caches
- Run simple benchmarks without decoding the numbers
  - compare write to cache vs write to disk
- Run the pool at 95% disk full
- Do random reads from widest raid-z
- Run a very large DB without tuning the recordsize
- Don't provision enough CPU
- Don't configure swap space
- Don't read the ZFS Best Practices

# How to get Great performance

- small files (<128K)
  - ufs allocates 1 inode per MB
  - netapps 1 / 32K
  - ZFS uses 1.2K to store 1K files !!!
  - Create 10s of files per single I/Os
  - \$ miss reads == single disk I/O
- ZFS does constant time snapshot
  - it's basically a noop to take a snapshot
  - snap deletions proportional to changes
  - snapshots helps simplify your business

# How to get Great performance

- Run ZFS in the storage back end (7000 Storage)
- Or provision for CPU usage.
- Configure enough RPM
  - 2 Mirrored 7.2 K RPM vs 1 x 15 K RPM in Raid-5
- Move Spindle Constrained setup to ZFS
  - write streaming + I/O aggregation
    - efficient use of spindles on writes,
    - 100% full stripes in storage
  - free spindles for reads
  - use a separate intent log (NVRAM or SSD or just N separate spindles) for an extra boost

# Solaris 10 Update 6

- Finally got write throttling, ZFS won't eat all of memory
  - Grows and shrink dance now as designed
  - Capping the ARC seems commonly done
  - ZFS reports accurate freemem, others cache data in freemem
- Cache flushes to SAN array partially solved
  - HDS, EMC with recent firmware are ok.
  - Can be tuned per array
  - Others ? set `zfs_nocacheflush` (cf evil tuning guide)
- Vdev level prefetching is auto tuning
  - no problems there

# Solaris 10 Update 6

- We have the separate intent log
  - one or a few disks, but preferably SSD or NVRAM/DRAM device

## Upcoming

- L2 ARC
  - on/off per dataset
- ARC
  - on/off per dataset, ~directio
- Storage 7000
  - Tracks Nevada



# Tuning is Evil

- Leave a trace, explain motivation
  - zfs\_nocacheflush (on storage arrays that do)
  - capping the ARC (to preserve large pages)
  - zfs\_prefetch\_disable (zfetch consuming cpus)
  - zfs\_vdev\_max\_pending (default 35, 10-16 for DB)
  - zil\_disable (NO!!! don't or face application corruptions)
- No tuning required
  - vdev prefetch (issue now fixed)

# ZFS Best Practices

- Tune recordsize only on fixed records DB files
- Mirror for performance
- 64-bit kernel (allows greater ZFS caches)
- configure swap (don't be scared by low memory)
- Don't slice up devices (confuses I/O scheduler)
- For raid-z[2] : don't go two wide (for random reads)
- Isolate DB log writer if that is critical (use few devices)
- Separate Root pool (system's identify) and data pools (system's function)

# ZFS Best Practices

- Don't mix legacy and non legacy shares (it's confusing)
- 1 FS per user (1 quota/reserv; user quota are coming)
- Rolling Snapshots (smf service)
- Instruct backup tool to skip .zfs
- Keep pool below 80% full (helps COW)



# MySQL Best Practices

- Match Recordsize with DB (16K)
- Use a separate intent log device within main zpool
- Find creative use of Snapshot/Clones send/recv
  - backups
  - master & slave architecture
- Use the ARC and L2ARC instead of disk RPM
  - a caching 7000 series serving masters & slaves
- NFS Directio and Jumbo Frames
  - save CPU cycles and memory for application
- Set innodb\_doublewrite=0
- Linux
  - innodb\_flush\_method = O\_DIRECT
  - echo noop > /sys/block/sde/queue/scheduler

# Tuning & Best Practices

- Tuning and BP wikis
  - [http://www.solarisinternals.com/wiki/index.php/ZFS\\_Best\\_Practices\\_Guide](http://www.solarisinternals.com/wiki/index.php/ZFS_Best_Practices_Guide)
  - [http://www.solarisinternals.com/wiki/index.php/ZFS\\_Evil\\_Tuning\\_Guide](http://www.solarisinternals.com/wiki/index.php/ZFS_Evil_Tuning_Guide)
  - [http://www.solarisinternals.com/wiki/index.php/ZFS\\_for\\_Databases](http://www.solarisinternals.com/wiki/index.php/ZFS_for_Databases)
  - <http://en.wikipedia.org/wiki/Zfs>
- ZFS Dynamics : In-Depth view
  - [http://blogs.sun.com/roch/entry/the\\_dynamics\\_of\\_zfs](http://blogs.sun.com/roch/entry/the_dynamics_of_zfs)
- Blue Prints
  - <http://wikis.sun.com/display/BluePrints/Main>
- Performance Savvy Bloggers
  - joyent (Ben Rockwood), Smugmug (Don Mcaskill), Neel (Oracle and MySQL), Media Temple



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# Networks

# Some Useful Numbers

- 1Gbit – theoretical 134MB/sec (megaBYTES)
- 10Gbit – theoretical 1.3GB/sec (gigaBYTES)

Bus	Technology	Bandwidth
PCI	32bit/33Mhz	133MB/sec
PCI	32bit/66Mhz	266MB/sec
PCI	64bit/66Mhz	530MB/sec
PCI-X	64bit/133Mhz	1GB/sec
PCI Express	1 lane, v1.X	250MB/sec
PCI Express	2 lanes, v1.X	500MB/sec
PCI Express	4 lanes, v1.X	1GB/sec
PCI Express	8 lanes, v1.x	2GB/sec (1.7GB/sec)

# Networks...

- Network performance and tuning is probably the most difficult to setup and measure
  - Hardware choices (NICs, PCI-Express attributes, platform-specific (SPARC versus X64) attributes), network switches
  - Software layers – TCP/IP stack, platform-specific device stack, platform-independent device stack...
  - OS releases (Solaris 10 updates versus NV)
    - software churn
  - Resource allocation – CPU to support load
  - Tuning methods - /etc/system and ndd(1M)
- Bandwidth is often the quoted performance metric
  - And it's important, but...
  - Many workloads care more about packets-per-second and latency



# NICs and Drivers

- The device name (ifconfig -a) is the driver
  - It's possible for multiple drivers to be available for the same hardware, i.e. configuring T2000 NICs with either e1000g or ipge (note: e1000g is better)



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# NICs and Drivers



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# NIC Tuneables



# Networks

- Key Observables
  - Link Utilization
  - Transmission, framing, checksum errors
  - Upstream software congestion
  - Routing
  - Over the wire latency
- What to measure
  - Link Bandwidth: nicstat
  - Link Speed: checkcable
  - Dropped upstream packets (nocanput)

## Networking - Tools

- netstat – kstat based, packet rates, errors, etc
- kstat – raw counters for NICs and TCP/UDP/IP
- nx.se – SE toolkit utility for bandwidth
- nicstat – NIC utilization
- snmpnetstat – network stats from SNMP
- checkcable – NIC status
- ping – host status
- traceroute – path to host, latency and hops
- snoop – network packets
- TTCP – workload generator
- pathchar – path to host analysis
- ntop – network traffic sniffer
- tcptop – DTrace tool, per process network usage
- tcpsnoop – DTrace tool, network packets by process
- dtrace – TCP, UDP, IP, ICMP, NIC drivers, etc....



## netstat(1)

```
# netstat -i 1
```

input			output			input			(Total)			output		
packets	errs	bge0	packets	errs	colls	packets	errs	bge0	packets	errs	colls	packets	errs	colls
15381402	0		1280618	0	0	15384170	0		1283386	0	0			
160	0	177	0	0	0	160	0	177	0	0	0			
213	0	205	0	0	0	213	0	205	0	0	0			
138	0	159	0	0	0	138	0	159	0	0	0			
159	0	172	0	0	0	159	0	172	0	0	0			
215	0	213	0	0	0	215	0	213	0	0	0			

## kstat(1)

```
# kstat -p bge:0:bge0:*bytes64
bge:0:bge0:obytes64      969276250
bge:0:bge0:rbytes64      1917373531
#
```



# TCP

```
# netstat -s | grep Bytes
```

tcpOutDataSegs	=862707	tcpOutDataBytes	=879539866
tcpRetransSegs	= 1189	tcpRetransBytes	=1159401
tcpInAckSegs	=473250	tcpInAckBytes	=879385959
tcpInInorderSegs	=694607	tcpInInorderBytes	=623233594
tcpInUnorderSegs	= 3926	tcpInUnorderBytes	=4877730
tcpInDupSegs	= 187	tcpInDupBytes	= 75281
tcpInPartDupSegs	= 6	tcpInPartDupBytes	= 7320
tcpInPastWinSegs	= 0	tcpInPastWinBytes	= 0

```
# kstat -n tcp
```

module: tcp	instance: 0
name: tcp	class: mib2
activeOpens	4809
attemptFails	22
connTableSize	56
connTableSize6	84
crtime	237.364807266
...	



# nicstat

```
$ nicstat 1
```

Time	Int	rKb/s	wKb/s	rPk/s	wPk/s	rAvs	wAvs	%Util	Sat
12:33:04	hme0	1.51	4.84	7.26	10.32	213.03	480.04	0.05	0.00
12:33:05	hme0	0.20	0.26	3.00	3.00	68.67	90.00	0.00	0.00
12:33:06	hme0	0.14	0.26	2.00	3.00	73.00	90.00	0.00	0.00
12:33:07	hme0	0.14	0.52	2.00	6.00	73.00	88.00	0.01	0.00
12:33:08	hme0	0.24	0.36	3.00	4.00	81.33	92.00	0.00	0.00
12:33:09	hme0	2.20	1.77	16.00	18.00	140.62	100.72	0.03	0.00
12:33:10	hme0	0.49	0.58	8.00	9.00	63.25	66.00	0.01	0.00
12:33:11	hme0	12.16	1830.38	185.06	1326.42	67.26	1413.06	15.09	0.00
12:33:12	hme0	19.03	3094.19	292.88	2229.11	66.53	1421.40	25.50	0.00
12:33:13	hme0	19.55	3151.87	301.00	2270.98	66.50	1421.20	25.98	0.00
12:33:14	hme0	11.99	1471.67	161.07	1081.45	76.25	1393.49	12.15	0.00
12:33:15	hme0	0.14	0.26	2.00	3.00	73.00	90.00	0.00	0.00
12:33:16	hme0	0.14	0.26	2.00	3.00	73.00	90.00	0.00	0.00
12:33:17	hme0	0.14	0.26	2.00	3.00	73.00	90.00	0.00	0.00

<http://blogs.sun.com/tmc>

<http://www.brendangregg.com>



# tcptop

<http://www.brendangregg.com/dtrace.html#DTraceToolkit>

```
# tcptop -C 10
```

Sampling... Please wait.

```
2005 Jul 5 04:55:25, load: 1.11, TCPin: 2 Kb, TCPout: 110 Kb
```

UID	PID	LADDR	LPORT	FADDR	FPORT	SIZE	NAME
100	20876	192.168.1.5	36396	192.168.1.1	79	1160	finger
100	20875	192.168.1.5	36395	192.168.1.1	79	1160	finger
100	20878	192.168.1.5	36397	192.168.1.1	23	1303	telnet
100	20877	192.168.1.5	859	192.168.1.1	514	115712	rcp

```
2005 Jul 5 04:55:35, load: 1.10, TCPin: 0 Kb, TCPout: 0 Kb
```

UID	PID	LADDR	LPORT	FADDR	FPORT	SIZE	NAME
0	242	192.168.1.5	79	192.168.1.1	54220	272	inetd
0	20879	192.168.1.5	79	192.168.1.1	54220	714	

in.fingerd

[ ... ]



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# tcpdump

```
# tcpdump -d
```

UID	PID	LADDR	LPORT	DR	RADDR	RPORT	SIZE	CMD
100	20892	192.168.1.5	36398	->	192.168.1.1	79	54	finger
100	20892	192.168.1.5	36398	<-	192.168.1.1	79	66	finger
100	20892	192.168.1.5	36398	->	192.168.1.1	79	54	finger
100	20892	192.168.1.5	36398	->	192.168.1.1	79	56	finger
100	20892	192.168.1.5	36398	<-	192.168.1.1	79	54	finger
100	20892	192.168.1.5	36398	<-	192.168.1.1	79	606	finger
100	20892	192.168.1.5	36398	->	192.168.1.1	79	54	finger
100	20892	192.168.1.5	36398	<-	192.168.1.1	79	54	finger
100	20892	192.168.1.5	36398	->	192.168.1.1	79	54	finger
100	20892	192.168.1.5	36398	->	192.168.1.1	79	54	finger
0	242	192.168.1.5	23	<-	192.168.1.1	54224	54	inetd
0	242	192.168.1.5	23	->	192.168.1.1	54224	54	inetd

...



# dtrace

```
# dtrace -n 'fbt:ip::entry { @[probefunc] = count(); }'  
dtrace: description 'fbt:ip::entry' matched 1875 probes
```

^C

```
tcp_set_rto                                     2  
tcp_timeout_cancel                             2  
tcp_timer_free                                2  
tcp_wput_data                                 2  
ip_input                                       3  
ip_loopback_src_or_dst                         3  
ip_tcp_input                                  3  
ipcl_classify_v4                             3  
ire_cache_lookup                            3  
squeue_enter_chain                          3  
tcp_find_pktinfo                           3  
tcp_input                                    3
```



# dtrace

```
# dtrace -n 'fbt:bge::entry { @[probefunc] = count(); }'  
dtrace: description 'fbt:bge::entry' matched 164 probes  
^C  

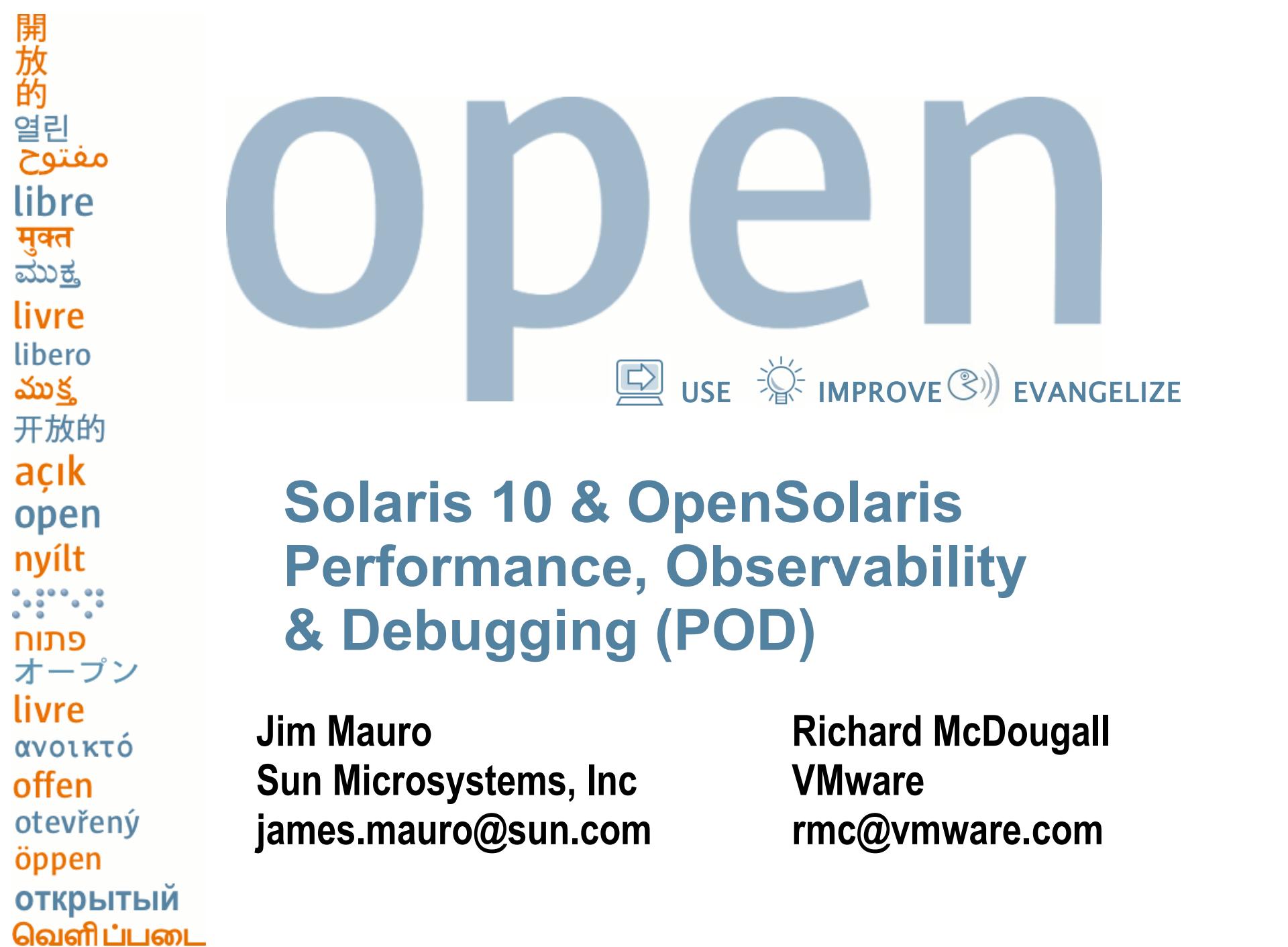

|                          |    |
|--------------------------|----|
| bge_atomic_renounce      | 1  |
| bge_atomic_claim         | 2  |
| bge_atomic_reserve       | 2  |
| bge_send                 | 2  |
| bge_m_tx                 | 3  |
| bge_atomic_shl32         | 6  |
| bge_chip_factotum        | 6  |
| bge_factotum_link_check  | 6  |
| bge_factotum_stall_check | 6  |
| bge_mbx_put              | 10 |
| bge_intr                 | 11 |
| bge_receive              | 11 |
| bge_recycle              | 11 |
| bge_chip_cyclic          | 12 |

  
...
```



# Summary

- Systems (hardware + software) are extremely complex
  - Understanding behavior requires a variety of skills
- Experience is your friend
  - Dive in!
- The Good News
  - Solaris/OpenSolaris is SECOND TO NONE when it comes to tools, utilities and observability
    - ...and some other things to
    - With few exceptions, typing the wrong thing won't hurt
      - exceptions include truss(1), dtrace with thousands of probes, dtrace PID provider with tens-of-thousands of probes
- With Solaris 10 and OpenSolaris, All questions can and will be answered
  - The Dark Ages are behind us...
  - Come to the light...



# open



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## Solaris 10 & OpenSolaris Performance, Observability & Debugging (POD)

**Jim Mauro**  
Sun Microsystems, Inc  
[james.mauro@sun.com](mailto:james.mauro@sun.com)

**Richard McDougall**  
VMware  
[rmc@vmware.com](mailto:rmc@vmware.com)

開放的  
열린  
مفتاح  
libre  
মুক্ত  
முक்த  
livre  
libero  
ମୁକ୍ତ  
开放的  
acık  
open  
nyílt  
ଓଡ଼ିଆ  
オープン  
livre  
ανοικτό  
offen  
otevřený  
öppen  
открытый  
வெளிப்படை



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# Supplemental Material



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# Memory Corruption & Memory Leak Detection



# The Problem

- Memory leaks
  - allocated memory is not freed when the program no longer needs it
  - Process address space size continues to grow
    - 32-bit processes can run out of address space
    - 64-bit processes can consume a LOT of memory
- Memory corruption
  - Overwriting a segment boundary
  - Unsynchronized writes into a mapped segment

# Memory Leak/Corruption Tools

- DTrace
  - Not the best tool for this job
  - pid provider can instrument malloc/free calls
    - track addresses and call stacks
- dbx
  - Has facilities for memory access errors and memory leak detection
- watchmalloc(3MALLOC)
  - binary replacement for libc malloc, free, etc
  - Environmental variable settings for debugging
- libumem(3LIB)
  - Extended feature set for debugging memory leaks and corruption
  - Used in conjunction with mdb(1)



# Memory Leak - DTrace

- Use the pid provider to instrument malloc/free calls
  - Post process to align mallocs & frees
- Grab a stack frame

```
#!/usr/sbin/dtrace -s
pid$target:$1:malloc:entry
{
    ustack();
}
pid$target:$1:malloc:return
{
    printf("%s: %x\n", probefunc, arg1);
}
pid$target:$1:free:entry
{
    printf("%s: %x\n", probefunc, arg0);
}
```



dbx

26 i = j;

(dbx) cont

hello world

Checking for memory leaks...

Actual leaks report (actual leaks: 1 total size: 32 bytes)

Memory Leak (mel):

Found leaked block of size 32 bytes at address 0x100101768

At time of allocation, the call stack was:

[1] memory\_leak() at line 19 in "hello.c"

[2] main() at line 31 in "hello.c"

Possible leaks report (possible leaks: 0 total size: 0 bytes)

Checking for memory use...

Blocks in use report (blocks in use: 1 total size: 12 bytes)

Block in use (biu):

Found block of size 12 bytes at address 0x100103778 (100.00% of

At time of allocation, the call stack was:

[1] memory\_use() at line 11 in "hello.c"

[2] main() at line 32 in "hello.c"

execution completed, exit code is 0

# watchmalloc(3MALLOC)

- Binary replacement for malloc, etc
  - Access via LD\_PRELOAD=watchmalloc.so.1
- Uses watchpoint facility of procfs (/proc)
- Imposes some constraints
  - See the man page
- Enabled via environmental variables
  - MALLOC\_DEBUG=WATCH,RW,STOP
    - WATCH – WA\_WRITE SIGTRAP
    - RW – enables WA\_READ & WA\_WRITE
    - STOP – Stop process instead of SIGTRAP



## watchmalloc(3MALLOC)

buf[11] 8060df3  
buf[12] 8060df4  
buf[13] 8060df5  
buf[14] 8060df6  
buf[15] 8060df7  
buf[16] 8060df8

Trace/Breakpoint Trap (core dumped)

opensolaris> pstack core

core 'core' of 841: ./mover 8

08050b25 main (2, 80479c8, 80479d4, 80479bc) + c5

080509cd \_start (2, 8047ae8, 8047af0, 0, 8047af2, 8047b08) + 7d

opensolaris> dbx mover core

For information about new features see `help changes'

To remove this message, put `dbxenv suppress\_startup\_message 7.7' in your .dbxrc

Reading mover

core file header read successfully

Reading ld.so.1

Reading watchmalloc.so.1

Reading libc.so.1

program terminated by signal TRAP (write access watchpoint trap)

0x08050b25: main+0x00c5: movb %al,0x00000000(%edx)

(dbx)



# libumem.so

- Binary replacement for malloc, etc
  - enable via LD\_PRELOAD
- Designed after the kernel heap allocator
  - Scalable. Less lock contention with threaded apps
- Also enables debug features for memory leak/access issues
  - special buffer management scheme for detecting memory issues

metadata section	memory available to the application	0xbb	memory not available to the application	redzone	debug metadata section

[http://developers.sun.com/solaris/articles/libumem\\_library.html](http://developers.sun.com/solaris/articles/libumem_library.html)

## libumem.so

- Enable with LD\_PRELOAD
- Set UMEM\_DEBUG & UMEM\_LOGGING
  - UMEM\_DEBUG=default
  - UMEM\_LOGGING=transaction
- Records Thread ID, Timestamp and stack trace for every memory transaction
- Fills allocated & freed segments within the buffer with special patterns
  - Detect use of uninitialized data and previously freed buffers
- Check redzone
- Debug metadata for buffer audit information



## libumem.so

- Detects
  - Buffer overruns
  - Multiple frees
  - Use of uninitialized data
  - Use of freed buffers
- Leak detection
  - Most useful when **UMEM\_DEBUG** is set to at least **default**
  - Can use **mdb(1)**'s ::findLeaks dcmd to search for memory leaks in core files and running processes
  - Get leaked memory summary, breakdown by stack trace



USE



IMPROVE



EVANGELIZE

# SunStudio



# SunStudio 12

- Compilers
  - C, C++, FORTRAN
- Thread Analyzer
- Performance Analyzer
- NetBeans IDE
- Add-on performance libraries
- Misc Tools
  - dbx
  - lint
  - cscope
  - etc...

# Thread Analyzer

- Detects data races and deadlocks in a multithreaded application
  - Points to non-deterministic or incorrect execution
  - Bugs are notoriously difficult to detect by examination
  - Points out actual and potential deadlock situations
- Process
  - Instrument the code with `-xinstrument=datarace`
  - Detect runtime condition with `collect -r all` [or race, detection]
  - Use graphical analyzer to identify conflicts and critical regions

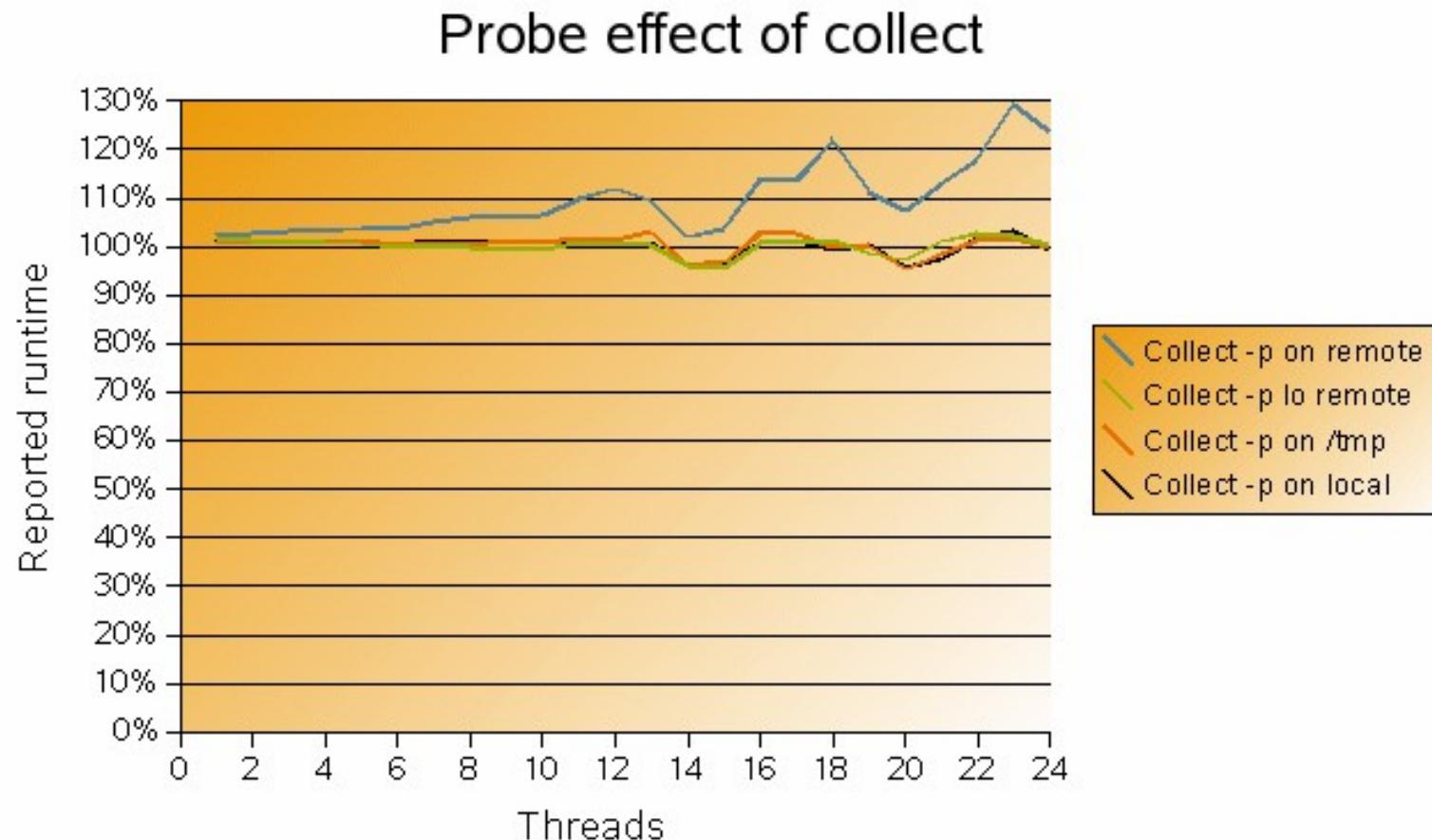
# Performance Analyzer

- Thread analyzer integrated into performance analyzer
  - Extentions to the .er files to accommodate THA data
  - collect command extensions
  - er\_print command extensions
- More extensive data collection
  - function, instruction count, dataspace profiling
  - attach to PID and collect data
- Probe effect can be mitigated
  - Reduce sampling rates when a lot of threads, or long-running collection

[http://blogs.sun.com/d/entry/analyzer\\_probe\\_effect](http://blogs.sun.com/d/entry/analyzer_probe_effect)

[http://blogs.sun.com/d/entry/analyzer\\_probe\\_effects\\_part\\_2](http://blogs.sun.com/d/entry/analyzer_probe_effects_part_2)

# Performance Analyzer Probe Effect



Graph courtesy of Darryl Gove, from [http://blogs.sun.com/d/entry/analyzer\\_probe\\_effect](http://blogs.sun.com/d/entry/analyzer_probe_effect)

# Performance Analyzer

- Collector Tool
  - collect(1)
  - profiles code and traces function calls
  - call stacks, microstates, hardware counters, memory allocation data, summary information
- Analyzer Tool
  - er\_print(1), analyzer(1)
  - Clock profiling metrics
  - Hardware counter metrics
  - Synchronization delay metrics
  - Memory allocation metrics
  - MPI tracing metrics

# Performance Analyzer

- Clock-based profiling
  - Thread state sampled/stored at regular intervals (SIGPROF)
  - Default resolution of 10milliseconds
  - High-res of 1ms possible
  - Low-res of 100ms for longer collections

User CPU time	LWP time spent running in user mode on the CPU.
Wall time LWP	time spent in LWP 1. This is usually the “wall clock time”
Total LWP time	Sum of all LWP times.
System CPU time	LWP time spent running in kernel mode on the CPU or in a trap state.
Wait CPU time	LWP time spent waiting for a CPU.
User lock time	LWP time spent waiting for a lock.
Text page fault time	LWP time spent waiting for a text page.
Data page fault time	LWP time spent waiting for a data page.
Other wait time	LWP time spent waiting for a kernel page, or time spent sleeping or stopped.



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# Performance Analyzer

- Hardware Counter Overflow Profiling Data
  -

# Hardware Counters

- PIC – Programmable Interval Counters
- Can be programmed to count hardware events (e.g. L2 cache miss, TLB miss, etc)
  - VERY processor-specific
  - Need to reference processor's PRM
- cpustat, cputrack
  - Solaris commands to setting PICs and tracking events
- Perf Analyzer & DTrace cpc provider
  - Use overflow profiling



## cpustat from a X4600

Use cputrack(1) to monitor per-process statistics.

CPU performance counter interface: AMD Opteron & Athlon64

event specification syntax:

[picn=]<eventn>[,attr[n][=<val>]]|[picn=]<eventn>[,attr[n][=<val>]],...]

Generic Events:

event[0-3]: PAPI\_br\_ins PAPI\_br\_msp PAPI\_br\_tkn PAPI\_fp\_ops  
PAPI\_fad\_ins PAPI\_fml\_ins PAPI\_fpu\_idl PAPI\_tot\_cyc  
PAPI\_tot\_ins PAPI\_I1\_dca PAPI\_I1\_dcm PAPI\_I1\_ldm  
PAPI\_I1\_stm PAPI\_I1\_ica PAPI\_I1\_icm PAPI\_I1\_icr  
PAPI\_I2\_dch PAPI\_I2\_dcm PAPI\_I2\_dcr PAPI\_I2\_dcw  
PAPI\_I2\_ich PAPI\_I2\_icm PAPI\_I2\_ldm PAPI\_I2\_stm  
PAPI\_res\_stl PAPI\_stl\_icy PAPI\_hw\_int PAPI\_tlb\_dm  
PAPI\_tlb\_im PAPI\_fp\_ins PAPI\_vec\_ins

See generic\_events(3CPC) for descriptions of these events

IC\_refill\_from\_L2 IC\_refill\_from\_system

IC\_itlb\_L1\_miss\_L2\_hit IC\_uarch\_resync\_snoop

IC\_instr\_fetch\_stall IC\_return\_stack\_hit

IC\_return\_stack\_overflow FR\_retired\_x86\_instr\_w\_excp\_intr

. . . . .

DC\_copyback DC\_dtib\_L1\_miss\_L2\_hit DC\_dtib\_L1\_miss\_L2\_miss

DC\_1bit\_ecc\_error\_found BU\_system\_read\_responses

BU\_quadwords\_written\_to\_system BU\_internal\_L2\_req

BU\_fill\_req\_missed\_L2 BU\_fill\_into\_L2

IC\_itlb\_L1\_miss\_L2\_miss FR\_retired\_fpu\_instr

NB\_mem\_ctrlr\_page\_access NB\_mem\_ctrlr\_page\_table\_overflow

NB\_mem\_ctrlr\_turnaround NB\_ECC\_errors NB\_sized\_commands

NB\_probe\_result NB\_gart\_events NB\_ht\_bus0\_bandwidth

NB\_ht\_bus1\_bandwidth NB\_ht\_bus2\_bandwidth NB\_sized\_blocks

NB\_cpu\_io\_to\_mem\_io NB\_cache\_block\_commands

attributes: edge pc inv cmask umask nouser sys

See Chapter 10 of the "BIOS and Kernel Developer's Guide for the  
AMD Athlon 64 and AMD Opteron Processors," AMD publication #26094



insts[/{0|1|2|3}],9999991 ('Instructions Executed', alias for FR\_retired\_x86\_instr\_w\_excp\_intr; events)  
ic[/{0|1|2|3}],100003 ('I\$ Refs', alias for IC\_fetch; events)  
icm[/{0|1|2|3}],100003 ('I\$ Misses', alias for IC\_miss; events)  
itlbh[/{0|1|2|3}],100003 ('ITLB Hits', alias for IC\_itlb\_L1\_miss\_L2\_hit; events)  
itlbt[/{0|1|2|3}],100003 ('ITLB Misses', alias for IC\_itlb\_L1\_miss\_L2\_miss; events)  
eci[/{0|1|2|3}],1000003 ('E\$ Instr. Refs', alias for BU\_internal\_L2\_req~umask=0x1; events)  
ecim[/{0|1|2|3}],10007 ('E\$ Instr. Misses', alias for BU\_fill\_req\_missed\_L2~umask=0x1; events)  
dc[/{0|1|2|3}],1000003 ('D\$ Refs', alias for DC\_access; load events)  
dcm[/{0|1|2|3}],100003 ('D\$ Misses', alias for DC\_miss; load events)  
dtlbh[/{0|1|2|3}],100003 ('DTLB Hits', alias for DC\_dtlb\_L1\_miss\_L2\_hit; load-store events)  
dtlbt[/{0|1|2|3}],100003 ('DTLB Misses', alias for DC\_dtlb\_L1\_miss\_L2\_miss; load-store events)  
ecd[/{0|1|2|3}],1000003 ('E\$ Data Refs', alias for BU\_internal\_L2\_req~umask=0x2; load-store events)  
ecd़[/{0|1|2|3}],10007 ('E\$ Data Misses', alias for BU\_fill\_req\_missed\_L2~umask=0x2; load-store events)  
fpadd[/{0|1|2|3}],1000003 ('FP Adds', alias for FP\_dispatched\_fpu\_ops~umask=0x1; events)  
fpmul[/{0|1|2|3}],1000003 ('FP Muls', alias for FP\_dispatched\_fpu\_ops~umask=0x2; events)  
fpu stall[/{0|1|2|3}],1000003 ('FPU Stall Cycles', alias for FR\_dispatch\_stall\_fpu\_full; CPU-cycles)  
memstall[/{0|1|2|3}],1000003 ('Memory Unit Stall Cycles', alias for FR\_dispatch\_stall\_ls\_full; CPU-cycles)

## Function Metrics

- Exclusive metrics – events inside the function itself, excluding calls to other functions
  - Use exclusive metrics to locate functions with high metric values
- Inclusive metrics – events inside the function and any functions it calls
  - Use inclusive metrics to determine which call sequence in your program was responsible for high metric values
- Attributed metrics – how much of an inclusive metric came from calls from/to another function; they attribute metrics to another function



# Using the Performance Analyzer...

```
# collect -p lo -d exper ./ldr 8 1 /zp/space
```

```
# collect -p lo -s all -d exper ./ldr 8 1 /zp/space
```

```
#collect -p lo -s all -t 10 -o synct.er -d exper ./ldr 8 1 /zp/space
```



# Run Time Checking (RTC)

- Detects memory access errors
- Detects memory leaks
- Collects data on memory use
- Works with all languages
- Works with multithreaded code
- Requires no recompiling, relinking or makefile changes